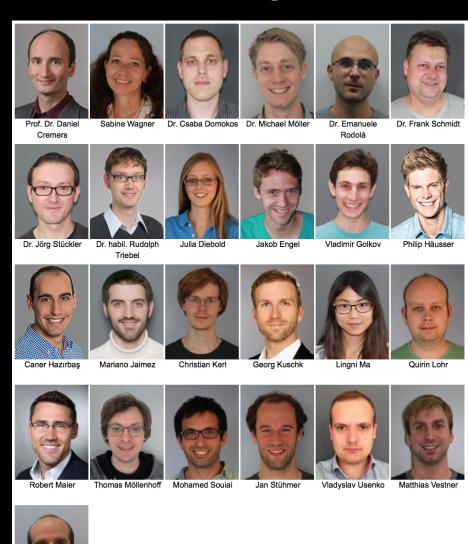
GPU Programming in Computer Vision

Thomas Möllenhoff, Robert Maier, Caner Hazırbaş

Introduction to Parallel Computing

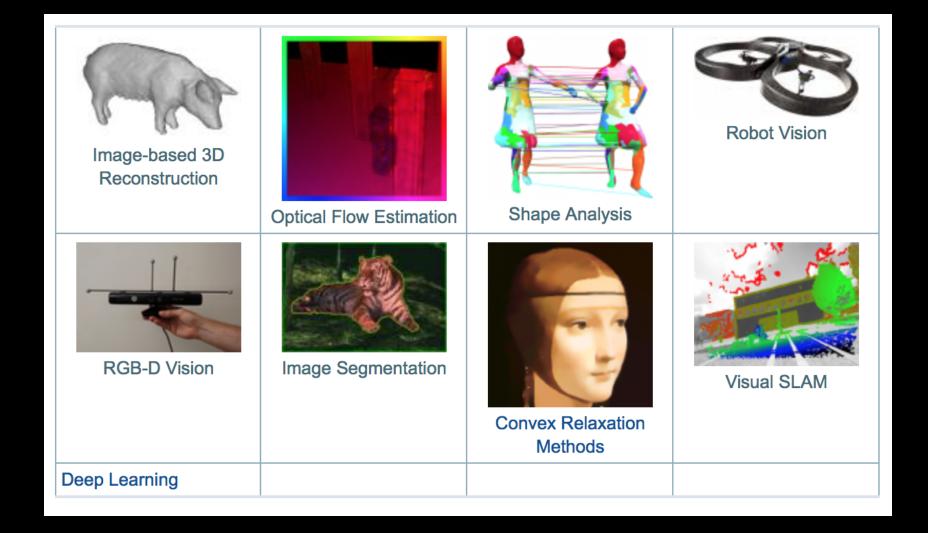
Technical University Munich, Computer Vision Group Summer Semester 2015, September 7 – October 6

Computer Vision Group



Thomas Windheuser

Our Research



This Course

- Parallel Programming with CUDA
- Computer Vision Basics
 - image filtering (convolution, diffusion, denoising)
 - basic algorithms
 - regularization (dealing with noise)
- Optimization + Numerics

Course Goals

- Learn how to program massively parallel processors and achieve
 - high performance
 - functionality and maintainability
 - scalability across future generations
- Acquire technical knowledge required to achieve above goals
 - principles and patterns of parallel programming
 - processor architecture features and constraints
 - programming API, tools and techniques
- Apply this knowledge to implement computer vision algorithms efficiently

Course Timeline: 07.09. – 06.10.

- September 7-11 (this week) : Lecture
 - 2-4h lectures (attendance mandatory)
 - programming exercises
 - groups of 2-3 students
- September 12-October 4: Student project
 - advanced applications
 - unsupervised
- October 5-6: Presentations

```
September 2015 >
Mo Di Mi Do Fr Sa So
31 1 2 3 4 5 6
7 8 9 10 11 12 13
14 15 16 17 18 19 20
21 22 23 24 25 26 27
28 29 30 1 2 3 4
5 6 7 8 9 10 11
```

Lecture Week

Lecture

- starts at 10:00 sharp each day
- attendance mandatory to pass the course

Exercises

- until 18:00 each day
- groups of 2-3 students
- present solutionson next day after the lecture

```
September 2015 >
Mo Di Mi Do Fr Sa So
31 1 2 3 4 5 6
7 8 9 10 11 12 13
14 15 16 17 18 19 20
21 22 23 24 25 26 27
28 29 30 1 2 3 4
5 6 7 8 9 10 11
```

"Work @ Home" ™

You can access your computer remotely:

