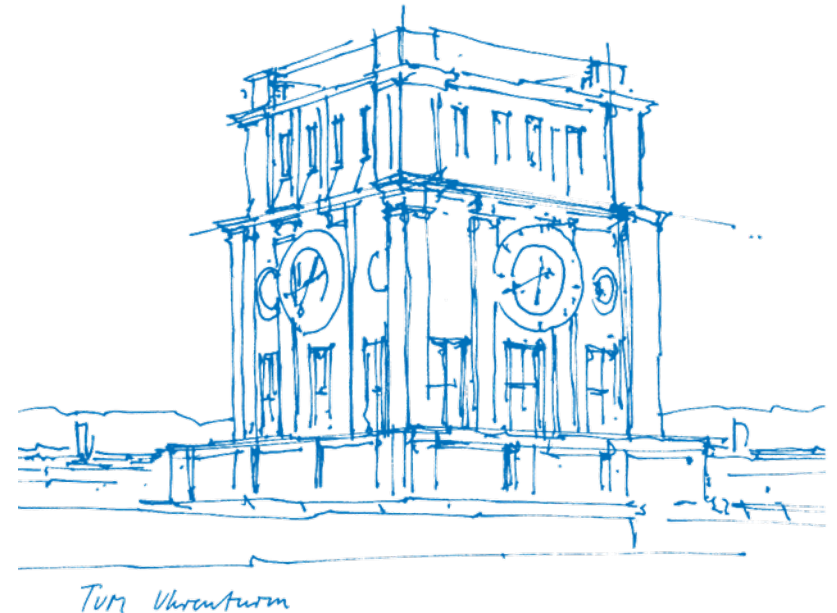




II : Graphical Model Representation

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Outline of the Chapter

- Bayesian network (directed graphical model).
- Markov random field (undirected graphical model).
- Independence assumption, representation power, parameterization, etc.



Bayesian Network

Bayesian Network

A **Bayesian network** (BN) is a *directed acyclic graph* $\mathcal{G} = (\mathcal{V}, \mathcal{E})$ together with:

- Random variables $X = (X_i)_{i \in \mathcal{V}}$ over \mathcal{V} ;
- A (joint probability) distribution P factorized as a product of conditional probability distributions (CPDs):

$$p(x) = \prod_{i \in \mathcal{V}} p(x_i | (x_j)_{j \in \text{Pa}_{\mathcal{G}}(i)}),$$

where $\text{Pa}_{\mathcal{G}}(i) = \{j \in \mathcal{V} : (j, i) \in \mathcal{E}\}$ consists of parents of i in \mathcal{G} .

Example "Student"

