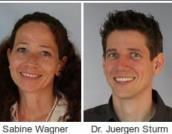
GPU Programming in Computer Vision

Introduction to Parallel Computing

Computer Vision Group









Mathieu Aubry

Prof. Dr. Daniel Cremers





Youngwook Kee

Christian Kerl

Maria Klodt



Julia Bergbauer

Quirin Lohr





Frank Steinbrücker Evgeny Strekalovskiy









RGB-D Sensors (Kinect)

Image Segmentation

3D Reconstruction from a Single Image

Convex Relaxation Methods



Jakob Engel





Jan Stühmer

Matthias Vestner Thomas Windheuser

Our Research is about

Optimization

Math in general



convex

- everything needs to broken down into functions, basic operations and numbers
- Numerics
 - continuous math on discrete hardware
- Programming (serial/parallel)
 - C/C++, CUDA, Matlab, ...
- Engineering

This Course covers

Parallel Programming (with CUDA)

- Computer Vision Basics
 - Image Filtering (Convolution, Diffusion)
 - Regularization (dealing with noise, unique solutions)

Optimization + Numerics

- Example Problems
 - Optical Flow Estimation
 - Superresolution

Course Goals

Learn how to program massively parallel processors and achieve

- High performance
- Functionality and maintainability
- Scalability across future generations
- Acquire technical knowledge required to achieve above goals
 - Principles and patterns of parallel programming
 - Processor architecture features and constraints
 - Programming API, tools and techniques
- Apply this knowledge to implement computer vision algorithms efficiently





Course Timeline

Aug. 26-30 (this week) : Lecture

- 4h lectures (attendance mandatory)
- programming exercises

Sep. 2-20: Student project

- optical flow/superresolution
- groups of 2-3 students
- unsupervised

Sep. 23-25: Presentations

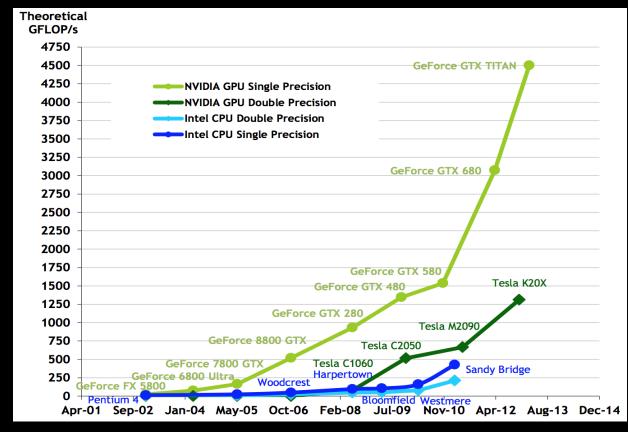
- 45 minutes presentation
- 25 minutes questions

September													
Mo.	Di. Mi. Do. Fr. Sa.												
26	27	28	29	30	31	1							
2	3	4	5	6	7	8							
9	10	11	12	13	14	15							
16	17	18	19	20	21	22							
23	24	25	26	27	28	29							
30	1	2	3	4	5	6							

Why Massively Parallel Processing?

A quiet revolution: Performance!

Computations: TFLOPs vs. 100 GFLOPs

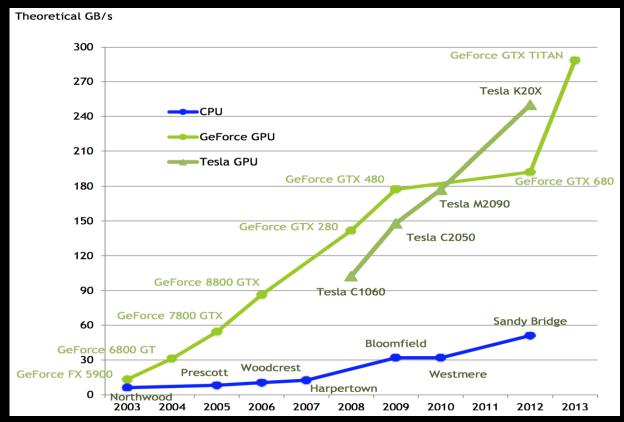


GPU in every PC – massive volume & impact

Why Massively Parallel Processing?

A quiet revolution: Performance!

