

Chapter 5

Operator Splitting Methods

Convex Optimization for Computer Vision
SS 2016

Relations

Monotone Operators

Fixed Point Iterations

Proximal Point
Algorithm

PDHG Revisited

Douglas-Rachford
Splitting

Michael Moeller
Thomas Möllenhoff
Emanuel Laude
Computer Vision Group
Department of Computer Science
TU München

Recap and Motivation

- **Last 3 lectures:** PDHG method for minimizing structured convex problems

$$\min_{u \in \mathbb{R}^n} G(u) + F(Ku)$$



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- Key ideas: monotone operators, fixed point iterations

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- Unintuitive overrelaxation, rather involved convergence analysis
- Next lectures: simple and unified convergence analysis of many different algorithms within a single approach
- Key ideas: monotone operators, fixed point iterations
- Give a new understanding of convex optimization algorithms



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- Concept: identifying functions with their *graph*



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- Another possible view: think of relations as a set valued functions, e.g., $\partial E : \mathbb{R}^n \rightarrow \mathcal{P}(\mathbb{R}^n)$



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Solve generalized equation (inclusion) problem

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- Finding saddle-points (\tilde{u}, \tilde{p}) of

$$PD(u, p) = G(u) - F^*(p) + \langle Ku, p \rangle$$

corresponds to the inclusion problem

$$0 \in \begin{bmatrix} \partial G & K^T \\ -K & \partial F^* \end{bmatrix} \begin{bmatrix} u \\ p \end{bmatrix}$$



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→ **Draw a picture for $E(u) = |u|$**





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The set-valued operator $T \subset \mathbb{R}^n \times \mathbb{R}^n$ is called **monotone** if

$$\langle u - v, Tu - Tv \rangle \geq 0, \quad \forall u, v \in \mathbb{R}^n. \quad \text{Notation}^1$$

An operator T is called **maximally monotone** if it is not contained in any other monotone operator.



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- Proximity operators of convex functions $\text{prox}_{\tau f} : \mathbb{R}^n \rightarrow \mathbb{R}^n$

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Calculus rules (exercise):

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- Lipschitz operators (and in particular nonexpansive operators) are single-valued (functions)



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- If F is nonexpansive (Lipschitz constant $L \leq 1$) and $\text{dom } T = \mathbb{R}^n$ then the set of fixed points $(I - F)^{-1}(0)$ is closed and convex (**exercise**)

Resolvent and Cayley Operators

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Facts:

- $0 \in Tx$ if and only if $x = J_{\lambda T}x = C_{\lambda T}x$
- If T is monotone, then $J_{\lambda T}$ and $C_{\lambda T}$ are nonexpansive



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- Allows simple and unified analysis

Iteration of Contraction Mappings

Contraction Mapping Theorem

Suppose that $F : \mathbb{R}^n \rightarrow \mathbb{R}^n$ is a contraction with Lipschitz constant $L < 1$. Then the fixed point iteration

$$u^{k+1} = Fu^k,$$

also called contraction mapping algorithm, converges to the unique fixed point of F .

→ Proof: see literature²



²This theorem is also known as the Banach fixed point theorem.

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- $I - \tau \nabla E$ is contractive for $\tau \in (0, 2/L)$



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- The mapping $F : \mathbb{R}^n \rightarrow \mathbb{R}^n$ is called *averaged* if $F = (1 - \theta)I + \theta T$, for some nonexpansive operator T and $\theta \in (0, 1)$



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Theorem: Krasnosel'skii-Mann

Let $F : \mathbb{R}^n \rightarrow \mathbb{R}^n$ be averaged, and denote the (non-empty) set of fixed points of F as U . Then the sequence (u^k) produced by the iteration

$$u^{k+1} = Fu^k$$

converges to a fixed point $u^* \in U$, i.e., $u^k \rightarrow u^*$.

→ Proof: board!



Example: gradient method

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with $\theta = \tau L/2 < 1$.



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- Hence, we get convergence of the gradient descent method from the previous theorem



Proximal Point Algorithm

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The Proximal Point Algorithm

- Recall our original goal of finding $u \in \mathbb{R}^n$ with

$$0 \in Tu,$$

for $T \subset \mathbb{R}^n \times \mathbb{R}^n$ monotone.



