

Chapter 1

Why Convex Optimization

1.1 Introduction

Optimization problems arise naturally in many computer vision and machine learning applications that estimate pixel values, motions, shapes, or model parameters from input images, videos, range sensors, or training data. By formalizing the problem into a concise mathematical form, we obtain an optimization problem whose solution are the model parameters that best fit the observed data and our prior knowledge of the physical world. The next step, finding a solution to the mathematical model, is far from trivial.

The bitter truth is that most optimization problems are unsolvable. Among the solvable ones, convex problems form a large subset that builds on solid mathematical properties and can be solved efficiently with algorithms that exploit these properties. Most commercial packages for optimization, however, use minimal assumptions on the structure of the optimization in order to fit a large class of problems, albeit in a poor manner, and lead to poor optimization strategies. The goal of this course is to present techniques that exploit the properties of convex optimization problems to develop efficient algorithms for a large set of computer vision and machine learning problems.

In many of these applications the process of creating a model takes a considerable amount of time and effort. Therefore, it is important to understand the properties of the model and the computational consequences of each decision. Very often we have to choose between a *good* model, which we cannot solve and a *bad* model, which can be solved efficiently. To distinguish between the two, it is necessary to be aware of some theory that explains what we can and what we cannot do with optimization problems, and how convexity plays a key role on the solvability of a problem.

This first chapter is a summary of Chapter 1 of *Introductory Lectures on Convex Optimization*, by Nesterov.

1.2 Limitations in General Optimization

Let us start by describing our optimization problem. Let $u \in \mathbb{R}^n$ be an n -dimensional real vector, $C \subset \mathbb{R}^n$ be a subset of \mathbb{R}^n , and E be a real-valued functions of u . We study different variants of the following general minimization problem:

$$\hat{u} \in \arg \min_{u \in C} E(u) \tag{1.1}$$

The function $E : \mathbb{R}^n \rightarrow \mathbb{R}$ is the objective function, while the set C is the feasible set. We consider a minimization problem by convention, but we can also consider a maximization problem with $-E$ as objective function.

There is a natural classification of the types of minimization problems that we will study: unconstrained problems where $C = \mathbb{R}^n$, smooth problems where E is differentiable, and non-smooth problems where E is not differentiable. We also distinguish two different types of solutions to the minimization problem.

Definition Global Minimum. u^* is a global solution of (1.1) if $E(u^*) \geq E(u)$ for all $u \in C$.

Definition Local Minimum. u^* is a local solution of (1.1) if there exists a $r > 0$ such that

$$E(u^*) \geq E(u) \quad \forall u \in C, \quad \|u - u^*\| < r.$$

Local minima are easier to find than global ones. For instance, given an estimate of the the minimizer u^0 , we can create a sequence $\{u^k\}$ that decreases the value of the energy at each step to find the local minimum. Formally, we say that these type of optimization methods create a relaxation sequence $\{E(u^k)\}$ that satisfies $E(u^{k+1}) \leq E(u^k)$ and always improves the initial value of the objective function. If E is bounded below on \mathbb{R}^n , then the sequence $\{E(u^k)\}$ converges to a local minimum. Let us formalize what we mean by convergence.

Definition We say that a sequence $\{a^k\} \subset \mathbb{R}^n$ converges to $\hat{a} \in \mathbb{R}^n$ if for all $\epsilon > 0$ there exists an $k_0 \in \mathbb{N}$ such that

$$\|a^k - \hat{a}\| < \epsilon \quad \forall k \geq k_0.$$

To implement the idea of relaxation we use another fundamental principle of numerical analysis, the approximation. The approximation replaces the original objective function E by a simplified objective function that is close to the original. When the function is differentiable, we usually resort to local approximations of the objective function based on its Taylor expansion at the current estimate to create linear and quadratic approximations of the objective.

Let $E(u)$ be differentiable at u^0 , then for $u \in \mathbb{R}^n$, we have

$$E(u) = E(u^0) + \langle \nabla E(u^0), u - u^0 \rangle + o(\|u - u^0\|) \quad \text{where} \quad \lim_{r \rightarrow 0} \frac{o(r)}{r} = 0.$$

Function $E(u; u^0) = E(u^0) + \langle \nabla E(u^0), u - u^0 \rangle$ is a linear approximation of E in a neighborhood of u^0 . Given an initial estimate of the minimizer u^0 , we can then use this linear approximation to reduce the value of $E(u)$ in a neighborhood of u^0 . In particular we can decide to iteratively step in the direction of maximum descent as follows:

$$\begin{aligned} u^1 &= u^0 - \tau \nabla E(u^0) \\ u^2 &= u^1 - \tau \nabla E(u^1) \\ &\dots \\ u^{k+1} &= u^k - \tau \nabla E(u^k). \end{aligned}$$