Bandwidth: ~5x



GPU in every PC – massive volume & impact

Serial Performance Scaling is Over

Cannot continue to scale processor frequencies no 10 GHz chips

Cannot continue to increase power consumption can't melt chip

Can continue to increase transistor density as per Moore's Law

How to Use Transistors?

Larger caches ... decreasing

- Instruction-level parallelism ... decreasing
 - out-of-order execution, speculation, …
- Data-level parallelism ... increasing
 - vector units, SIMD execution, ...
 - Intel SSE, GPUs, ...
- Thread-level parallelism ... increasing
 - multithreading, multicore, manycore

Design difference: CPU vs. GPU

- Different goals produce different designs
 - CPU must be good at everything, parallel or not
 - GPU assumes work load is highly parallel
- CPU: minimize latency experienced by 1 thread
 - big on-chip caches
 - sophisticated control logic

GPU: maximize throughput of all threads

- skip big caches, multithreading hides latency
- share control logic across many threads, SIMD
- create and run thousands of threads

Design difference: CPU vs. GPU

minimize latency

Different goals produce different designs

- CPU must be good at everything, parallel or not
- GPU assumes work load is highly parallel

Control	ALU	ALU						
	ALU	ALU						
Cache								
Cuono								
DRAM		C	ORAM					
CF			-	βP				

maximize throughput

Enter the GPU

- Massively parallel
- Affordable supercomputing



NVIDIA GPUs

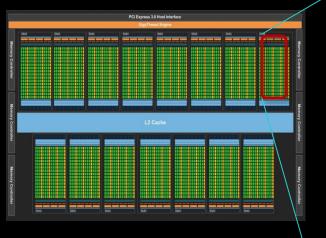
Compute Capability

- version number of the hardware architecture
- core architecture and incremental improvements

Arch	CC	GPUs	Features (e.g.)
	1.0	8800 GTX, Tesla C870	Basic functionality
Tesla	1.1	9800 GTX, Quadro FX 580	Atomics in global mem
(2007)	1.2	GT 240, Quadro FX 1800M	Atomics in shared mem
	1.3	GTX 285, Tesla C1060	Double precision
Fermi (2010)	2.0	GTX 480/580, Tesla C2070	Memory cache
	2.1	GTX 460, GTX 560 Ti	More cores (hardware)
Kepler	3.0	GTX 680/770, Tesla K10	Power efficiency, Many cores
(2012)	3.5	GTX 780/Titan, Tesla K20	Dynamic Parallelism, Hyper-Q

NVIDIA GPUs: Current

Kepler GPU



- 15 multiprocessors (up to)
- 192 Cuda Cores per SM
 - 2880 Cores in total (up to)

SMX																				
-	Instruction Cache Warp Scheduler Warp Scheduler Warp Scheduler Warp Scheduler																			
Di	Dispatch Dispatch			Dispatch Dispatch					Di	spatc	Dispatch Dispatch									
	+ +											Dispat								
	Register File (65,536 x 32-bit)																			
Core	Core Core Core DP Unit Core				Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DPU	Init	LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP U	Init	LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP U	Init	LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP U	Init	LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP U	Init	LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP U	Init	LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP U	Init	LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP U	Init	LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP U	Init	LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP U	Init	LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP U	Jnit	LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP U	Init	LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP U	Init	LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP U	Init	LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP U	Init	LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit		SFU			Core	DP Unit	Core	Core	Core	DP U	Init	LD/ST	SFU
	Interconnect Network 64 KB Shared Memory / L1 Cache																			
-	48 KB Read-Only Data Cache																			
	Tex		Tex		Tex Tex			(Tex Tex					Tex				Tex		
	Tex		Tex			Tex		Tex	٢		Tex		Tex		Tex				Tex	

Enter CUDA ("Compute Unified Device Architecture")

- Scalable parallel programming model
 - exposes the computational horsepower of GPUs
- Abstractions for parallel computing
 - Iet programmers focus on parallel algorithms
 - not mechanics of a parallel programming language
- Minimal extensions to familiar C/C++ environment to run code on the GPU
 - Low learning curve

CUDA: Scalable parallel programming

Provide straightforward mapping onto hardware

- good fit to GPU architecture
- maps well to multi-core CPUs too
- Execute code by many threads in parallel