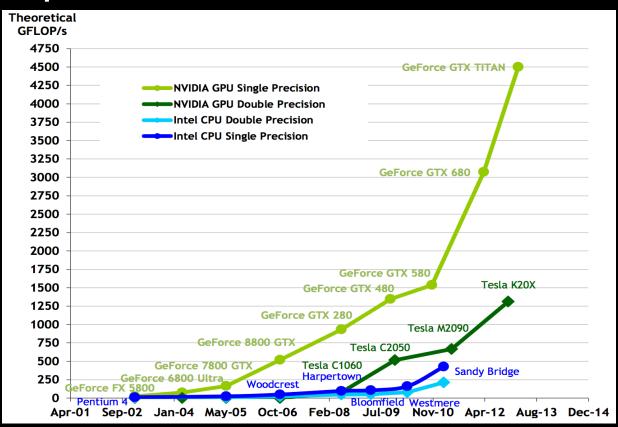
```
ssh -X p123@atradig789.informatik.tu-muenchen.de
```

- p123: replace with your login
- atradig789: replace with your computer name
 - type hostname to find out the name

- Works from within Linux or Mac
 - for Mac: install XQuartz first (X11 server)

Why Massively Parallel Processing?

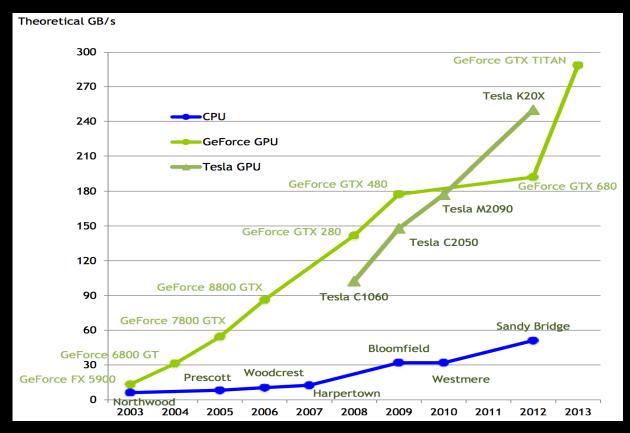
- A quiet revolution: Performance!
 - computations: TFLOPs vs. 100 GFLOPs



GPU in every PC – massive volume & impact

Why Massively Parallel Processing?

- A quiet revolution: Performance!
 - bandwidth: ~5x



GPU in every PC – massive volume & impact

Serial Performance Scaling is Over

- Cannot continue to scale processor frequencies
 - no 10 GHz chips

- Cannot continue to increase power consumption
 - can't melt chip

- Can continue to increase transistor density
 - as per Moore's Law

How to Use Transistors?

- Larger caches ... decreasing
- Instruction-level parallelism ... decreasing
 - out-of-order execution, speculation, ...
- Data-level parallelism ... increasing
 - vector units, SIMD execution, ...
 - Intel SSE, GPUs, ...
- Thread-level parallelism ... increasing
 - multithreading, multicore, manycore

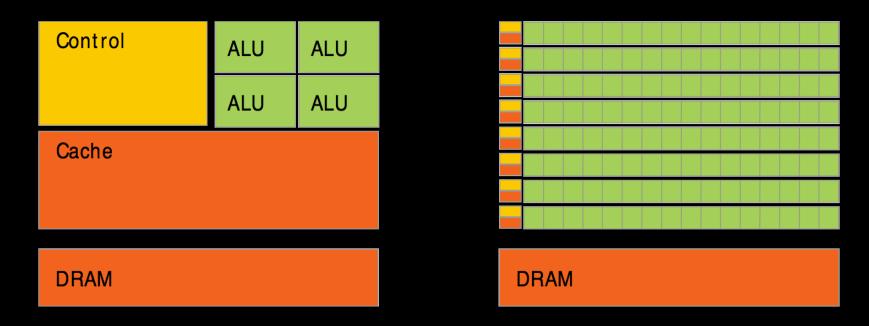
Design Difference: CPU vs. GPU

- Different goals produce different designs
 - CPU must be good at everything, parallel or not
 - GPU assumes work load is highly parallel
- CPU: minimize latency experienced by 1 thread
 - big on-chip caches
 - sophisticated control logic

- GPU: maximize throughput of all threads
 - skip big caches, multithreading hides latency
 - share control logic across many threads, SIMD
 - create and run thousands of threads

Design Difference: CPU vs. GPU

- Different goals produce different designs
 - CPU must be good at everything, parallel or not
 - GPU assumes work load is highly parallel



CPU minimize latency

GPU maximize throughput

Enter the GPU

Massively parallel



NVIDIA GPUs

- Compute Capability
 - version number of the hardware architecture
 - core architecture and incremental improvements

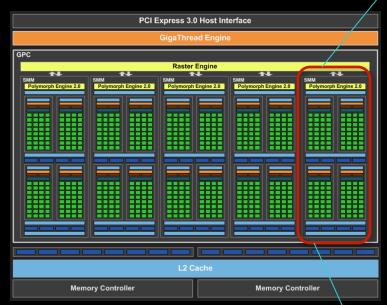
Arch	CC	GPUs	Features (e.g.)
Tesla (2007)	1.0	8800 GTX, Tesla C870	Basic functionality
	1.1	9800 GTX, Quadro FX 580	Atomics in global mem
	1.2	GT 240, Quadro FX 1800M	Atomics in shared mem
	1.3	GTX 285, Tesla C1060	Double precision
Fermi (2010)	2.0	GTX 480/580, Tesla C2070	Memory cache
	2.1	GTX 460, GTX 560 Ti	More cores (hardware)
Kepler (2012)	3.0	GTX 680/770, Tesla K10	Power efficiency, Many cores
	3.5	GTX 780/Titan, Tesla K20	Dynamic Parallelism, Hyper-Q
Maxwell (2014)	5.0	GTX 750, GTX 750 Ti	135% performance/core 200% performance/watt

