

Frank R. Schmidt Matthias Vestner

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### Geometric deep learning on graphs and manifolds

**Short Course** 



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### Michael Bronstein

- Associate Professor USI Lugano
- Associate Professor Tel Aviv University
- Principal Engineer at the Intel Perceptual Computing
  - authored the first book on non-rigid shape analysis
  - www.inf.usi.ch/bronstein/





10. Isometries, Rigid Alignment





**Isometries** 

## Geodesic distance



Let  ${\mathcal M}$  be a manifold. We define the geodesic distance between two points  $x, y \in \mathcal{M}$  as

$$d_{\mathcal{M}}(x,y) = \inf\{ \text{length}(c) | c : [0,1] \to \mathcal{M}, c(0) = x, c(1) = y \}.$$

- For the manifolds we consider (compact) there exists a minimizer (not nec. unique)
- Using the first fundamental form the length of curves can be measured in the parameter domain
- every submanifold comes with a natural metric induced by the first fundamental form
- we ommit the proof that d is actually a metric





## **Isometries**

 $d_{\mathcal{M}}(x,y) = d_{\mathcal{N}}(\Phi(x),\Phi(y))$  for all points  $x,y \in \mathcal{M}$ .

A mapping  $\Phi: \mathcal{M} \to \mathcal{N}$  between two shapes (manifolds) is an isometry if

$$a_{\mathcal{M}}(x,y) = a_{\mathcal{N}}(\Psi(x), \Psi(y))$$
 for all points  $x, y \in \mathcal{M}$ .





If such a mapping exists  ${\mathcal M}$  and  ${\mathcal N}$  are called isometric. Many shape matching approaches assume that the shapes to be matched are (nearly) isometric. The task then becomes to find the (almost-)isometry  $\Phi$ .

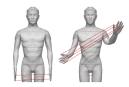


### Intrinsic symmetry



Most of the shapes we consider come with an intrinsic symmetry  $S:\mathcal{M}\to\mathcal{M}$ , such that

$$d_{\mathcal{M}}(x,y) = d_{\mathcal{M}}(S(x),S(y)), \quad \forall x,y \in \mathcal{M}, S \neq \mathrm{id}.$$



As a consequence isometries are not unique. Let  $\Phi:\mathcal{M}\to\mathcal{N}$  be an isometry and  $S:\mathcal{M}\to\mathcal{M}$  be an intrinsic symmetry. Then  $\Phi\circ S^{-1}$  is also an isometry:

$$d_{\mathcal{M}}(x,y)=d_{\mathcal{M}}(S^{-1}(x),S^{-1}(y))=d_{\mathcal{N}}(\Phi\circ S^{-1}(x),\Phi\circ S^{-1}(y))$$

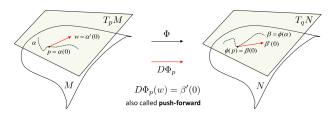


### **Push Forward**





We can define the differential of a map between manifolds as we did with coordinate maps. Given a map  $\Phi: \mathcal{M} \to \mathcal{N}$  the differential is a linear map  $D\Phi_p:T_p\mathcal{M}\to T_q\mathcal{N}$  which maps tangent vectors at  $p\in\mathcal{M}$  to tangent vectors at  $q = \Phi(p) \in \mathcal{N}$ .



## **Equivalent definition**

A diffeomorphism  $\Phi: \mathcal{M} \to \mathcal{N}$  is an isometry iff it preserves angles:

$$\langle v, w \rangle_{T_p \mathcal{M}} = \langle D\Phi_p v, D\Phi_p w \rangle_{T_q \mathcal{N}}$$

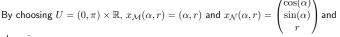
**Proof (only one direction):** Let  $c:[0,1] \to \mathcal{M}$  be a shortest curve connecting  $p \in \mathcal{M}$  and  $q \in \mathcal{M}$ :  $d_{\mathcal{M}}(p,q) = L(c) = \int_0^1 \|\dot{c}(t)\| dt$ . Then the curve  $d = \Phi \circ c : [0,1] \to \mathcal{N}$  has length

$$L(d) = \int_0^1 \left\| \frac{d}{dt} (\Phi \circ c(t)) \right\| dt = \int_0^1 \left\| D\Phi_{c(t)} \dot{c}(t) \right) \right\| dt = \int_0^1 \left\| \dot{c}(t) \right\| dt = L(c)$$

Since there is no shorter curve connecting  $\Phi(p)$  and  $\Phi(q)$  (why?), it follows  $d_{\mathcal{N}}(p, q) = d_{\mathcal{N}}(\Phi(p), \Phi(q)).$ 

Reason: the length of all (not only shortest) curves are preserved

## Example (Cylinder)

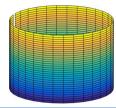


observing

$$g_{\mathcal{M}}(\alpha, r) = g_{\mathcal{N}}(\alpha, r)$$

for all 
$$(\alpha,r)\in U$$

we know that the stripe U of  $\mathbb{R}^2$  and the (sliced!) cylinder are isometric.