Scale to 100s of cores & 10,000s of threads

- GPU threads are lightweight create / switch is free
- GPU needs 1000s of threads for full utilization

Outline of CUDA Basics

- Kernels and Thread Hierarchy
- Execution on the GPU
- Memory Management
- See the Programming Guide for the full API

BASIC KERNELS AND THREAD HIERARCHY

CUDA Definitions

Device: GPU

executes code in parallel

Host: CPU

manages execution on the device

Kernel: C/C++ function executed on the device

- executed by many threads
- each thread executes the same sequential program
- each thread is free to execute a unique code path

Quick Example

CPU: Process subtasks serially one by one:

```
for( int i=0; i<n; i++ )
{
    c[i] = a[i] + b[i];
}</pre>
```

GPU: Process each subtask in its own thread:

```
global___void vecAdd (float* a, float* b, float* c)
{
    int i = threadIdx.x + blockDim.x * blockIdx.x;
    c[i] = a[i] + b[i];
}
Each thread knows its index
```

Launch enough threads to cover all data

Thread Hierarchy

Kernel threads are grouped into blocks

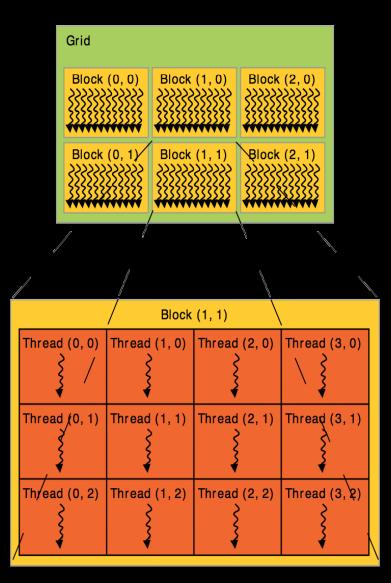
up to 512 (CC 1.x), 1024 (CC 2.x), or 2048 (CC 3.x) threads per block

Idea: Threads from the same block can cooperate

- synchronize their execution,
- communicate via shared memory
- threads from different blocks cannot cooperate
- Allows transparent scaling to different GPUs
- All kernel blocks together form a grid

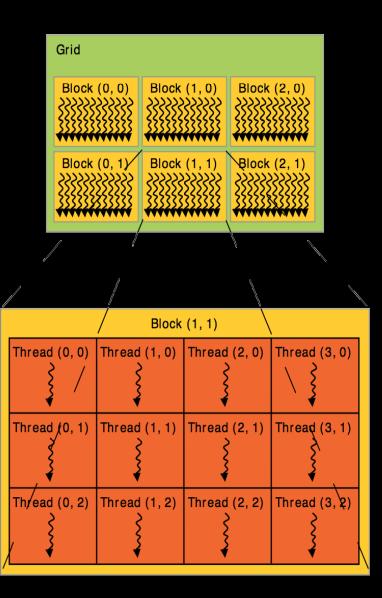
Thread Hierarchy

- # threads per block: up to 512 (CC 1.x), up to 1024 (CC 2.x), up to 2048 (CC 3.x)
- Blocks can be 1D, 2D, or 3D
- Grids can be 1D, 2D, or 3D
 - CC 1.x: only 1D or 2D
- Dimensions set at launch
 - Can be different for each grid



IDs and Dimensions

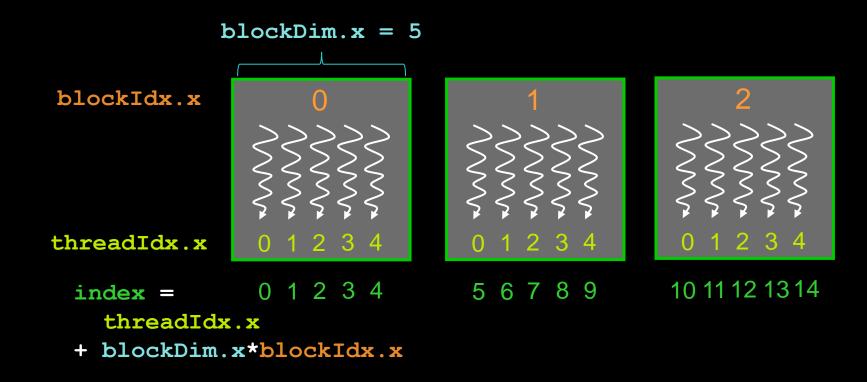
- Threads:
 - 3D IDs, unique within a block
- Blocks:
 - 3D IDs, unique within a grid
- Built-in variables:
 - threadIdx, blockIdx
 - blockDim, gridDim