NVIDIA GPUs

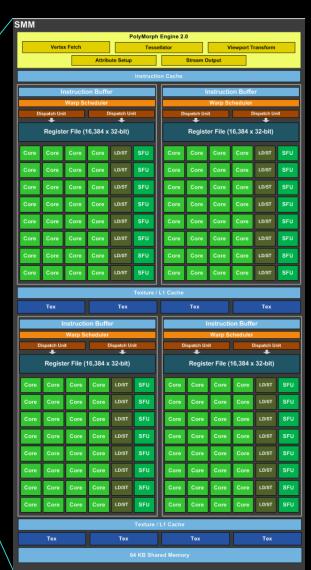
- Compute Capability
 - version number of the hardware architecture
 - core architecture and incremental improvements
- List of features for each Compute Capability:
 - see NVIDIA Programming Guide: Appendix G.1

NVIDIA GPUs: Current Architecture

1st gen Maxwell GPU



- 5 multiprocessors (up to)
- 128 Cuda Cores per SM
 - 640 Cores in total (up to)



Enter CUDA ("Compute Unified Device Architecture")

- Scalable parallel programming model
 - exposes the computational horsepower of GPUs
- Abstractions for parallel computing
 - let programmers focus on parallel algorithms
 - not mechanics of a parallel programming language
- Minimal extensions to familiar C/C++ environment to run code on the GPU
 - low learning curve

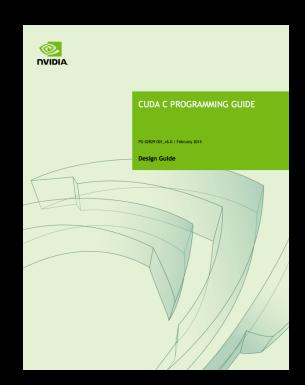
CUDA: Scalable Parallel Programming

- Provide straightforward mapping onto hardware
 - good fit to GPU architecture
 - maps well to multi-core CPUs too
- Execute code by many threads in parallel
- Scale to 100s of cores & 10,000s of threads
 - GPU threads are lightweight create / switch is free
 - GPU needs 1000s of threads for full utilization

Reference: CUDA Programming Guide

- CUDA comes with excellent documentation
 - doc/pdf in the CUDA folder, have a look!

- CUDA Programming Guide
 - one of the best CUDA references
 - covers every CUDA feature
 - provides in-depth explanations



- Also: list of all CUDA functions:
 - CUDA_Runtime_API.pdf

Outline of CUDA Basics

- Kernels and Thread Hierarchy
- Execution on the GPU
- Memory Management
- Error Handling And Compiling

See the Programming Guide for the full API

BASIC KERNELS AND THREAD HIERARCHY

CUDA Definitions

- Device: GPU
 - executes code in parallel
- Host: CPU
 - manages execution on the device
- Kernel: C/C++ function executed on the device
 - executed by many threads
 - each thread executes the same sequential program
 - each thread is free to execute a unique code path

Quick Example

CPU: Process subtasks serially one by one:

```
for( int i=0; i<n; i++ )
{
    c[i] = a[i] + b[i];
}</pre>
```

GPU: Process each subtask in its own thread:

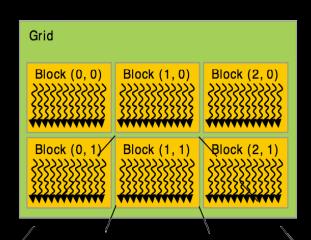
launch enough threads to cover all data

Thread Hierarchy

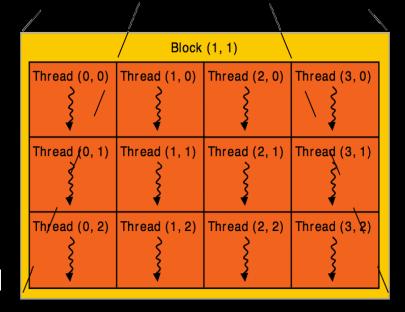
- Kernel threads are grouped into blocks
 - up to 512 or 1024 threads per block
- Idea: Threads from the same block can cooperate
 - synchronize their execution
 - communicate via shared memory
 - threads from different blocks cannot cooperate
- Allows transparent scaling to different GPUs
- All kernel blocks together form a grid

Thread Hierarchy

threads per block: up to 512 (CC 1.x), up to 1024 (CC>=2.0)

