

# Combinatorial Optimization in Computer Vision

## Chapter 8: Labelling with Ordering Constraint

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# Labelling Problems with Ordering Constraint

**Problem:** In general, the labelling problem with ordering constraint deals with the assignment of a label in  $\mathcal{L}$  for each pixel of an image such that the final labelling  $x : V \rightarrow \mathcal{L}$  satisfies certain **geometric constraints**.

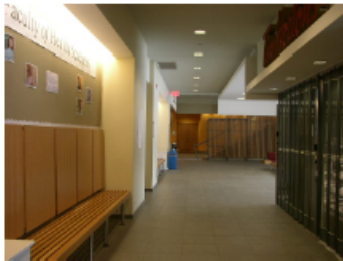
Examples include:

- $\mathcal{L}$  comprises parts of a class of objects. Then one might want to impose constraints as: „a *car wheel* cannot be below a *car roof*“.
- „The *sky* is above the *street*.“

# Geometric Scene Labelling

Here, we will discuss an approach by Liu, Samarabandu and Veksler who consider the specific problem of finding a multilabel segmentation  $x : V \rightarrow \mathcal{L}$  with

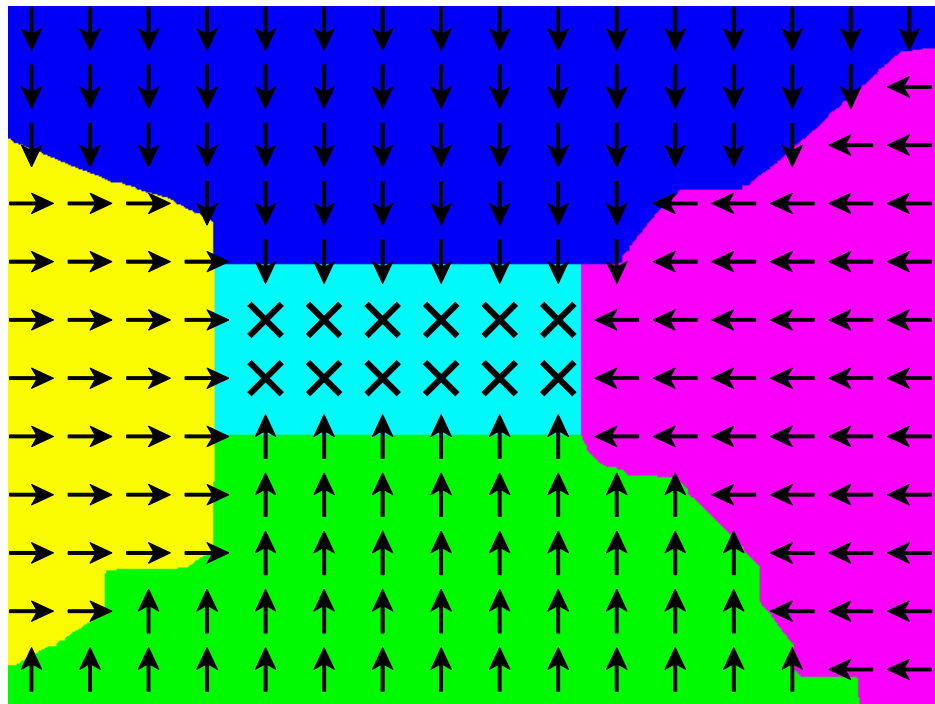
$$\mathcal{L} := \{C = \text{center}, L = \text{left}, R = \text{right}, B = \text{bottom}, T = \text{top}\}.$$



# Constraints on Geometric Scene Labels

There are some obvious constraints on these labels like

- „left“ cannot appear right from „center“ or „right“
- „top“ cannot appear below „center“ or „bottom“



# MRF Model

The resulting problem is formulated using an MRF model

$X = (X_v)_{v \in V}$  where  $V$  is the set of pixels and the  $X_v$  take values in  $\mathcal{L}$ .

The corresponding Gibbs potential is given by

$$E(x) = \underbrace{\sum_{v \in V} \phi_v(x(v))}_{\text{data term}} + \underbrace{\sum_{v \sim w} a_{vw} \psi(x(v), x(w))}_{\text{regularizer}} + \underbrace{\sum_{v \sim w} \theta_{v,w}(x(v), x(w))}_{\text{constraints}}.$$

Here, the data term is learned from a database of images and the Potts model is used for the regularizer.

# Constraints in the Gibbs Potential

The constraints are incorporated into the Gibbs potential by penalizing two neighboring label assignments, which violate a constraint, with penalty  $\infty$ .

