Combinatorial Optimization in Computer Vision Lecture: F. R. Schmidt and C. Domokos Exercises: T. Möllenhoff and T. Windheuser Winter Semester 2015/2016 Computer Vision Institut für Informatik

## Weekly Exercises 7

Room: 02.09.023 Tuesday, 09.12.2015, 14:15-15:45 Submission deadline: Tuesday, 09.12.2015, 11:15, Room 02.09.023

### Convexity

# (2 Points)

**Exercise 1** (2 Points). Show that the following functions are convex:

- a)  $f : \mathbb{R} \to \mathbb{R}, \ x \mapsto |x|^p, \ p \in [1, \infty).$  c)  $f : \mathbb{R}_{\geq 0} \to \mathbb{R}, \ x \mapsto x \log(x).$
- b)  $f : \mathbb{R} \to \mathbb{R}, x \mapsto \log(1 + \exp(x)).$  d)  $f : \mathbb{R} \to \mathbb{R}, x \mapsto \max(0, x).$

#### Solution.

a) We have f(x) = g(h(x)), with  $g : \mathbb{R}_{\geq 0} \to \mathbb{R}, x \mapsto x^p$  and h(x) = |x|. Since g is increasing and convex on  $\mathbb{R}_{\geq 0}$  and h is convex, their composition is convex as well.

b) 
$$f'(x) = \frac{\exp(x)}{1 + \exp(x)}, \ f''(x) = \frac{\exp(x)(1 + \exp(x)) - \exp(x)\exp(x)}{(1 + \exp(x))^2} = \frac{\exp(x)}{(1 + \exp(x))^2} \ge 0.$$

c) 
$$f'(x) = x\frac{1}{x} + \log(x) = \log(x), \ f''(x) = \frac{1}{x} \ge 0 \text{ for } x \ge 0$$

d) We want to prove that

$$\max(0, \alpha x + (1 - \alpha)y) \le \alpha \max(0, x) + (1 - \alpha) \max(0, y).$$

Clearly it holds

$$0 \le \alpha \max(0, x) + (1 - \alpha) \max(0, y),$$

and also

 $\alpha x + (1 - \alpha)y \le \alpha \max(0, x) + (1 - \alpha)\max(0, y).$ 

Hence the claim follows.

## Multi-Label Problems and Submodularity (8 Points)

Let  $\mathcal{L} = \{1, \ldots, \ell\} \subset \mathbb{N}, \ell \in \mathbb{N}$ , be a totally ordered label set and  $n \in \mathbb{N}$ . The total order on  $\mathcal{L}$  induces a partial order on  $\mathcal{L}^n$ . It is easy to check that  $\mathcal{L}^n$  is a distributive

lattice, where meet  $\wedge$  and join  $\vee$  are the component-wise minimum and maximum, respectively. I.e.

$$(x \wedge y)_i = (\min\{x_i, y_i\})_i \text{ and} (x \vee y)_i = (\max\{x_i, y_i\})_i$$
 (1)

for any  $x, y \in \mathcal{L}^n$ .

**Definition.** A function  $c: \mathcal{L}^n \to \mathbb{R}$  is called *submodular* if for any  $x, y \in \mathcal{L}^n$  it holds

$$c(x \wedge y) + c(x \vee y) \le c(x) + c(y).$$
(2)

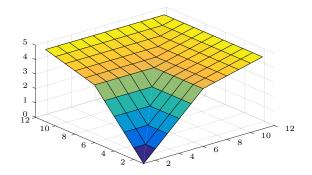
Exercise 2 (4 Points).

- a) Let  $g : \mathbb{R} \to \mathbb{R}$  be a convex function. Prove that the function  $c : \mathcal{L} \times \mathcal{L} \to \mathbb{R}, c(x, y) = g(|x y|)$ , is submodular.
- b) Let  $f_i : \mathcal{L} \to \mathbb{R}$  be some functions and let  $c_{i,j} : \mathcal{L} \times \mathcal{L} \to \mathbb{R}$  be some pair-wise submodular functions. Prove that the energy  $E : \mathcal{L}^n \to \mathbb{R}$ , given by

$$E(x) = \sum_{i} f_i(x_i) + \sum_{i,j} c_{i,j}(x_i, x_j),$$
(3)

is submodular.

c) Let  $a, b, c \in \mathbb{R}_{\geq 0}$  be some constants. Show that  $c : \mathcal{L} \times \mathcal{L} \to \mathbb{R}$ , where  $c(x, y) = \min\{a \max\{x, y\}, b \max\{x, y\} + c\}$ , is submodular but not necessarily convex.



**Exercise 3** (4 Points). Let  $E : \mathbb{B}^n \to \mathbb{R}$  be a quadratic submodular energy and  $x \in \mathbb{B}^n$ . Further, let  $x^1 \in \mathbb{B}^n$  be the result of the 1-expansion starting from x and let  $x^0 \in \mathbb{B}^n$  be the result of the 0-expansion starting from  $x^1$ .

- a) Show that there is a  $x^* \in \arg \min_z E(z)$  such that  $x^* \leq x^1$ .
- b) Show that  $x^0 \in \arg \min_z E(z)$ .