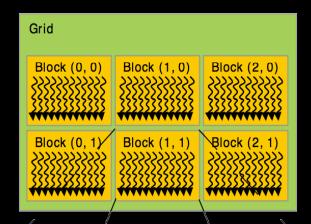


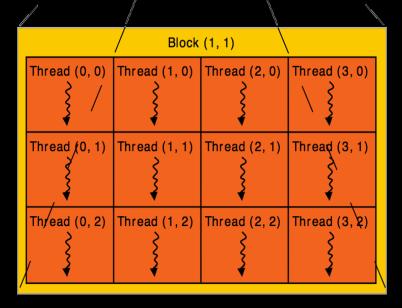
- Blocks can be 1D, 2D, or 3D
- Grids can be 1D, 2D, or 3D
 - CC 1.x: only 1D or 2D
- Dimensions set at launch
 - can be different for each grid



IDs and Dimensions

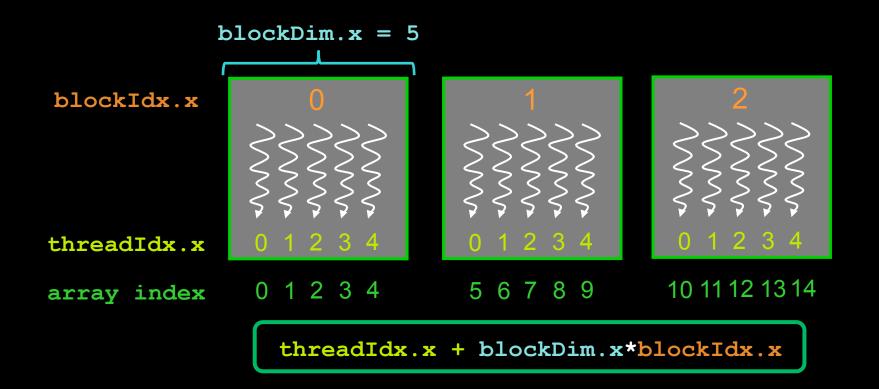
- Threads:
 - 3D IDs, unique within a block
- Blocks:
 - 3D IDs, unique within a grid
- Built-in variables:
 - threadIdx, blockIdx
 - blockDim, gridDim





Array Accesses: Index Calculation

- Obtain unique array index from block/thread IDs
 - threadIdx, blockIdx
 - blockDim, gridDim



Kernel Launch

Usual C/C++ function call, with an additional specification of grid and block sizes:

```
mykernel <<< grid, block >>> (...);
```

- dim3 grid; dim3 block;
 - three ints: block.x, block.y, block.z
- Kernel is launched by the CPU
 - CC >= 3.5: kernels can launch other kernels
- Executed on the GPU

Example: One-dimensional Kernel

```
global void mykernel (int *a, int n)
    int ind = threadIdx.x + blockDim.x * blockIdx.x;
    if (ind < n) a[ind] = a[ind] + 1;
int main()
    dim3 block = dim3(128,1,1); // 128 threads
    // ensure enough blocks to cover n elements (round up)
    \frac{\text{dim3 grid} = \text{dim3}( (n + \text{block.x} - 1) / \text{block.x}, 1, 1);}
    mykernel <<<grid, block>>> (d a, n);
    // Also possible:
    // launch 4 blocks, each with 128 threads
    mykernel <<<4,128>>> (d a, n);
```

Example: Two-dimensional Kernel

```
global void mykernel (int *a, int w, int h)
    int x = threadIdx.x + blockDim.x * blockIdx.x;
    int y = threadIdx.y + blockDim.y * blockIdx.y;
    int ind = x + w*y;
    if (x \le w \le v \le h) a[ind] = a[ind] + 1;
int main()
    \frac{\text{dim3 block}}{\text{dim3 (32,8,1)}}; \frac{1}{32*8} = 256 \text{ threads}
    // ensure enough blocks to cover w * h elements (round up)
    dim3 grid = dim3((w + block.x - 1) / block.x,
                        (h + block.y - 1) / block.y, 1);
    mykernel <<<grid,block>>> (d A, dimx, dimy);
```

Always Check Validity of Indices

- There may be more threads than array elements
 - Always test whether the indices are within bounds

```
global void mykernel (int *a, int n)
  int ind = threadIdx.x + blockDim.x * blockIdx.x;
  if (ind < n) a[ind] = a[ind] + 1;
global void mykernel (int *a, int w, int h)
  int x = threadIdx.x + blockDim.x * blockIdx.x;
  int y = threadIdx.y + blockDim.y * blockIdx.y;
  int ind = x + w*y;
  if (x \le x \le y \le n) a[ind] = a[ind] + 1;
```

Exercise: IDs of Threads and Blocks

```
kernel<<<4,4>>> (d a);
```

```
global void kernel (int *a)
   int idx = threadIdx.x + blockDim.x * blockIdx.x;
   a[idx] = 7;
                                     Output: 7777 7777 7777 7777
global void kernel (int *a)
   int idx = threadIdx.x + blockDim.x * blockIdx.x;
   a[idx] = blockIdx.x;
                                     Output: 0 0 0 0 1 1 1 1 1 2 2 2 2 3 3 3 3
global void kernel(int *a)
   int idx = threadIdx.x + blockDim.x * blockIdx.x;
   a[idx] = threadIdx.x;
                                     Output: 0 1 2 3 0 1 2 3 0 1 2 3 0 1 2 3
```