For example, if  $v$  is the left neighbor pixel of  $w$ , then the table for  $\theta(\alpha, \beta)$  looks like (the label of  $v$  is indicated in the row)

	$L$	$R$	$C$	$T$	$B$
$L$	0	$\infty$	0	0	0
$R$	$\infty$	0	$\infty$	$\infty$	$\infty$
$C$	$\infty$	0	0	$\infty$	$\infty$
$T$	$\infty$	0	$\infty$	0	$\infty$
$B$	$\infty$	0	$\infty$	$\infty$	0

# Minimization of Gibbs Potential

The Gibbs potential

$$E(x) = \underbrace{\sum_{v \in V} \phi_v(x(v))}_{\text{data term}} + \underbrace{\sum_{v \sim w} a_{vw} \psi(x(v), x(w))}_{\text{regularizer}} + \underbrace{\sum_{v \sim w} \theta_{v,w}(x(v), x(w))}_{\text{constraints}}.$$

can be approximately minimized using the expansion algorithm.

However, in this case, there is no quality guarantee (because the pairwise potential is not a metric).

# Order-Preserving Moves

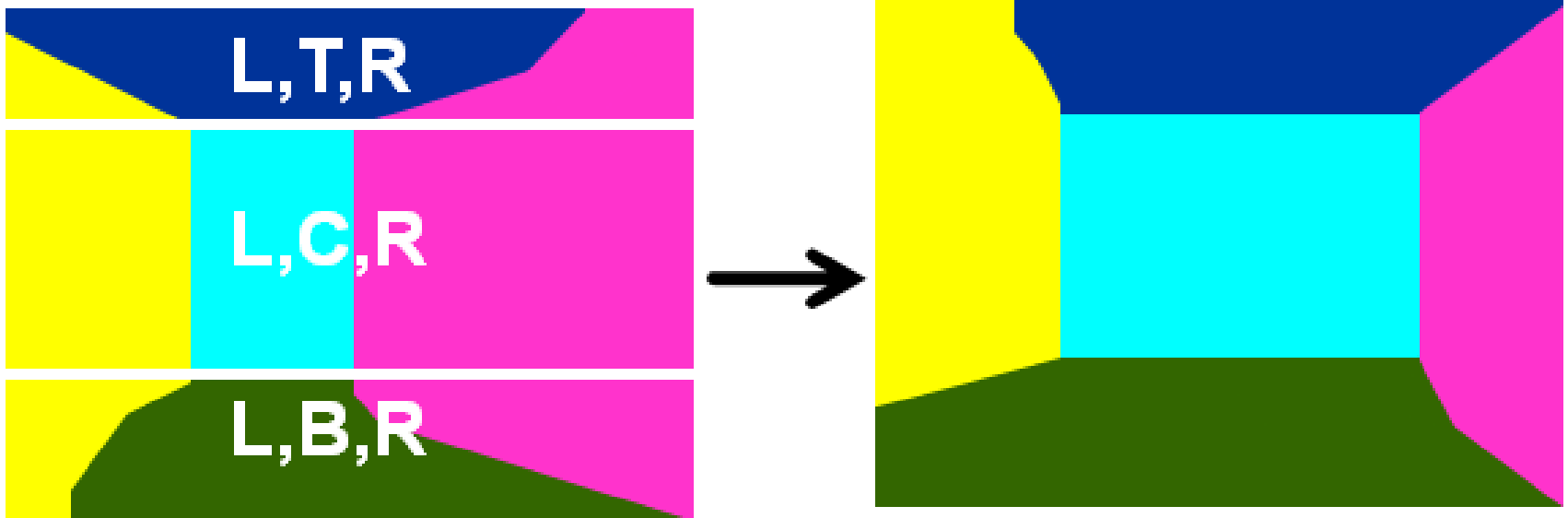
It turns out that the expansion algorithm converges to bad local minima in this example. Thus, a new class of moves is introduced, the so-called **order-preserving moves**.