This gives us a very simple algorithm know as gradient descent. We will see in this course that under certain conditions, the algorithm creates a relaxation sequence that decreases the value of the objective function and converges to a point $\hat{u} \in \mathbb{R}^n$. This a point then satisfies $\hat{u} = \hat{u} - \tau \nabla E(\hat{u}) \Rightarrow \nabla E(\hat{u}) = 0$. This is a necessary condition for optimality, as the next theorem shows.

Theorem 1. First-order Optimality Condition. *Let u^* be a local minimum of differentiable function $E(u)$. Then $\nabla E(u^*) = 0$.*

Proof. Since u^* is a local minimum of $E(u)$, then there exists $r > 0$ such that for all v with $\|v - u^*\| \leq r$, we have $E(v) \geq E(u^*)$. Since E is differentiable, this implies that

$$E(v) = E(u^*) + \langle \nabla E(u^*), v - u^* \rangle + o(\|v - u^*\|) \geq E(u^*).$$

Thus, for all s we have $\langle \nabla E(u^*), s \rangle \geq 0$. If we consider the directions s and $-s$, we get $\nabla E(u^*) = 0$. \square

Note that we have proved only a necessary condition of a local minimum. The points satisfying this condition are called the stationary points of function. In order to see that such points are not always the local minima, it is enough to look at function $E(u) = u^3$. The optimality condition $E'(u) = 3u^2 = 0$ suggests that 0 should be a local minimum, even though the function is decreasing for any $u < 0$ and can thus not have a minimum at 0. The point 0 is in fact a stationary point, not a maximum or minimum.

To discern between local minima and stationary points of a function, let us introduce the second-order approximation. Let function $E(u)$ be twice differentiable with Hessian $\nabla^2 E(u)$ at u . Then

$$E(v) = E(u) + \langle \nabla E(u), v - u \rangle + \frac{1}{2} \langle \nabla^2 E(u)(v - u), v - u \rangle + o(\|v - u\|^2).$$

The function $E(v; u) = E(u) + \langle \nabla E(u), v - u \rangle + \frac{1}{2} \langle \nabla^2 E(u)(v - u), v - u \rangle$ is the quadratic (or second-order) approximation of function E at u . Note that the Hessian is a symmetric matrix that can be seen as a derivative of the vector function ∇E . As a result, using a linear approximation to each component of ∇E , we have

$$\nabla E(v) = \nabla E(u) + \nabla^2 E(u)(v - u) + o(\|v - u\|).$$

Using the second-order approximation, we can write down the second-order optimality conditions.

Theorem 2. Second-order Pptimality Condition *Let u^* be a local minimum of twice differentiahte function $E(u)$. Then $\nabla E(u^*) = 0$ and $\nabla^2 E(u^*)$ is symmetric and positive semi-definite, that we denote by $\nabla^2 E(u^*) \succeq 0$.*

Proof. Since u^* is a local minimum of function E , there exists $r > 0$ such that

$$E(u) \geq E(u^*) \quad \forall u \quad \text{with} \quad \|u - u^*\| < r.$$

The first order optimality condition gives us $\nabla E(u^*) = 0$ and, as a result

$$E(u) = E(u^*) + \langle \nabla^2 E(u^*)(v - u^*), v - u^* \rangle + o(\|y - u^*\|^2) \geq E(u^*).$$

Thus, $\langle \nabla^2 E(u^*)(v - u^*), v - u^* \rangle \geq 0$. Letting $s = v - u$ we have $\langle \nabla^2 E(u^*)s, s \rangle \geq 0$, which implies positive semi-definiteness. \square

This second-order characteristic of a local minimum is also sufficient.