Code Executed on GPU: Functions

Special qualifiers to declare GPU functions:

```
__global__ : kernels
launched by CPU to run on the GPU
must return void
```

- __device__ : auxiliary GPU functions can only be called on the GPU called from __global__ or __device__ functions
- host : "normal" CPU C/C++ functions can only be called on the CPU
- __host___ device__ : qualifiers can be combined callable from CPU and from GPU

Code Executed on GPU: Restrictions

- C/C++ with some restrictions
 - only access to GPU memory
 - not to CPU memory
 - can access "pinned" CPU memory (special allocation needed)
 - from CUDA 6 and CC 3.0: GPU can access CPU memory
 - no access to host functions
 - no variable number of arguments
 - no static variables in functions or classes

Code Executed on GPU: Features

- Many C/C++ features available for GPU code
 - templates
 - recursion (CC >=2.0)
 - overloading
 - function overloading
 - operator overloading
 - classes
 - stack allocation
 - heap allocation (CC >= 2.0)
 - inheritance, virtual functions (CC >= 2.0)
 - function pointers (CC >= 2.0)
 - printf() formatted output (CC >= 2.0)
- Vector variants of basic types
 - float2, float3, float4, double2, int4, char2, etc.
 - float2 a=make float2(1,2); a.x=10; a.y=a.x;

Blocks: Must Be Independent

- Any possible ordering of blocks should be valid
 - presumed to run to completion without pre-emption
 - can run in any order (order is unspecified)
 - can run concurrently OR sequentially
- Blocks may coordinate but not synchronize
 - shared queue pointer: OK
 - shared lock: BAD ... can easily deadlock
- Independence requirement gives scalability

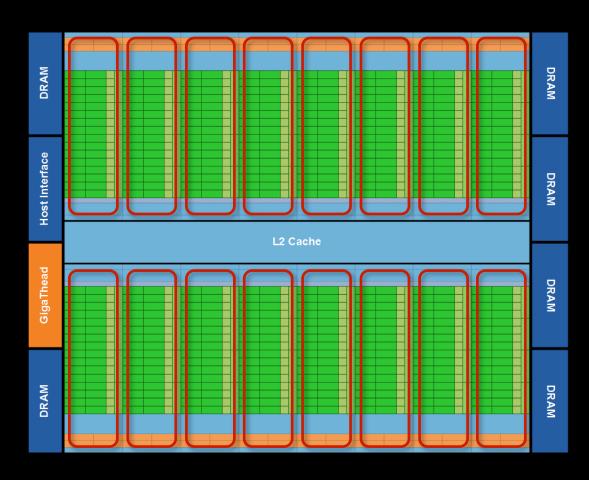
Execution of Kernels: Asynchronous

- Kernel launches are asynchronous w.r.t. CPU
 - after kernel launch, control immediately returns
 - CPU is free to do other work while the GPU is busy
- Kernel launches are queued
 - kernel doesn't start until previous kernels are finished
 - concurrent kernels possible for CC >= 2.0 (given enough resources)
- Explicit synchronization if needed
 - cudaDeviceSynchronize()

EXECUTION ON GPU

NVIDIA GPU Architecture

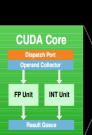
Fermi GPU (CC 2.x)

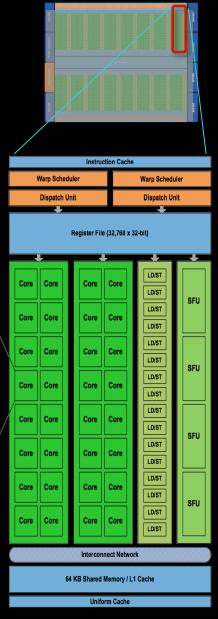


- 16 independent multiprocessors (SMs)
- No shared resources except global memory
- No synchronization, always work in parallel

Single Fermi SM Multiprocessor

- 32 CUDA Cores per SM (512 total)
 - arithmetic/logic operations
- 16 memory load/store units
 - (slow) access to off-chip GPU mem
- 4 Special Function Units
 - 1/X, 1/SQRT(X), SIN, COS, EXP, ...
- 64 KB on-chip shared memory
 - shared amongst CUDA cores
 - enables thread communication