## **Eigenvalues and -vectors**



In the next weeks (starting today) we will make a lot use of the concept of eigenvalues and eigenvectors. Let us briefly revise the definitions and fundamental properties of eigenvalues and -vectors

**Definition.** Given a matrix  $\mathbf{A} \in \mathbb{C}^{n \times n}$ . If a pair  $\lambda \in \mathbb{C}, 0 \neq \mathbf{v} \in \mathbb{C}^n$  satisfies

$$\mathbf{A}\mathbf{v} = \lambda \mathbf{v}$$

we call v an eigenvector of A and  $\lambda$  its corresponding eigenvalue.

The n Eigenvalues  $\{\lambda_1, \ldots \lambda_n\}$  are the roots of  $\mathbf{A}$ 's characteristic polynomial

$$p_{\mathbf{A}}(\lambda) = \det(\mathbf{A} - \lambda \mathbf{I}) = \prod_{i=1}^{n} (\lambda - \lambda_i)$$

The fundamental theorem of algebra guarantees that this polynomial has exactly nroots (counted by multiplicity)

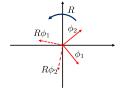
### Symmetric matrices

Eigendecompositions

**Theorem.** Given a symmetric matrix  $\mathbf{A} \in \mathbb{R}^{n \times n}$  all its eigenvalues are real and the corresponding eigenvectors can be chosen, such that

$$\langle \mathbf{v}_i, \mathbf{v}_j \rangle = \begin{cases} 0 & i \neq j \\ 1 & i = j \end{cases}$$

Notice that even under the assumption of orthonormality, the choice of eigenfunctions is not unique.



# $g_{\mathcal{N}}(u) = \langle Dx_{\mathcal{N}}(u), Dx_{\mathcal{N}}(u) \rangle$

Preserving intrinsics

Let  $\Phi: \mathcal{M} \to \mathcal{N}$  be an isometry,  $p \in \mathcal{M}$  and  $x_{\mathcal{M}}: \mathbb{R}^2 \supset U \to \mathcal{M}$  be a coordinate map of a neighborhood V of p. Then  $x_{\mathcal{N}}:=\Phi\circ x_{\mathcal{M}}:U\to \mathcal{N}$  is a coordinate map of the neighborhood  $\Phi(V) \subset \mathcal{N}$  of  $\Phi(p). \;\;$  For the first fundamental forms

$$g_{\mathcal{N}}(u) = \langle Dx_{\mathcal{N}}(u), Dx_{\mathcal{N}}(u) \rangle$$
  
=  $\langle D\Phi_{x_{\mathcal{M}}(u)} Dx_{\mathcal{M}}(u), D\Phi_{x_{\mathcal{M}}(u)} Dx_{\mathcal{M}}(u) \rangle = g_{\mathcal{M}}(u)$ 

Thus intrinsic properties of the shapes are preserved under isometries.



 $q_M: U \to \mathbb{R}^{2\times 2}, \ q_N: U \to \mathbb{R}^{2\times 2}$  we observe:



# Eigendecompositions

## Ugly facts about eigenvectors



Even for real matrices  $\mathbf{A} \in \mathbb{R}^{n \times n}$  the spectrum (set of eigenvalues) can contain complex values:

$$\mathbf{A} = \begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix}$$
$$p_{\mathbf{A}}(\lambda) = (\lambda^2 + 1) = (\lambda - i)(\lambda + i)$$

If v is a corresponding eigenvector of  $\lambda$ , then also every vector  $\mathbf{w} = \alpha \mathbf{v} \ (\alpha \neq 0)$  is an eigenvector to  $\lambda$ :

$$\mathbf{A}\mathbf{w} = \mathbf{A}\alpha\mathbf{v} = \alpha\mathbf{A}\mathbf{v} = \alpha\lambda\mathbf{v} = \lambda\alpha\mathbf{v} = \lambda\mathbf{w}$$