Array Accesses: Indexing

Obtain unique array index from block/thread IDs

- threadIdx, blockIdx
- blockDim, gridDim



Kernel launch

Usual C/C++ function call, with an additional specification of grid and block sizes:

mykernel <<< gridSize, blockSize >>> (...);

- dim3 gridSize, dim blockSize
 - three int'S: blockSize.x, blockSize.y, blockSize.z
- Launched on the host side

CC 3.x: kernels can launch other kernels

Code executed on GPU: Restrictions

C/C++ with some restrictions

- Only access to GPU memory, cannot access CPU memory
 - (but access to "pinned" host memory, requires special allocation)
- No access to host functions
- No variable number of arguments
- No static variables

Code executed on GPU: Features

Many C/C++ features available on the GPU

- Templates
- Operator overloading
- Classes, inheritance
- Recursion (CC >=2.0)
- Function pointers (CC >= 2.0)
- new / delete (CC >= 2.0)
- Dynamic polymorphism, virtual functions (CC >= 2.0)
- Even printf() ! (CC >= 2.0)

Vector variants of basic types

- float2, float3, float4, double2, int4, char2, etc.
- float2 a = make_float2(1,2); a.x = 10; a.y = a.x;

Code executed on GPU: Specifiers

Special qualifiers to declare GPU functions:

- global__: kernels
 - Iaunched by CPU to run on the GPU
 - must return void
- device : auxiliary GPU functions
 - can only be called on the GPU
 - called from __global__ or __device__ functions
- host____: "normal" CPU C/C++ functions
 can only be called on the CPU

host and device qualifiers can be combined

Example: Vector Addition Kernel

```
Compute vector sum c = a + b
// Each thread performs one pair-wise addition
  global void vecAdd (float* a, float* b, float* c)
{
    int i = threadIdx.x + blockDim.x * blockIdx.x;
    c[i] = a[i] + b[i];
}
int main()
{
    • • •
    // Run grid of N/256 blocks of 256 threads each
    vecAdd <<< N/256, 256 >>> (d A, d B, d C);
}
```

Example: 2D Indexing

```
global void kernel (int *a, int dimx, int dimy)
{
    int x = threadIdx.x + blockDim.x * blockIdx.x;
    int y = threadIdx.y + blockDim.y * blockIdx.y;
    int ind = x + dim x + y;
    a[ind] = a[ind]+1;
}
int main()
{
    • • •
    dim3 \ block = dim3(32, 8, 1);
    dim3 grid = dim3( dimx/block.x, dimy/block.y, 1);
    kernel <<<grid,block>>> (d A, dimx, dimy);
}
```

Kernel Variations and Output

```
kernel<<<1,4>>>(d_a);
```

```
global void kernel( int *a )
    int idx = blockIdx.x*blockDim.x + threadIdx.x;
    a[idx] = 7;
                                          Output: 7777777777777777777777
}
 global void kernel( int *a )
    int idx = blockIdx.x*blockDim.x + threadIdx.x;
    a[idx] = blockIdx.x;
}
                                          Output: 0 0 0 0 1 1 1 1 2 2 2 2 3 3 3 3
 global
           void kernel( int *a )
ł
    int idx = blockIdx.x*blockDim.x + threadIdx.x;
    a[idx] = threadIdx.x;
                                          Output: 0 1 2 3 0 1 2 3 0 1 2 3 0 1 2 3
}
```

Blocks must be independent

- Any possible interleaving of blocks should be valid
 - presumed to run to completion without pre-emption
 - can run in any order (order is unspecified)
 - can run concurrently OR sequentially
- Blocks may coordinate but not synchronize
 - shared queue pointer: OK
 - shared lock: BAD ... can easily deadlock
- Independence requirement gives scalability

Execution of Kernels

Kernel launches are asynchronous w.r.t. CPU

- After kernel launch, control immediately returns
- CPU is free to do other stuff while the GPU is busy