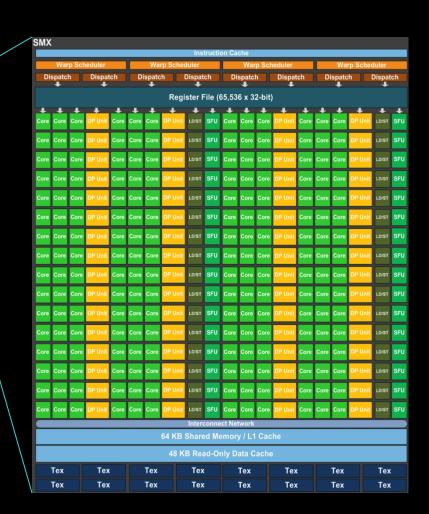
NVIDIA GPU Architecture: Current

Kepler GPU (CC 3.x)



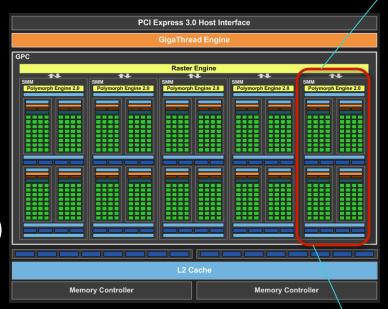
15 multiprocessors (up to)

- 192 Cuda Cores per SM
 - 2880 Cores in total (up to)

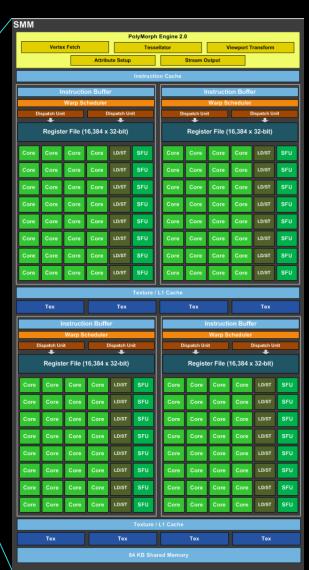


NVIDIA GPU Architecture: Current

Maxwell GPU (CC 5.0)



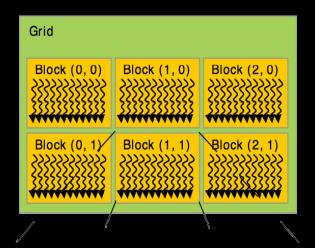
- 5 multiprocessors (up to)
- 128 Cuda Cores per SM
 - 640 Cores in total (up to)

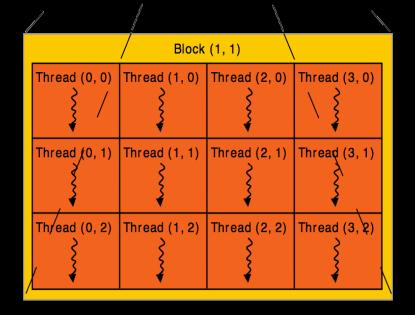


Warps: Key Architectural Idea

- SIMT (Single Instruction Multiple Thread) execution
 - threads run in groups of 32 called warps
- All 32 threads in a warp execute the same instruction
 - always, no matter what (even if threads diverge)
- Threads are executed warp-wise by the GPU
 - for each warp, the 32 threads are executed in parallel
 - warps are executed one after another
 - but several warps can run simultaneously
 - up to 2 for CC 2.x, up to 6 for CC 3.x

Thread Hierarchy

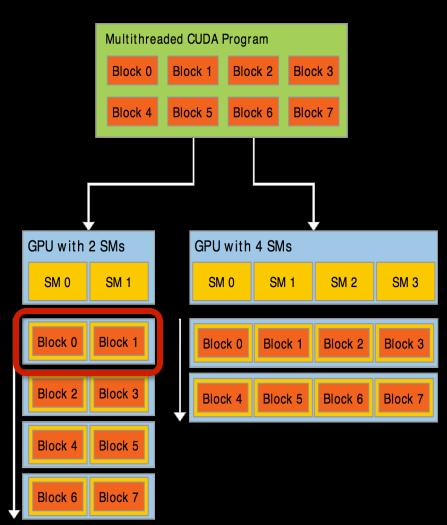




Execution of Kernels on the GPU

Blocks are distributed across the Multiprocessors (SMs)

- Active blocks
 - are currently executed
 - reside on a multiprocessor
 - resources allocated
 - executed until finished
- Waiting blocks
 - wait to be executed
 - not yet assigned to a SM



Blocks Execute on Multiprocessors

Each block is executed on one Multiprocessor (SM)

- cannot migrate
- reason for block independence
- Several blocks per SM possible
 - if enough resources available
 - SM resources are divided among all blocks
- Block threads share SM resources
 - SM registers are divided up among the threads
 - SM shared memory can be read/written by all threads