If  $\mathbf{v}_i, \mathbf{v}_j$  are eigenvectors to  $\lambda_i = \lambda_j$  then every linear combination  $\mathbf{w} = \alpha_i \mathbf{v}_i + \alpha_j \mathbf{v}_j$  of them is also an eigenvector (same holds for eigenvalues with higher multiplicity):

$$\mathbf{A}\mathbf{w} = \mathbf{A}(\alpha_i \mathbf{v}_i + \alpha_j \mathbf{v}_j) = \alpha_i \mathbf{A} \mathbf{v}_i + \alpha_j \mathbf{A} \mathbf{v}_j = \alpha_i \lambda_i \mathbf{v}_i + \alpha_j \lambda_i \mathbf{v}_j = \lambda_i \mathbf{w}$$



### Eigendecomposition 1

Isometries Eigendecompositions



If the eigenvectors  $\mathbf{v}_1,\dots,\mathbf{v}_n$  of A span  $\mathbb{R}^n$ , we can  $\operatorname{\mathbf{decompose}}$   $\mathbf{A}$  as

$$\mathbf{A} = \mathbf{V} \mathbf{\Lambda} \mathbf{V}^{-1}$$

with

$$\mathbf{V} = \begin{pmatrix} | & & | \\ v_1 & \cdots & v_n \\ | & & | \end{pmatrix} \qquad \qquad \mathbf{\Lambda} = \begin{pmatrix} \lambda_1 & \dots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \dots & \lambda_n \end{pmatrix}$$

$$\mathbf{A} = \begin{pmatrix} \lambda_1 & \dots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \dots & \lambda_n \end{pmatrix}$$

Let  $\mathbf{x} = \sum \alpha_i \mathbf{v}_i = \mathbf{V} \alpha$  be an arbitrary vector. Then

$$\mathbf{V}\boldsymbol{\Lambda}\mathbf{V}^{-1}\mathbf{x} = \mathbf{V}\boldsymbol{\Lambda}\begin{pmatrix} \alpha_1 \\ \vdots \\ \alpha_n \end{pmatrix} = \mathbf{V}\begin{pmatrix} \lambda_1\alpha_1 \\ \vdots \\ \lambda_n\alpha_n \end{pmatrix} = \sum_i \alpha_i\lambda_i\mathbf{v}_i = \sum_i \alpha_i\mathbf{A}\mathbf{v}_i = \mathbf{A}\mathbf{x}$$

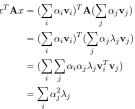
If the eigenvectors are orthonormal, we have  $\mathbf{V}^{-1} = \mathbf{V}^T.$ 

A symmetric matrix  ${f A}$  is positive definit iff all its eigenvalues are positive Let  $\{v_i\}$  be orthonormal eigenvectors of  $\mathbf{A}$  and  $0 \neq x = \sum_i \alpha_i \mathbf{v}_i$  an arbitrary vector.

Then

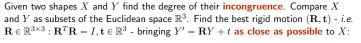
$$x^{T}\mathbf{A}x = (\sum_{i} \alpha_{i} \mathbf{v}_{i})^{T} \mathbf{A} (\sum_{j} \alpha_{j} \mathbf{v}_{j})^{T} \mathbf{A} (\sum_{j} \alpha_{j} \lambda_{j} \mathbf{v}_{j})^{T} (\sum_{j} \alpha_{j} \lambda_{j} \mathbf{v}_{j})^{T} (\sum_{j} \alpha_{j} \lambda_{j} \mathbf{v}_{i}^{T} \mathbf{v}_{j})$$
$$= (\sum_{i} \sum_{j} \alpha_{i} \alpha_{j} \lambda_{j} \mathbf{v}_{i}^{T} \mathbf{v}_{j})$$
$$= \sum_{j} \alpha_{j}^{2} \lambda_{j}$$

which is positive iff all  $\lambda_i$  are positive.



## Rigid Alignment

## Iterative closest point



$$d_{\mathsf{ICP}}(X,Y) = \min_{\mathbf{R},\mathbf{t}} d(\mathbf{R}Y + \mathbf{t}, X)$$

**Minimum:** extrinsic similarity of X and Y

 $\label{eq:minimizer:best rigid alignment between $X$ and $Y$} \label{eq:minimizer:best rigid}$ 

ICP is a family of algorithms differing in

- the choice of the shape-to-shape distance d
- the choice of the numerical minimization algorithm

### ICP algorithm



Given are a point cloud  $X = \{x_i\}$ , and an either discrete or continuous Y.