$$P(D, I, G, S, L) = P(D)P(I)P(G|D, I)P(S|I)P(L|G).$$

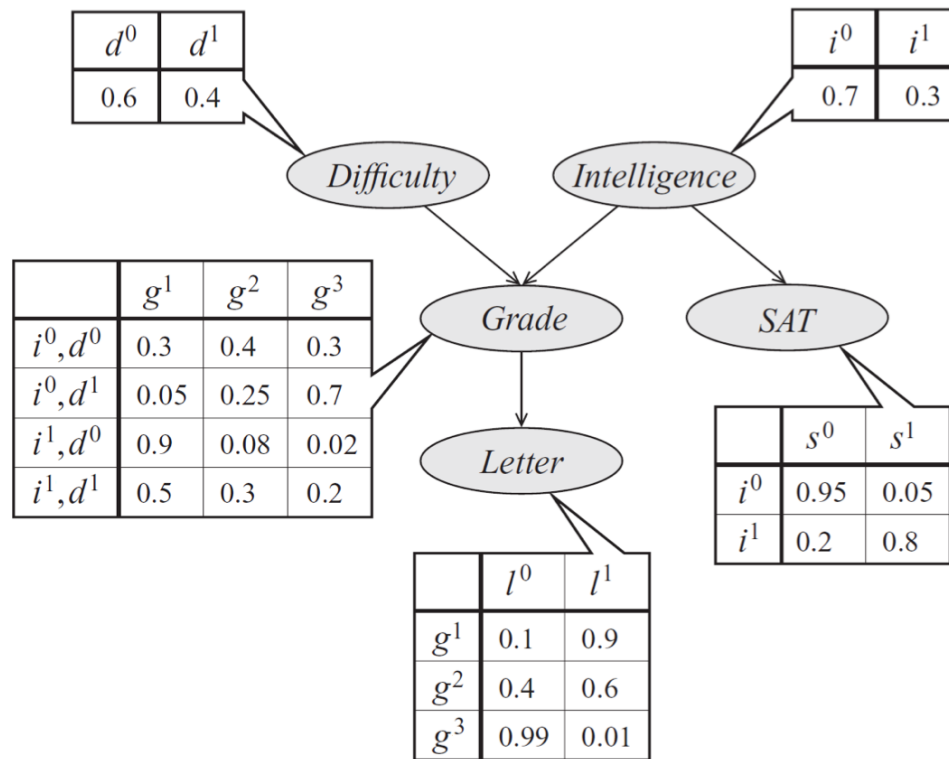


Figure: Bayesian network represented in probability tables.

Model Complexity

Consider BN representation for RVs $(X_i)_{i=1}^n$.

- If each RV X_i takes at most d outcomes and has at most k parents, then representation of

$$p(x_i | (x_j)_{j \in \text{Pa}_{\mathcal{G}}(i)})$$

requires $O(d^{k+1})$ free parameters.

- Since the joint distribution for $(X_i)_{i=1}^n$ is a product of n CPDs, the overall model complexity for BN is $O(nd^{k+1})$.
- Compared to a naive representation for the joint distribution which requires $O(d^n)$ parameters (typically $n \gg k$).

The reduction of complexity is due to the underlying independence assumptions.

Independence in BN

- For a distribution P for RVs (X_i) , we denote by $\mathcal{I}(P)$ the set of all **independence assumptions (assertions)** that hold in P :

$$\mathcal{I}(P) = \{(X_i \perp X_j \mid X_k)\}.$$

Recall conditional independence: $X_i \perp X_j \mid X_k$ iff

$$p(x_i, x_j \mid x_k) = p(x_i \mid x_k)p(x_j \mid x_k).$$

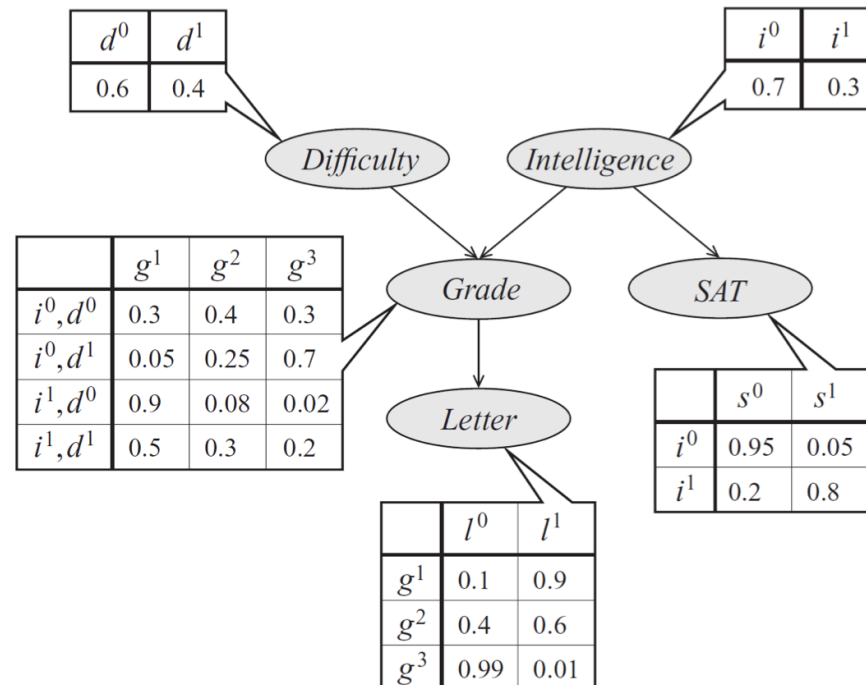
- BN \mathcal{G} implies **local independencies**:

$$\mathcal{I}_\ell(\mathcal{G}) = \left\{ \left(X_i \perp (X_j)_{j \in \text{NonDes}_{\mathcal{G}}(i)} \mid (X_k)_{k \in \text{Pa}_{\mathcal{G}}(i)} \right) \right\},$$

where $\text{NonDes}_{\mathcal{G}}(i)$ denotes the non-descendants of i (including i itself) in \mathcal{G} .