Kernel launches are queued

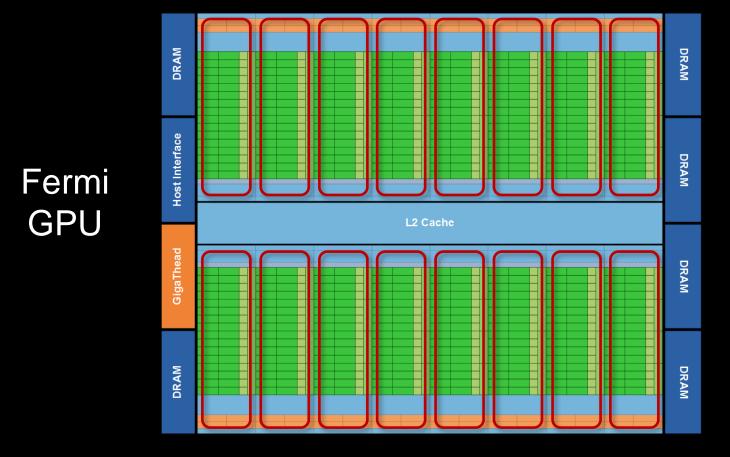
- Kernel doesn't start until previous kernels are finished
- Concurrent kernels possible for CC >= 2.0 (given enough resources)

Explicit synchronization if needed

cudaDeviceSynchronize()

EXECUTION ON GPU

NVIDIA GPU Architecture



- 16 independent multiprocessors (SMs)
- No shared resources except global memory
- No synchronization, always work in parallel

Single Fermi SM Multiprocessor

32 CUDA Cores per SM (512 total)

arithmetic/logic operations

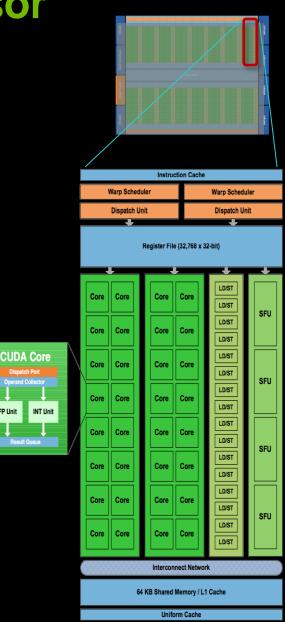
16 memory load/store units

(slow) access to off-chip GPU mem

4 Special Function Units 1/X, 1/SQRT(X), SIN, COS, EXP, ...

64 KB on-chip shared memory

- shared amongst CUDA cores
- enables thread communication



NVIDIA GPU Architecture: Current

Kepler GPU



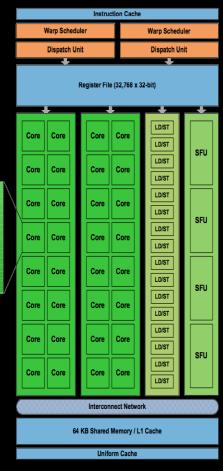
- 15 multiprocessors (up to)
- 192 Cuda Cores per SM
 - 2880 Cores in total (up to)

SMX																			
-	Instruction Cache Warp Scheduler Warp Scheduler Warp Scheduler																		
Di	Dispatch Dispatch			Dispatch Dispatch				Dispatch Dispatch				Dispatch Dispatch							
	Register File (65,536 x 32-bit)																		
Core		Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core		Core	DP Unit	Core	Core	Core	DP Ur	1it LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP Ur	iit LD/S1	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP Ur	iit LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP Ur	it LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP Ur	iit LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP Ur	iit LD/S1	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP Ur	iit LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP Ur	it LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP Ur	uit LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP Ur	lit LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LDIST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP Ur	iit LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP Ur	it LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP Ur	iit LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP Ur	iit LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit	LD/ST	SFU	Core	Core	Core	DP Unit	Core	Core	Core	DP Ur	nit LD/ST	SFU
Core	Core	Core	DP Unit	Core	Core	Core	DP Unit		SFU			Core	DP Unit	Core	Core	Core	DP Ur	iit LD/S1	SFU
	Interconnect Network																		
	64 KB Shared Memory / L1 Cache																		
	48 KB Read-Only Data Cache																		
<u> </u>	Tex		Tex			Tex		Tex		Tex		Tex		Tex			Tex		
	Tex		Tex			Tex		Tex	(Tex		Te>	(Tex		Te	ĸ