- Initialize Y
- Until convergence
  - Findest the best point-to-point correspondence  $y_i = \operatorname{argmin}_{y \in Y} \|x_i y\|$
  - Minimize the misalignment between corresponding points:  $(\mathbf{R}, \mathbf{t}) = \operatorname{argmin}_{\mathbf{R}, \mathbf{t}} \sum_{i} ||(Rx_i + t) - y_i||^2$

  - Update  $Y = \mathbf{R}Y + \mathbf{t}$

## Eigenvectors and optimization

There is a fundamental relation between the eigenvectors of a spd matrix  ${f A}$  and the optimization problem

$$\max \quad \mathbf{x}^T \mathbf{A} \mathbf{x}$$

s.t. 
$$\langle \mathbf{x}, \mathbf{x} \rangle = 1$$

Let  $\{\mathbf v_i\}$  be orthonormal eigenvectors of  $\mathbf A$  (ordered from big to small) and  $\mathbf{x} = \sum_i \alpha_i \mathbf{v}_i$ , satisfying  $\langle \mathbf{x}, \mathbf{x} \rangle = 1$ . First we observe that

$$\langle \mathbf{x}, \mathbf{x} \rangle = \langle \sum_{i} \alpha_{i} \mathbf{v}_{i}, \sum_{j} \alpha_{j} \mathbf{v}_{j} \rangle = \sum_{i} \alpha_{i}^{2} = 1$$

For the objective we get:

$$\mathbf{x}^T \mathbf{A} \mathbf{x} = (\sum_i \alpha_i \mathbf{v}_i)^T A \sum_j \alpha_j \mathbf{v}_j = \sum_i \alpha_i^2 \lambda_i \leqslant \lambda_1 = \mathbf{v}_1^T \mathbf{A} \mathbf{v}_1$$

Thus maximizing the quadratic function is equivalent to finding the principal eigenvector v<sub>1</sub>

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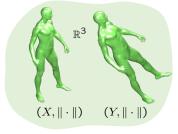
## **Euclidean isometry**



 $(X,d_X)$ 



 $(Y, d_Y)$ 



Intrinsic isometry

Two different metric spaces

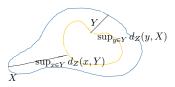
**Euclidean isometry** 

Part of the same metric space

## Shape-to-shape distance

The Hausdorff distance  $d_H$  between two subsets  $X,Y\subset Z$  of a metric space  $(Z, d_Z)$  is defined as

$$d_H(X,Y) = \max\{\sup_{y \in Y} d_Z(y,X), \sup_{x \in X} d_Z(x,Y)\}$$



Non-symmetric version:  $\sup_{y \in Y} d_Z(y, X)$ 

The "maximum"-version is sensitive to outliers. A variant is to use some kind of average:  $d(X,Y) = \int_{Y} d_Z^2(y,X) dn$ .

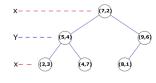


To perform the nearest neighbor search it is beneficial to make use of efficient datastructures such as k-d trees.

**KD** Tree

- binary tree
- each non-leaf node encodes a hyperplane
- Construction in  $O(n \log n)$
- Average query time  $O(\log n)$
- Course of dimensionality: in efficient for  $n < 2^k$





## Optimal Rigid alignment 1

For simplicity we assume that X and Y are centered at the origin:  $\sum_{i} x_{i} = \sum y_{i} = 0$ . Thus the second term in

$$\sum_{i} \left\| \mathbf{R} x_i - t - y_i \right\|^2 = \sum_{i} \left\| \mathbf{R} x_i - y_i \right\|^2 - 2 \langle t, \sum_{i} (\mathbf{R} x_i - y_i) \rangle + n \left\| t \right\|^2$$

vanishes and it follows t = 0 (or in general  $t = \sum_i x_i - \sum_i y_i$ ). It remains to find the orthogonal matrix  ${f R}$  minmimizing

$$\sum_{i} \|\mathbf{R}x_{i}\| - 2\sum_{i} y_{i}^{T} \mathbf{R}x_{i} + \sum_{i} \|y_{i}\|^{2} = \sum_{i} \|x_{i}\| - 2\sum_{i} y_{i}^{T} \mathbf{R}x_{i} + \sum_{i} \|y_{i}\|^{2}$$

The first and the last term are independent of  ${f R}$ , we thus have to maximize  $\sum_{i} y_{i}^{T} \mathbf{R} x_{i}$ .