Example "Student"

$$\mathcal{I}_\ell(\mathcal{G}) = \left\{ \left(X_i \perp (X_j)_{j \in \text{NonDes}_\mathcal{G}(i)} \mid (X_k)_{k \in \text{Pa}_\mathcal{G}(i)} \right) \right\}.$$



In the above example, we have: $(L \perp \{I, D, S\} \mid G), (G \perp S \mid \{I, D\}) \in \mathcal{I}_\ell(\mathcal{G})$.

Beyond Local Independence

- Does \mathcal{G} encode other independence assertions besides $\mathcal{I}_\ell(\mathcal{G})$? (Yes.)
- How to identify whether a specific independence assertion holds in \mathcal{G} ?

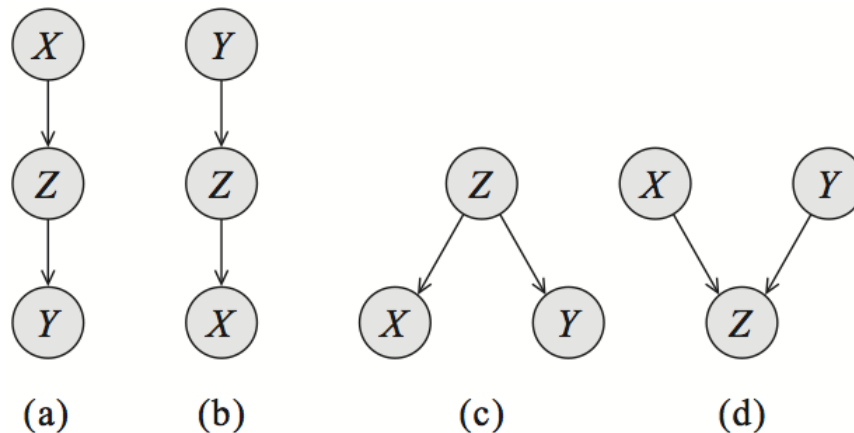
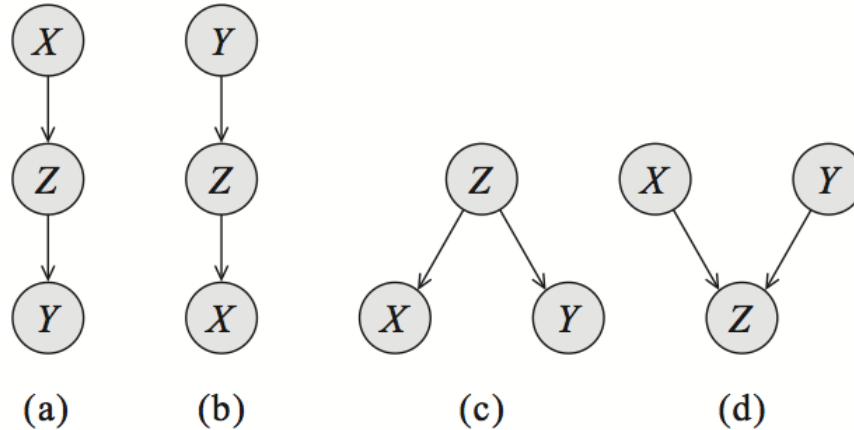


Figure: Two-edge trails from X to Y via Z . (d) is called the **v-structure**.

In the above figure, dependence flows from X to Y if the **trail** $X \leftrightarrow Z \leftrightarrow Y$ is **active**. This is the case if:

- In (a)–(c), Z is unobserved.
- In a v-structure such as (d), Z or one of descendants of Z is observed.

Active Trail



Let $X_1 \leftrightarrow X_2 \leftrightarrow \dots \leftrightarrow X_n$ be a trail in a BN \mathcal{G} , and Z be a set of observed nodes (RVs). The trail is **active** given Z if

- Whenever there is a v-structure (case (d)) in the trail $X_{i-1} \leftrightarrow X_i \leftrightarrow X_{i+1}$, then X_i or one of its descendants are in Z .
- No other node along the trail is in Z .