**Horizontal Moves:** Region  $C$  can change in horizontal but not in vertical direction.

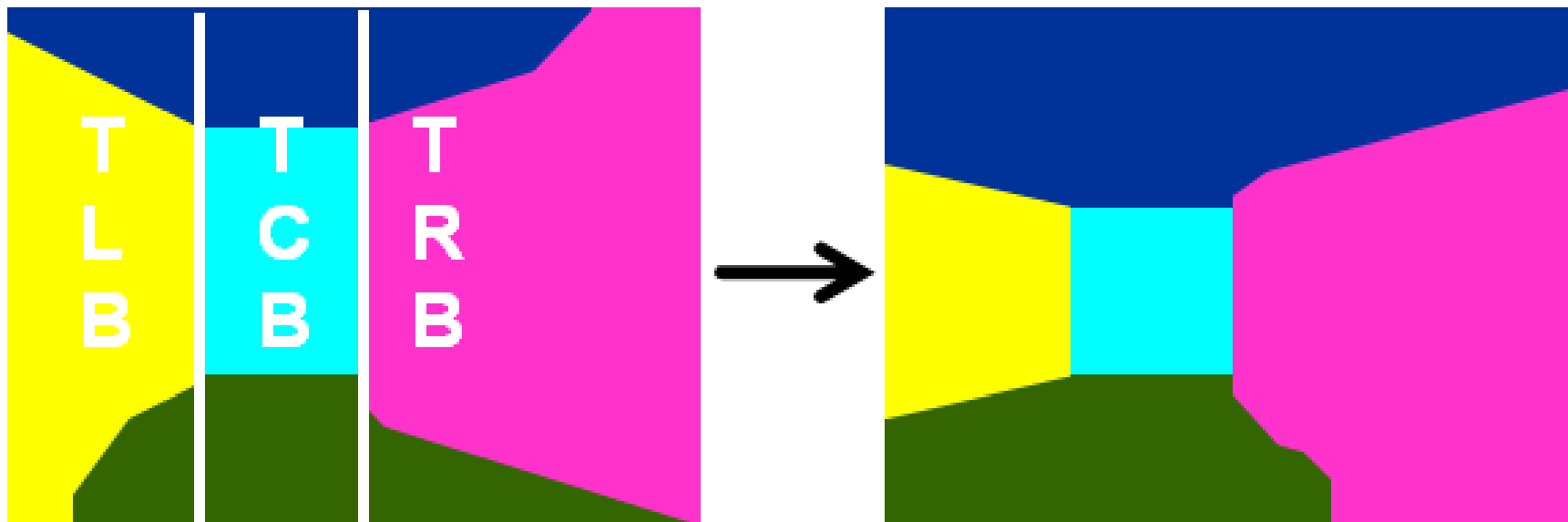
**Vertical Moves:** Region  $C$  can change in vertical but not in horizontal direction.



# Horizontal Moves



# Vertical Moves



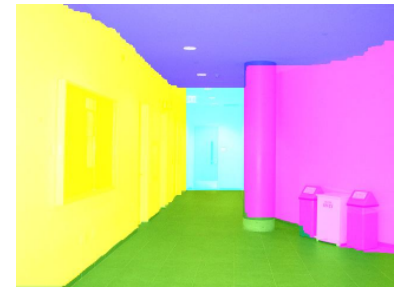
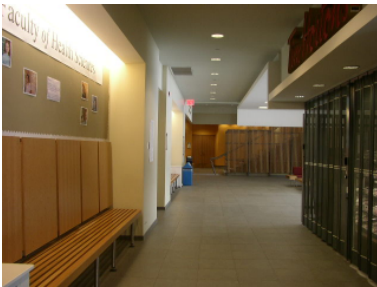
# Order-Preserving Move Algorithm

1. Initialize  $x(v) \leftarrow C \quad \forall v \in V$ .
2.  $x_h \leftarrow$  optimal horizontal move from  $x$   
 $x_v \leftarrow$  optimal vertical move from  $x$
3. If  $E(x_h) < E(x_v)$ , set  $x \leftarrow x_h$ , else  $x \leftarrow x_v$ .
4. Repeat
  - $x_h \leftarrow$  optimal horizontal move from  $x$
  - $x \leftarrow$  optimal vertical move from  $x_h$until  $E(x)$  stops decreasing.

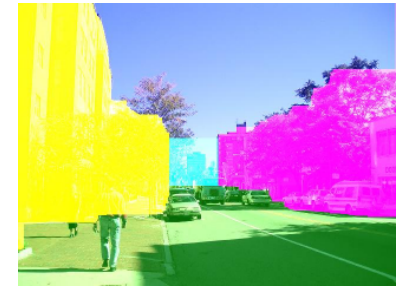
# Comments

- Similarly to the expansion algorithm, this method produces a **local optimum**.
- Unfortunately, no quality guarantee is available.
- The core step is the search for the optimal horizontal or vertical move. In each of these, a multilabel problem over more than 2 labels has to be solved. A result by Schlesinger and Flach shows that this can be done in polynomial time by computing a cut in an appropriate graph.

# Results: Indoor Images



# Results: Outdoor Images



# Scene Walk using the Geometric Scene Labels



# Summary

- Frequently we want to impose geometric constraints on labellings.
- Liu, Samarabandu and Veksler propose a move-make algorithm for computing local optima of the corresponding MRF energy.
- Similarly to the expansion algorithm, the core step consists in computing the optimal move from a given labelling. Again, this is solved by computing a cut in an appropriate graph.