Key Architectural Ideas

SIMT (Single Instruction Multiple Thread) execution

- threads run in groups of 32 called warps
- warp threads execute same instructions
- HW automatically handles divergence
- Hardware multithreading
 - Allocate resources for many more threads than CUDA Cores
 - HW schedules which warp(s) to run next
- Any non-waiting warp can run
 - switching between warps is free



CUDA Core

Result Queu

INT Unit

FP Unit

Execution of Kernels on the GPU

Each block is executed on one SM

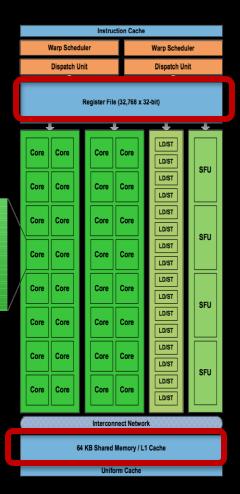
- cannot migrate
- reason for block independence

Block threads share SM resources

- SM registers are divided up among the threads
- SM shared memory can be read/written by all threads

Several blocks per SM possible

- if enough resources available
- SM resources are divided among all blocks

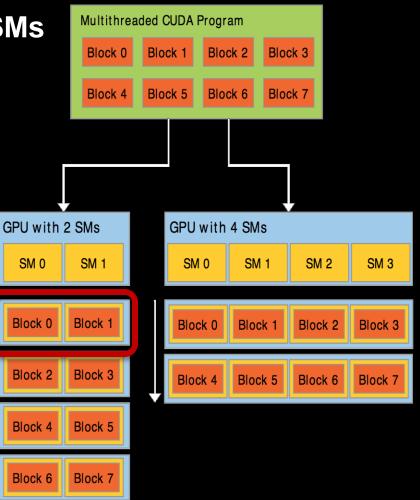


CUDA Core

INT Unit

Execution of Kernels on the GPU

- Blocks are distributed across SMs
- At each moment, one or more blocks are active
 - reside on a multiprocessor
 - resources allocated
 - executed until finished
- Others wait to be executed
 - not yet assigned to a SM



Execution on each Multiprocessor

On each SM, all blocks which reside on it are divided into warps (groups of 32 threads)

At each clock cycle:

- Each warp scheduler chooses a warp which is ready to be executed
- The next instruction of these warps are issued to the CUDA Cores
 - or to load/store units
 - or to special function units
 - or to texture units

Execution on each Multiprocessor

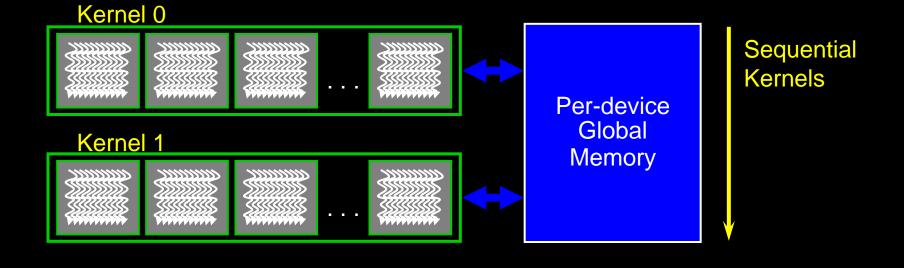
time

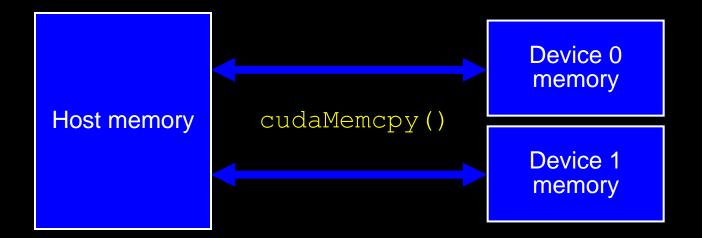
	Warp Scheduler	Warp Sc	Warp Scheduler					
	Dispatch Unit	Dispato	Dispatch Unit					
	Register File (32,768 x 32-bit)							
	Core Core Co	LD/ST						
	Core Core Co	ore Core LD/ST	SFU					
CUDA Core Dispatch Port Operand Collector	Core Core Co	LD/ST	SFU					
FP Unit INT Unit	Core Core Co	Core LD/ST						
Result Queue	Core Core Co	Core LD/ST	SFU					
	Core Core Co	Core LD/ST						
	Core Core Co	LD/ST	SFU					
	Core Core Co	Core LD/ST						
	Interconnect Network							
	64 KB Shared Memory / L1 Cache							
		Uniform Cache						