Intuitively, information/dependence flows from X_1 to X_n (and vice versa) through the active trail $X_1 \leftrightarrow X_2 \leftrightarrow \dots \leftrightarrow X_n$.

D-separation, Global Independence

Let X, Y, Z be three sets of nodes in a BN \mathcal{G} . If there is no active trail between any node in X and Y given Z , we say X and Y are **d-separated** given Z .

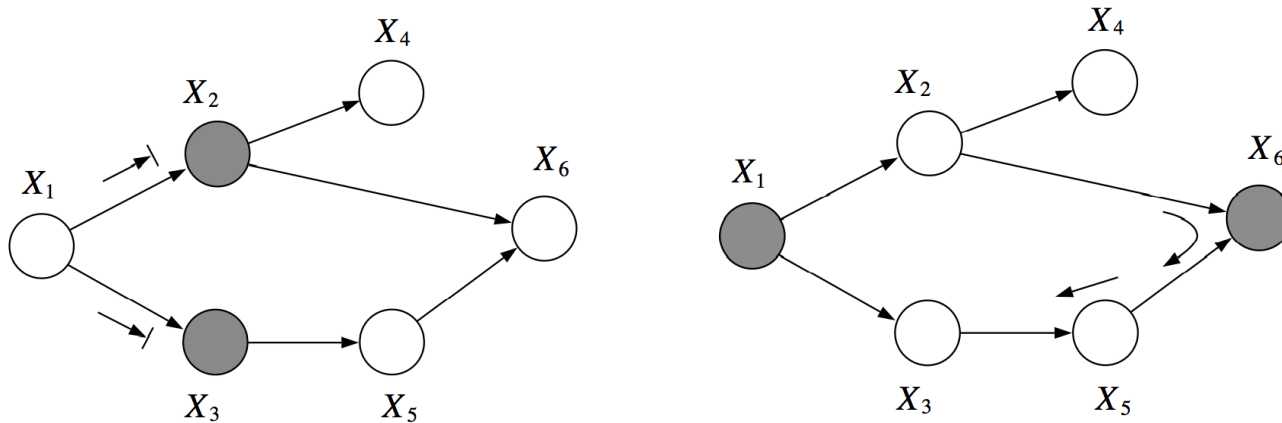


Figure: (left) X_1 and X_6 are d-sep. given $\{X_2, X_3\}$; (right) X_2 and X_3 are not d-sep. given $\{X_1, X_6\}$.

We denote by $\mathcal{I}(\mathcal{G})$ the set of **global Markov independencies**:

$$\mathcal{I}(\mathcal{G}) = \{(X \perp Y \mid Z) : X \text{ and } Y \text{ are d-separated given } Z\}.$$

Facts about D-separation

- F1. (soundness)** If a distribution P factorizes according to \mathcal{G} , then $\mathcal{I}(\mathcal{G}) \subset \mathcal{I}(P)$. In this case, we call \mathcal{G} an **I-map** for P .
- F2. (sharpness)** If nodes X and Y are not d-separated given Z in \mathcal{G} , then X and Y are dependent given Z in some distribution P that factorizes over \mathcal{G} .
- F3. (completeness)** When a distribution P factorizes according to \mathcal{G} , $\mathcal{I}(\mathcal{G}) = \mathcal{I}(P)$ does not necessarily hold. Obviously, one can add superfluous edges to \mathcal{G} s.t. $\mathcal{I}(\mathcal{G}) \subsetneq \mathcal{I}(P)$.

$p(b a)$	b_0	b_1
a_0	0.4	0.6
a_1	0.4	0.6

Figure: A possible I-map for P is $A \rightarrow B$, but $\emptyset = \mathcal{I}(\mathcal{G}) \subsetneq \mathcal{I}(P)$. Such cases almost surely do not happen.

Remark: For almost all P (except for a set of measure zero in the space of CPD parameterizations) for which \mathcal{G} is an I-map, we have $\mathcal{I}(\mathcal{G}) = \mathcal{I}(P)$.