**Solution.** First let us recall the definition of the  $\alpha$ -Expansion starting from z. Given a current labeling  $z \in \mathcal{L}^n$  and label  $\alpha \in \mathcal{L}$  we obtain a minimizer  $y^* \in \mathbb{B}^n$  of the following energy:

$$E^{z,\alpha}(y) = \sum_{i=1}^{n} f_i^{z,\alpha}(y_i) + \sum_{i=1}^{n} \sum_{j \in \mathcal{N}(i)} f_{ij}^{z,\alpha}(y_i, y_j),$$

where

$$f_i^{z,\alpha}(0) = f_i(z_i), \qquad f_i^{z,\alpha}(1) = f_i(\alpha), \\
 f_{ij}^{z,\alpha}(0,0) = f_{ij}(z_i, z_j), \qquad f_{ij}^{z,\alpha}(0,1) = f_{ij}(z_i, \alpha), \\
 f_{ij}^{z,\alpha}(1,0) = f_{ij}(\alpha, z_j), \qquad f_{ij}^{z,\alpha}(1,1) = f_{ij}(\alpha, \alpha).$$

Subsequently, we set

$$z_i^{\text{new}} = \begin{cases} \alpha, & \text{if } y_i = 1, \\ z_i, & \text{otherwise.} \end{cases}$$

Clearly it holds that  $E^{z,\alpha}(y^*) = E(z^{\text{new}})$  by definition. Since  $y^*$  is a global minimizer,  $z^{\text{new}} \in \mathcal{L}^n$  is the optimal solution that can be obtained by setting some variables to the label  $\alpha$ .

a) Let us define the sets

$$A = \{ i \mid x_i^1 = 1 \}, \tag{4}$$

$$B = \{i \mid x_i^* = 1\},\tag{5}$$

$$C = \{i \mid x_i^0 = 1\},\tag{6}$$

$$X = \{ i \mid x_i = 1 \}, \tag{7}$$

$$S = A \cap B. \tag{8}$$

In the following we identify sets with binary vectors in  $\mathbb{B}^n$  as defined above, i.e., we interchangingly use E(A) for  $E(x^1)$ .

Since E is submodular, we have

$$E(A \cup B) \le E(A) + \underbrace{E(B) - E(S)}_{\le 0} \le E(A).$$
(9)

Since A is obtained by the 1-expansion starting from X, we have that  $X \subset A$  and that A is the set with the overall minimum energy containing X, i.e. we have that:

$$E(A) \le E(Y), \ \forall Y \supset X.$$

Since A is the set with the overall minimum energy containing X and clearly  $A \cup B$  also contains X and  $E(A \cup B) \leq E(A)$  we have that

$$E(A) = E(A \cup B).$$

Hence, it follows from (9) that

$$E(S) \le E(B),$$

which in turn shows that  $S \in \operatorname{argmin}_{z} E(z)$  since B is the global minimizer of E. Since  $S \subset A$  we have that  $x^* \leq x^1$  for some global minimizer  $x^* \in \operatorname{argmin}_{z} E(z)$ .

b) According to the definition of the 0-expansion starting from A, C is the set with the minimal energy having the property  $\overline{A} \subset \overline{C}$ , i.e.

$$E(C) \leq E(Y), \ \forall \overline{Y} \supset \overline{A}.$$

Since from a) we have  $B \subset A$  we have that  $\overline{B} \supset \overline{A}$  and hence  $E(C) \leq E(B)$ . Hence C is a global minimizer, i.e.  $x^0 \in \operatorname{argmin}_z E(z)$ .

## Programming

(10 Points)

**Exercise 4** (Image Denoising, 10 Points). Given a noisy input image  $I \in \mathcal{L}^n$  consisting of n pixels and  $\mathcal{L} = \{1, \ldots, \ell\}$  intensities, the goal of this exercise is to compute a denoised version  $x \in \mathcal{L}^n$  of I (see Figure 1).

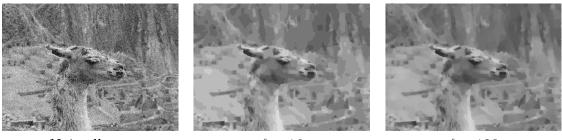
To this end, we want to minimize the following energy:

$$E(x) = \sum_{i=1}^{n} f_i(x_i) + \sum_{i=1}^{n} \sum_{j \in \mathcal{N}(i)} f_{ij}(x_i, x_j).$$
(10)

The unary potentials  $f_i : \mathcal{L} \to \mathbb{R}$  are chosen accordingly to the statistical noise model. In this exercise we assume Gaussian noise, i.e.  $f_i(x_i) = \frac{1}{2}(x_i - I_i)^2$ . The pairwise potentials  $f_{ij} : \mathcal{L} \times \mathcal{L} \to \mathbb{R}$  are used to model the prior knowledge about the input image. Here we pick  $f_{ij}(x_i, x_j) = \lambda |x_i - x_j|^p$  for some parameters  $\lambda > 0$ ,  $p \ge 1$ . Chose a simple 4-connected neighbourhood for  $\mathcal{N}$ .

Download  $07\_supp.zip$  from the lecture website and use denoising.cpp as a starting point which loads an image, converts it into grayscale and adds Gaussian noise with standard deviation  $\sigma$ . Extend the program in the following ways:

- a) Formulate the multilabel energy minimization problem (10) as a pseudo-Boolean optimization problem using the construction presented in the lecture. Use Kolmogorov's code<sup>1</sup> for graph cuts to find a global minimizer of the submodular pseudo-Boolean energy equivalent to (10). The graph construction of the dataterm is already provided in the code denoising.cpp and your task is to add the edges for the regularizer.
- b) Formulate the energy minimization problem (10) as a multilabel optimization problem in OpenGM and apply the  $\alpha$ -Expansion and  $\alpha$ - $\beta$ -Swap algorithms (these are already implemented in OpenGM). Compare the results from the different algorithms to the globally optimal result of exercise a).



Noisy llama

 $\ell = 16$ 

 $\ell = 128$ 

Figure 1: Examplary denoising results for p = 1 and  $\lambda = 0.25$  and different number of labels.

 $<sup>^{1}</sup>$ http://pub.ist.ac.at/~vnk/software.html