## Optimal Rigid alignment 3



$$\sum_{i} y_i^T \mathbf{R} x_i = \sum_{i} \operatorname{tr}(\mathbf{R} x_i y_i^T) = \operatorname{tr}(\mathbf{R}^T \sum_{i} y_i x_i^T) = \operatorname{tr}(\mathbf{R}^T \mathbf{M})$$

If  $\mathbf{M}$  has full rank, we can construct

$$S = \sqrt{\mathbf{M}^T \mathbf{M}} \quad \mathbf{U} = \mathbf{M} \mathbf{S}^{-1}$$

such that  $\mathbf{M} = \mathbf{U}\mathbf{S}$  is a decomposition of  $\mathbf{M}$  in an orthogonal matrix  $\mathbf{U}$  and a positive definit matrix S, thus  $tr(\mathbf{R}^T\mathbf{M}) = tr(\mathbf{R}^T\mathbf{U}\mathbf{S})$ . The orthogonal matrix  $\mathbf{R}$ maximizing this term equals  $\mathbf{U}$ .

 $\mathbf{proof:}\ \ \mathbf{Let}\ \mathbf{S} = \sum_i \lambda_i v_i v_i^T$  (eigenvalues and eigenvectors). Then

$$\operatorname{tr}(\mathbf{R}^T\mathbf{U}\mathbf{S}) = \operatorname{tr}(\mathbf{R}^T\mathbf{U}\sum_i \lambda_i v_i v_i^T) = \operatorname{tr}(\sum_i \lambda_i (\mathbf{R}^T\mathbf{U}v_i)^T v_i) \leqslant \operatorname{tr}(\sum_i \lambda_i v_i v_i^T)$$

## Principal component analysis 1

Given a pointset  $X = \{x_i\}_{i=1}^n$  we want to align it with a rigid motion  $X \to \mathbf{R}(X - \mathbf{t})$  such that:

- the center of mass lies at the origin
- the direction in which the pointset expands the most should be the  $x_1$ -axis and so forth

By translating the center of mass of the point set to the origin, the first goal is easily achieved:  $\mathbf{t} = \sum_i x_i$ 

Now that the pointset is centered at the origin it remains to find the orthogonal matrix  $\ensuremath{\mathbf{R}}$  aligning it with the axis. Assume we know the three principal components  $d_1, d_2, d_3$  with which it is aligned before the

rotation, then choosing  $\mathbf{R} = \begin{pmatrix} 1 & 1 & 1 \\ d_1 & d_2 & d_3 \\ 1 & 1 & 1 \end{pmatrix}$ 







 $\sum_{i} y_i^T \mathbf{R} x_i = \sum_{i} \operatorname{tr}(\mathbf{R} x_i y_i^T) = \operatorname{tr}(\mathbf{R}^T \sum_{i} y_i x_i^T) = \operatorname{tr}(\mathbf{R}^T \mathbf{M})$ 

If  $\mathbf{M}$  has full rank, we can construct

$$S = \sqrt{\mathbf{M}^T \mathbf{M}} \quad \mathbf{U} = \mathbf{M} \mathbf{S}^{-1}$$

where the square root of a symmetric positive definit (spd) matrix  $\mathbf{A} = \mathbf{V} \mathbf{\Lambda} \mathbf{V}^T$  is defined as

$$\sqrt{\mathbf{A}} = \mathbf{V}\sqrt{\mathbf{\Lambda}}\mathbf{V}^T$$

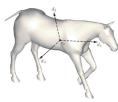
such that it holds  $\sqrt{\mathbf{A}}^T \sqrt{\mathbf{A}} = \mathbf{A}$ . One can show that the optimal choice is  $\mathbf{R}=\mathbf{U}.$ 

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## Drawbacks of ICP

Although a very simpe algorithm, ICP relies on a good initialization.

As an alternative to ICP and/or an initializier one could bring the shapes into a "canonical" pose. This canonical pose can be found using principal component analysis (PCA).



## Principal component analysis 2

We are looking for a direction  $\mathbf{d}$  ( $\|\mathbf{d}\|=1)$  maximizing

$$\sum_{i=1}^{n} \langle \mathbf{d}, \mathbf{x}_i \rangle^2 = \sum_{i=1}^{n} \mathbf{d}^T \mathbf{x}_i \mathbf{x}_i^T \mathbf{d}$$



The covariance matrix  $\Sigma_X$  is defined as  $\Sigma_X = \sum_{i=1}^n \mathbf{x}_i \mathbf{x}_i^T$  and is spd. So we can rewrite the objective as

$$\max \quad \mathbf{d}^T \mathbf{\Sigma}_X \mathbf{d}$$
  
s.t.  $\langle \mathbf{d}, \mathbf{d} \rangle = 1$ 

We have seen that this objective is maximized by the principal eigenvector of  $\Sigma_X$ .  $d_2$  and  $d_3$  are then the eigenvectors corresponding to second and third eigenvalue (when ordered by magnitude).