I-equivalence

- Two BNs \mathcal{G}_1 and \mathcal{G}_2 are said to be **I-equivalent** if $\mathcal{I}(\mathcal{G}_1) = \mathcal{I}(\mathcal{G}_2)$.
- The **skeleton** of a BN $\mathcal{G} = (\mathcal{V}, \mathcal{E})$ is an undirected graph $(\mathcal{V}, \mathcal{E}')$ such that $\{X, Y\} \in \mathcal{E}'$ whenever $(X, Y) \in \mathcal{E}$.
- Fact: If two BNs have the same skeleton and the same set of v-structures, then they are I-equivalent.

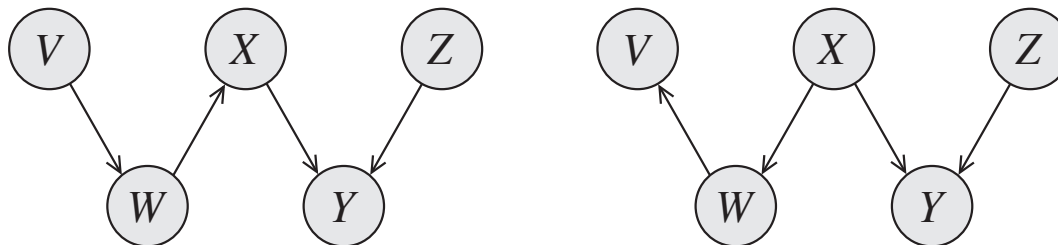


Figure: Example of two I-equivalent BNs.

Perfect Map and Counterexamples

- We say a BN \mathcal{G} is a **perfect map** for a distribution P if $\mathcal{I}(\mathcal{G}) = \mathcal{I}(P)$.
- Certain independencies cannot be expressed perfectly by BN.

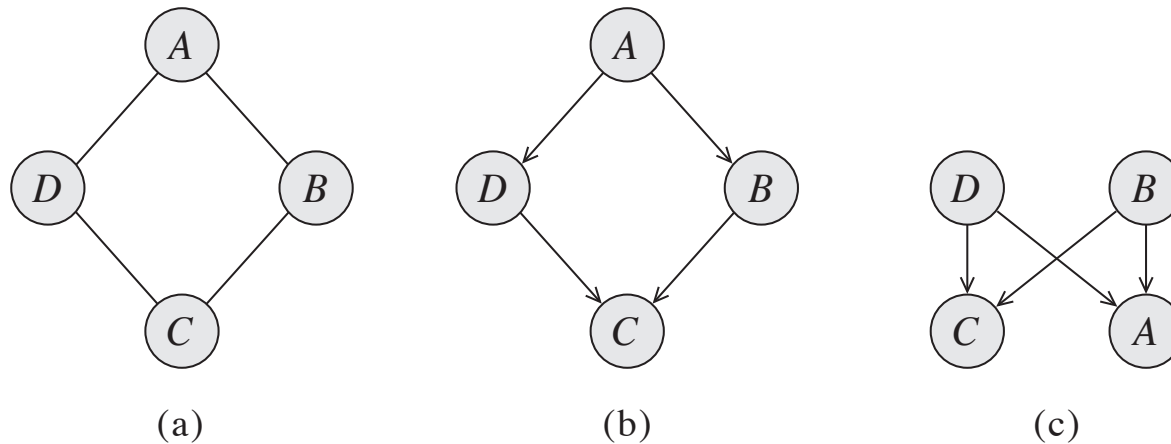


Figure: A counterexample where a perfect map does not exist.

- (a) Desired independence assertions: $A \perp C \mid \{B, D\}$, $B \perp D \mid \{A, C\}$.
- (b) In this BN: $(A \perp C \mid \{B, D\}) \in \mathcal{I}(\mathcal{G})$, but $(B \perp D \mid \{A, C\}) \notin \mathcal{I}(\mathcal{G})$.
- (c) Again, $(A \perp C \mid \{B, D\}) \in \mathcal{I}(\mathcal{G})$, but $(B \perp D \mid \{A, C\}) \notin \mathcal{I}(\mathcal{G})$.

Topics which are not covered here ...

- Algorithm for detecting d-separation in a BN \mathcal{G} .
- Algorithm for finding minimal I-map \mathcal{G} for a given distribution P .
- Algorithm for finding perfect map \mathcal{G} (if exists) for a given distribution P .
- Further reading: Koller & Friedman, Chapter 3.