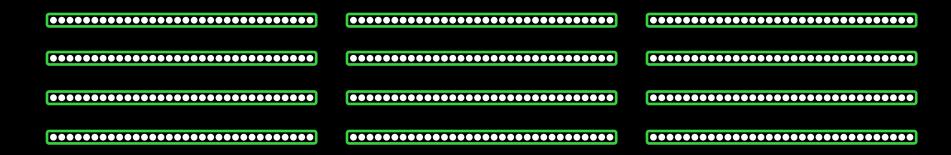


• Assume there are three blocks on one SM, with 128 threads per block:

block 0 block 1 block 2

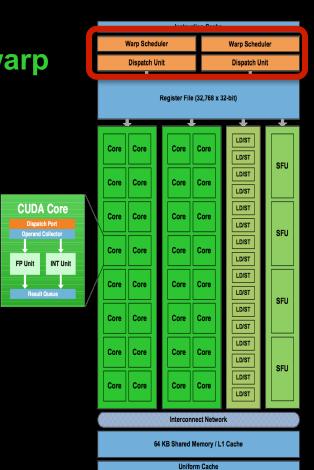
128 128 128 128 threads threads

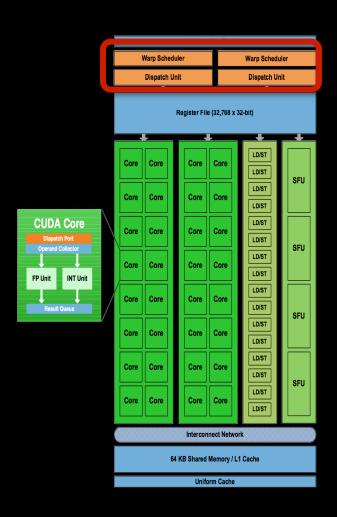
- Threads from all blocks are divided into warps
- In our example:
 - 4 warps from every block (128 threads/32)
 - 12 warps overall on SM (3 blocks * 4 warps/block)
 - 12*32 = 384 threads

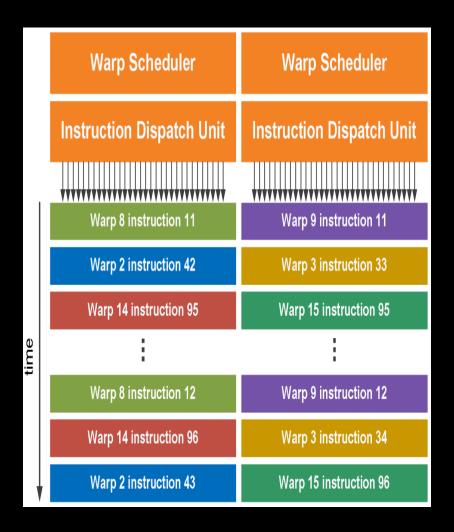


- Resources are allocated for all potential warps
 - the state of every potentially executable warp is always present on the Multiprocessor, until finished
 - overall many more potentially executable threads than CUDA Cores possible
- Therefore:
 - switching between warps is free
 - any non-waiting warp can run

- At each clock cycle
 - each warp scheduler chooses a warp which is ready to be executed
- For each chosen warp
 - the next instruction is executed for all 32 threads of the warp
 - issued for execution to
 - CUDA Cores
 - or load/store units
 - or special function units
 - or texture units

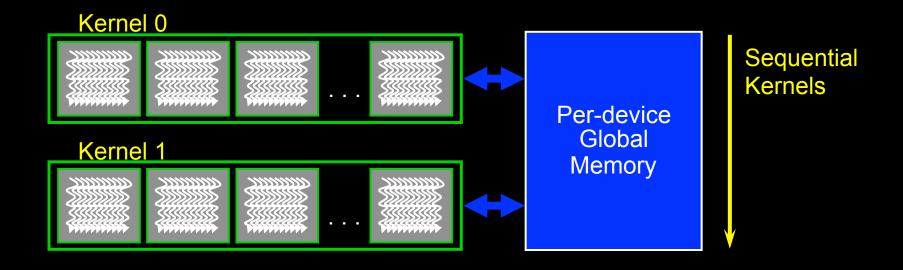


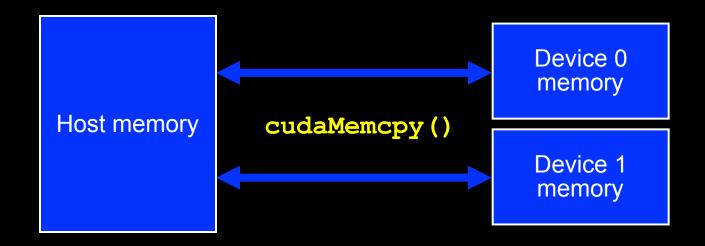




MEMORY MANAGEMENT

GPU Memory





GPU Memory

- CPU and GPU have separate memory spaces
 - data is moved across PCle bus
 - use functions to allocate/set/copy memory on GPU
 - very similar to corresponding C functions
- Pointers are just addresses
 - cannot tell from pointer if memory is on GPU or CPU
 - but possible for CC>=2.0: unified virtual addressing
 - must exercise care when dereferencing:
 - crash if GPU dereferences pointer to CPU memory
 - and vice versa

Allocation / Release

- Host (CPU) manages device (GPU) memory:
 - cudaMalloc (void **pointer, size_t nbytes)
 - cudaMemset (void *pointer, int value, size_t count)
 - cudaFree (void* pointer)

```
int n = 1024;
size_t nbytes = n*sizeof(int);
int *d_a = NULL;
cudaMalloc(&d_a, nbytes);
cudaMemset(d_a, 0, nbytes);
cudaFree(d_a);
```