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Definition: Proximal Point Algorithm (PPA)³

Given some maximally monotone operator $T \subset \mathbb{R}^n \times \mathbb{R}^n$, and some sequence $(\lambda_k) > 0$. Then the iteration

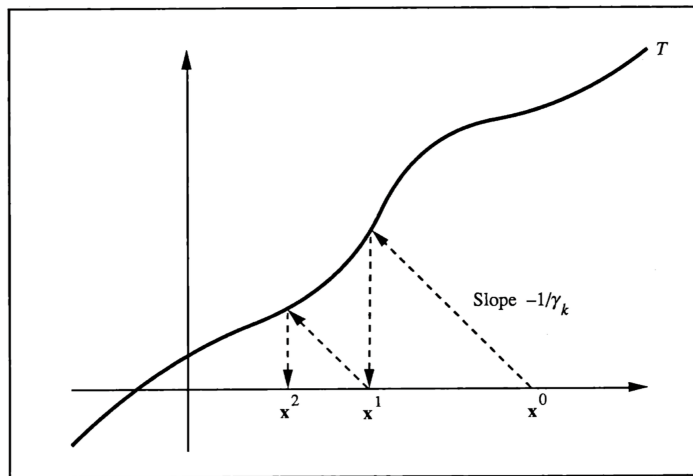
$$u^{k+1} = (I + \lambda_k T)^{-1} u^k,$$

is called the *proximal point algorithm*.

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Intuition of the Proximal Point Algorithm ⁴



⁴Eckstein, Splitting methods for monotone operators with applications to parallel optimization, 1989, pp. 42



Convergence of Proximal Point Algorithm

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- The resolvent $J_{\lambda T} = (I + \lambda T)^{-1}$ is an averaged operator



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$$C_{\lambda T} := 2J_{\lambda T} - I \Leftrightarrow J_{\lambda T} = \frac{1}{2}I + \frac{1}{2}C_{\lambda T}$$



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PDHG Revisited

PDHG as Proximal Point Method

- Remember that for convex-concave saddle point problems

$$PD(u, p) = G(u) - F^*(p) + \langle Ku, p \rangle$$

we have the following:

$$(\tilde{u}, \tilde{p}) = \arg \min \max_{u, p} PD(u, p) \Leftrightarrow \begin{bmatrix} 0 \\ 0 \end{bmatrix} \in \underbrace{\begin{bmatrix} \partial G(\tilde{u}) + K^T \tilde{p} \\ -K\tilde{u} + \partial F^*(\tilde{p}) \end{bmatrix}}_{=: T(\tilde{u}, \tilde{p})}$$



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- Stack primal and dual variables into vector $z = (u, p)^T$:

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- Plugging things in yields

$$u^k - u^{k+1} \in \lambda \partial G(u^{k+1}) + \lambda K^T p^{k+1}$$

$$p^k - p^{k+1} \in \lambda \partial F^*(p^{k+1}) - \lambda K u^{k+1}$$



PDHG as Proximal Point Method

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leads to:

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- Almost looks like the PDHG method, step size λ
- Problem:** cannot implement this algorithm, since updates in u^{k+1} and p^{k+1} depend on each other



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- This is called generalized or customized PPA:

$$0 \in M(z^{k+1} - z^k) + Tz^{k+1} \Leftrightarrow z^{k+1} = (M + T)^{-1} Mz^k$$



Convergence of Customized Proximal Point Method

- For symmetric, positive definite M , we can write $M = L^T L$, L invertible (Cholesky decomposition)

Operator Splitting
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⁵Bauschke, Combettes, Convex Analysis and Monotone Operator Theory in Hilbert Spaces, Theorem 24.5

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- Apply classical PPA to operator $T' = L^{-T} \circ T \circ L^{-1}$

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- Hence customized PPA inherits convergence from classical proximal point

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Convergence of PDHG

- When is the step size matrix symmetric positive definite?