Warp Scheduler	Warp Scheduler			
Instruction Dispatch Unit	Instruction Dispatch Unit			
Warp 8 instruction 11	Warp 9 instruction 11			
Warp 2 instruction 42	Warp 3 instruction 33			
Warp 14 instruction 95	Warp 15 instruction 95			
Warp 8 instruction 12	Warp 9 instruction 12			
Warp 14 instruction 96	Warp 3 instruction 34			
Warp 2 instruction 43	Warp 15 instruction 96			

MEMORY MANAGEMENT

Memory Model





Memory Spaces

CPU and GPU have separate memory spaces

- Data is moved across PCIe bus
- Use functions to allocate/set/copy memory on GPU
 - Very similar to corresponding C functions

Pointers are just addresses

- Can't tell from the pointer value whether the address is on CPU or GPU
 - possible if CC >= 2.0 using unified addressing
- Must exercise care when dereferencing:
 - Dereferencing CPU pointer on GPU will likely crash
 - Same for vice versa

GPU Memory Allocation / Release

Host (CPU) manages device (GPU) memory:

- cudaMalloc (void ** pointer, size_t nbytes)
- cudaMemset (void * pointer, int value, size_t count)
- cudaFree (void* pointer)

```
int n = 1024;
int nbytes = 1024*sizeof(int);
int * d_a = 0;
cudaMalloc( (void**)&d_a, nbytes );
cudaMemset( d_a, 0, nbytes);
cudaFree(d_a);
```

Data Copies

cudaMemcpy(void *dst, void *src, size_t nbytes, enum cudaMemcpyKind direction);

- returns after the copy is complete
- blocks CPU thread until all bytes have been copied
- doesn't start copying until previous CUDA calls complete
- non-blocking copies are also available

enum cudaMemcpyKind

- cudaMemcpyHostToDevice
- cudaMemcpyDeviceToHost
- cudaMemcpyDeviceToDevice

Example: Host code for vecAdd

// allocate and initialize host (CPU) memory
float *h_A = ..., *h_B = ...; *h_C = ...(empty)

// allocate device (GPU) memory
float *d A, *d B, *d C;
cudaMalloc((void**) &d A, N * sizeof(float));
cudaMalloc((void**) &d B, N * sizeof(float));
cudaMalloc((void**) &d C, N * sizeof(float));

// copy host memory to device
cudaMemcpy(d_A, h_A, N * sizeof(float),
cudaMemcpyHostToDevice));
cudaMemcpy(d_B, h_B, N * sizeof(float),
cudaMemcpyHostToDevice));

// execute grid of N/256 blocks of 256 threads each
vecAdd<<<N/256, 256>>>(d_A, d_B, d_C);

Example: Host code for vecAdd (2)

// execute grid of N/256 blocks of 256 threads each
vecAdd<<<N/256, 256>>>(d_A, d_B, d_C);

// copy result back to host memory
cudaMemcpy(h_C, d_C, N * sizeof(float),
cudaMemcpyDeviceToHost));

// do something with the result ...

// free device (GPU) memory
cudaFree(d_A);
cudaFree(d_B);
cudaFree(d_C);

CUDA Short Summary

Thread Hierarchy

- thread smallest executable unity
- block group of threads, shared memory for collaboration
- grid consists of several blocks
- warp group of 32 threads

Keyword extensions for C/C++

 kernel - function called b 	ov CPU, executed on GPU
	y or of choodica on or o

device

global

__host__

<<<...>>>

- function called by GPU and executed on GPU
- [optional] function called and executed by CPU
- kernel launch, chevrons specify grid and block sizes

Compilation:

nvcc <filename>.cu -o <executable>