Data Copies Between GPU and CPU

- cudaMemcpy (void *dst, void *src, size_t nbytes, cudaMemcpyKind direction);
 - blocks the CPU thread until all bytes have been copied
 - non-blocking variants are also available
 - doesn't start copying until all previous CUDA calls complete
- cudaMemcpyKind:
 - cudaMemcpyHostToDevice
 - cudaMemcpyDeviceToHost
 - cudaMemcpyDeviceToDevice

```
cudaMemcpy(dev_ptr, host_ptr, n*sizeof(float),
    cudaMemcpyHostToDevice);
```

Example Host Code

```
// allocate and initialize host (CPU) memory
float *h a = ..., *h b = ...; *h c = ... (empty)
// allocate device (GPU) memory
float *d a, *d b, *d c;
cudaMalloc( &d a, n * sizeof(float) );
cudaMalloc( &d b, n * sizeof(float) );
cudaMalloc( &d c, n * sizeof(float) );
// copy host memory to device
cudaMemcpy( d a, h a, n * sizeof(float), cudaMemcpyHostToDevice );
cudaMemcpy( d b, h b, n * sizeof(float), cudaMemcpyHostToDevice );
// launch kernel
dim3 block = dim3(128,1,1);
\dim 3 \text{ grid} = \dim 3((n + block.x - 1) / block.x, 1, 1);
vecAdd \ll grid_block >>> (d a, d b, d c);
// copy result back to host (CPU) memory
cudaMemcpy( h c, d c, n * sizeof(float), cudaMemcpyDeviceToHost );
// do something with the result...
// free device (GPU) memory
cudaFree(d a);
cudaFree(d b);
cudaFree(d c);
```

Use float by Default

- GPUs can handle double since CC>=1.3
- But float operations are still much faster
 - by an order of magnitude
 - so use double only if float is not enough
- Avoid using double where not needed:
 - Add 'f' suffix to float literals:
 - 0.f, 1.0f, 3.1415f are of type float
 - 0.0, 1.0, 3.1415 are of type double
 - Use float version of math functions:
 - expf/logf/sinf/sqrtf/etc. take and return float
 - exp / log / sin / sqrt / etc. take and return double

Blocks Size: How To Choose?

- Number of threads/block should be multiple of 32
 - because threads are always executed in groups of 32
- Rules of thumb:
 - not too small or too big: between 128 and 256 threads
 - start with dim3 (32,8,1), i.e. 256 threads
 - experiment with similar sized "power-of-2"-blocks:
 - **(64,4,1), (128,2,1), (32,4,1), (64,2,1)**
 - **(32,16,1), (64,8,1), (128,4,1), (256,2,1)**
 - measure the run time and choose the best block size!

ERROR HANDLING AND COMPILING

Error Handling

- Checking for errors is crucial for programming GPUs
- cudaError t cudaGetLastError()
 - returns the code for the last error
 - resets the error flag back to cudaSuccess
 - cudaPeetAtLastError(): get error code without resetting it
 - if everything OK: cudaSuccess
- char* cudaGetErrorString(cudaError t code)
 - returns a C-string describing the error

```
cudaMalloc(&d_a, n*sizeof(float));
cudaError_t e = cudaGetLastError();
if (e!=cudaSuccess)
{
    cerr << "ERROR: " << cudaGetErrorString(e) << endl;
    exit(1);
}</pre>
```

Error Handling

- Kernel execution is asynchronous
 - first wait for the kernel to finish by cudaDeviceSynchronize()
 - only then call cudaGetLastError()
 - otherwise it will be called too soon, the error may not have yet occured
 - kernel launch itself may produce errors due to invalid configurations
 - too many threads/block, too many blocks, too much shared memory requested
- Kernels may produce subtle memory corruption errors
 - may get unnoticed even after cudaDeviceSynchronize()
 - subsequent CUDA calls may or may not fail because of such an error
 - if they do fail, they were not the origin of the error
- It helps to keep track of the previous x CUDA calls
 - x=1, or x=2, or x=10

Compiling

CUDA files have ending .cu: squareArray.cu

- NVidia CUDA Compiler: nvcc
 - handles the CUDA part
 - hands over pure C/C++ part to host compiler

```
nvcc -o squareArray squareArray.cu
```

Additional info about the kernels by option

```
--ptxas-options=-v
```

CUDA Short Summary

Thread Hierarchy

thread - smallest executable unity

block - group of threads, shared memory for collaboration

grid - consists of several blocks

warp - group of 32 threads

Keyword extensions for C/C++

__global__ - kernel - function called by CPU, executed on GPU

__device___ - function called by GPU and executed on GPU

__host__ - [optional] - function called and executed by CPU

<><...>>> - kernel launch, chevrons specify grid and block sizes

Compilation:

nvcc -o <executable> <filename>.cu --ptxas-options=-v