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Lemma (Pock-Chambolle-2011⁶)

Let $\theta = 1$, T and Σ symmetric positive definite maps satisfying

$$\left\| \Sigma^{\frac{1}{2}} K T^{\frac{1}{2}} \right\|^2 < 1,$$

then the block matrix

$$M = \begin{bmatrix} T^{-1} & -K^T \\ -\theta K & \Sigma^{-1} \end{bmatrix}$$

is symmetric and positive definite.

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Summary

- Customized proximal point algorithms yield a whole family of methods, many choices of M are conceivable

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- **Next lecture:** Douglas-Rachford splitting and ADMM

Organizational Remarks

Exams:

- **Important:** Registration deadline 30.06. in TUMonline!

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Douglas-Rachford Splitting

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Motivation

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$$0 \in Tu \Leftrightarrow u = (I + \lambda T)^{-1}u$$



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- Resolvents $J_{\lambda A} = (I + \lambda A)^{-1}$ and $J_{\lambda B} = (I + \lambda B)^{-1}$ can be more easily evaluated than $J_{\lambda T}$

Splitting methods

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$$0 \in Au + Bu \Leftrightarrow C_A C_B v = v, u = J_B v$$

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Douglas-Rachford
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\rightarrow **Draw a picture for $T = \partial \iota_{C_1} + \partial \iota_{C_2}$!**

- *Peaceman-Rachford* splitting is undamped iteration

$$v^{k+1} = C_A C_B v^k$$



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- Always converges if there exists a solution $0 \in Au^* + Bu^*$, since it's fixed point iteration of averaged operator

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Douglas-Rachford Splitting (DRS)



- The Douglas-Rachford iteration $v^{k+1} = \left(\frac{1}{2}I + \frac{1}{2}C_A C_B\right) v^k$ can be written as

$$u_b^{k+1} = J_B(v^k),$$

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- u_a^k and u_b^k can be thought of estimates to a solution
- v^k running sum of residuals, drives u_a^k and u_b^k together

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Application to Convex Optimization

- Let's apply DRS to minimize

$$\min_{u \in \mathbb{R}^n} G(u) + F(u)$$



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$$0 \in \tau \partial G(u) + \tau \partial F(u)$$



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- Find zero of $T = A + B$, $A = \tau \partial G$, $B = \tau \partial F$
- The algorithm becomes (after slight simplifications):

$$\begin{aligned} u^{k+1} &= \text{prox}_{\tau G}(v^k), \\ v^{k+1} &= \text{prox}_{\tau F}(2u^{k+1} - v^k) + v^k - u^{k+1}. \end{aligned}$$

- We can rewrite the step in v^{k+1} using Moreau's Identity

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$$\begin{aligned} v^{k+1} &= \text{prox}_{\tau F}(2u^{k+1} - v^k) + v^k - u^{k+1} \\ &= u^{k+1} + \tau \text{prox}_{(1/\tau)F^*}((2u^{k+1} - v^k)/\tau) \end{aligned}$$



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- Looks familiar? :-)
- Applying DRS on the primal problem $\min_u G(u) + F(u)$ is equivalent to PDHG!

Optimization Problems with Compositions

- Ideally we'd like to solve problems of the form

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- Even problems looking very complicated at first glance can be split up into many simple substeps

Option 1: Graph Projection Splitting

- We want to minimize for $K : \mathbb{R}^n \rightarrow \mathbb{R}^m$

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- Proximal operator for \tilde{F} is projection onto the graph of $Ku = w$ (solving a least squares problem)

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- Other possibility: factorization caching

⁸N. Parikh, S. Boyd, Block Splitting for Distributed Optimization, 2014

Option 2: DRS for Problems with Compositions

- Consider the dual problem to $\min_u G(u) + F(Ku)$

$$\min_p G^*(-K^*p) + F^*(p) = (G^* \circ -K^*)(p) + F^*(p)$$



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- Reorder slightly with new variable w^{k+1}

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Option 2: DRS for Problems with Compositions

- Even more simple:

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$$0 \in \partial G(u^{k+1}) + \sigma K^T(Ku^{k+1} + \frac{1}{\sigma}(p^k - \sigma Ku^k))$$

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- Adding and subtracting $K^T p^{k+1}$ to first line yields

$$0 \in \partial G(u^{k+1}) + K^T p^{k+1} + \sigma K^T K(u^{k+1} - u^k) - K^T(p^{k+1} - p^k)$$

$$0 \in \partial F^*(p^{k+1}) - Ku^{k+1} - K(u^{k+1} - u^k) + \frac{1}{\sigma}(p^{k+1} - p^k)$$



Relation to PDHG

- Previous iterations can be written as PPA, $z = (u, p)^T$:

$$0 \in \underbrace{\begin{bmatrix} \partial G & K^T \\ -K & \partial F^* \end{bmatrix} \begin{bmatrix} u^{k+1} \\ p^{k+1} \end{bmatrix}}_{Tz^{k+1}} + \underbrace{\begin{bmatrix} \sigma K^T K & -K^T \\ -K & \frac{1}{\sigma} I \end{bmatrix}}_M \underbrace{\begin{bmatrix} u^{k+1} - u^k \\ p^{k+1} - p^k \end{bmatrix}}_{z^{k+1} - z^k}$$



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Proximal Point
Algorithm

PDHG Revisited

Douglas-Rachford
Splitting

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Alternating Direction Method of Multipliers (ADMM)

- Recall this formulation

$$u^{k+1} = \operatorname{argmin}_u G(u) + \frac{\sigma}{2} \left\| Ku + \frac{v^k}{\sigma} \right\|^2,$$

$$p^{k+1} = \operatorname{prox}_{\sigma F^*}(v^k + 2\sigma Ku^{k+1}),$$

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- Apply Moreau's identity to step in p^{k+1}

$$u^{k+1} = \operatorname{argmin}_u G(u) + \frac{\sigma}{2} \left\| Ku + \frac{v^k}{\sigma} \right\|^2,$$

$$p^{k+1} = v^k + 2\sigma Ku^{k+1} - \sigma \operatorname{prox}_{\sigma F}\left(\frac{v^k}{\sigma} + 2Ku^{k+1}\right),$$

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Alternating Direction Method of Multipliers (ADMM)

- Make new variable for $\text{prox}_{\sigma F}$ -step, write prox as argmin:

$$u^{k+1} = \underset{u}{\text{argmin}} G(u) + \frac{\sigma}{2} \left\| Ku + \frac{v^k}{\sigma} \right\|^2,$$

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- Replacing the variable v^k in the u^{k+1} update yields

$$u^{k+1} = \underset{u}{\text{argmin}} G(u) + \frac{\sigma}{2} \left\| Ku + \frac{p^k - \sigma Ku^k}{\sigma} \right\|^2,$$

Alternating Direction Method of Multipliers (ADMM)

- Replace variable p^k in all update steps

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- Rewrite as:

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Alternating Direction Method of Multipliers (ADMM)

- Using the following fact we can further rewrite the updates:

$$\operatorname{argmin}_a \frac{\sigma}{2} \left\| a - \frac{b}{\sigma} \right\|^2 = \operatorname{argmin}_a - \langle a, b \rangle + \frac{\sigma}{2} \|a\|^2$$



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- Reintroduce $p^{k+1} = v^k + \sigma Ku^{k+1}$, can be rewritten as:

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Alternating Direction Method of Multipliers (ADMM)

- Let $\bar{w}^{k+1} = w^k$:

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- Change order of first two iterates:

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Alternating Direction Method of Multipliers (ADMM)

- Final update equations:

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$$L_{\text{aug}}^{\tau}(u, w, p) = G(u) + F(w) + \langle p, Ku - w \rangle + \frac{\tau}{2} \|Ku - w\|^2$$

⁹Boyd et al., Distributed optimization and statistical learning via the alternating direction method of multipliers, 2011



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- It has gained enormous popularity recently ⁹, over 3458 citations in 5 years

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- Many relations exist between the primal-dual algorithms, often special cases of one another
- Depending on the problem structure, better to use either Graph Projection/DRS/ADMM or PDHG (more next week!)
- **Rule of thumb:** Graph Projection/DRS/ADMM few expensive iterations, PDHG many cheap iterations