Theorem 3. *Let function $E(u)$ be twice differentiable on \mathbb{R}^n and let u^* satisfy $\nabla E(u^*) = 0$ and $\nabla^2 E(u^*) \succ 0$. Then u^* is a strict local minimum of E .*

Proof. In a small neighborhood of u^* , $E(u)$ can be represented as

$$E(u) = E(u^*) + \langle \nabla^2 E(u^*)(u - u^*), u - u^* \rangle + o(\|u - u^*\|^2).$$

Since $\lim_{r \rightarrow 0} \frac{o(r)}{r} = 0$, there exists a value \bar{r} such that for all $r \in [0, \bar{r}]$ we have

$$|o(r)| \leq \frac{r}{4} \lambda_1,$$

where $\lambda_1 > 0$ is the smallest eigenvalue of matrix $\nabla^2 E(u^*)$. As $\nabla^2 E(u^*)$ is symmetric and positive definite, it has positive eigenvalues $\lambda_1, \lambda_2, \dots, \lambda_n > 0$ and orthogonal eigenvectors q_1, q_2, \dots, q_n , such that $\nabla^2 E(u^*) = \sum_{1 \leq i \leq n} \lambda_i q_i^T q_i$ and $\|q_i^T v\| = \|v\|$ for all $v \in \mathbb{R}^n$. As a result,

$$E(u) \geq E(u^*) + \frac{\lambda_1}{2} \|u - u^*\|^2 + o(\|u - u^*\|^2) \geq E(u^*) + \frac{\lambda_1}{4} \|u - u^*\|^2 \geq E(u^*). \quad (1.2)$$

\square

For general optimization problems, we thus require second-order differentiability to formulate necessary and sufficient optimality conditions. The optima described by these conditions is, moreover, only local. This is quite disappointing because most applications in computer vision and machine learning have objective functions that are not differentiable, where these general optimality conditions are meaningless. Even in the rare cases where second-order derivatives exists, computing the Hessian is not feasible because the size of the problem is too large. For these reasons, we resort to the field of convex optimization. Convex optimization is a fairyland where the objective function does not need to be differentiable, optimality conditions are not only necessary but sufficient, and the algorithms scale well with the size of the problem.

Chapter 2

Convex Analysis

2.1 Convex Optimization

We start this section with the unconstrained minimization problem

$$\min_{u \in \mathbb{R}^n} E(u). \tag{2.1}$$

In the general situation we cannot do too much: even when the function is smooth, the gradient method converges only to a stationary point of function E and second-order differentiability is necessary to derive optimality conditions for a local minimum that are necessary and sufficient. To make the problem tractable we introduce a key assumption on the kind of functions E that we minimize. In particular, we call for the following property: for any E differentiable, the first-order optimality condition should be necessary and **sufficient** for a point to be a **global solution** of (2.1). Convex functions come with this guarantee.

Definition A function $E: \mathbb{R}^n \rightarrow \mathbf{R}$ is convex if and only if for any $u, v \in \mathbb{R}^n$ and $\theta \in [0, 1]$

$$E(\theta u + (1 - \theta)v) \leq \theta E(u) + (1 - \theta)E(v).$$

E is strictly convex if the inequality is strict for all $\theta \in (0, 1)$, $v \neq u$.

The definition of convex functions implicitly assumes that it is possible to evaluate the function at any point of the segment

$$[u, v] = \{z = \theta u + (1 - \theta)v : 0 \leq \theta \leq 1\}.$$

As a result, it is natural to consider a set that contains the whole segment between any two points in the set. Such sets are called convex.

Definition Convex Sets. The set C is convex if for any $u, v \in C$ and $\theta \in [0, 1]$, $\theta u + (1 - \theta)v \in C$.

We can then include this notion in the definition of convex functions with restricted domain.

Definition The **domain** of a function $E: \mathbb{R}^n \rightarrow \mathbb{R}$ is the set

$$\text{dom}(E) = \{u \in \mathbb{R}^n : E(u) < \infty\}$$

We can now extend the definition of convexity to functions.

Definition Convex Function. The function $E: \mathbb{R}^n \rightarrow \overline{\mathbb{R}} = \mathbb{R} \cup \{\infty\}$ is convex if

- its domain $\text{dom}(E)$ is a convex set.
- For all $u, v \in \text{dom}(E)$ and all $\theta \in [0, 1]$ it holds that

$$E(\theta u + (1 - \theta)v) \leq \theta E(u) + (1 - \theta)E(v).$$

E is **strictly convex** if the inequality is strict for all $\theta \in (0, 1)$, $v \neq u$.

In the following we assume that the domain of E is not empty, that is, the function E is proper.

Definition Function $E : \mathbb{R}^n \rightarrow \bar{\mathbb{R}}$ is **proper** if its domain is not empty.

This course will investigate convex minimization problems, they are characterized by the form

$$\hat{u} \in \arg \min_{u \in C} E(u), \tag{2.2}$$

where C is a convex set and E is a convex function. To write such a problem in our familiar unconstrained optimization form, we define the **extended real-valued function** \tilde{E} by introducing the constraint $u \in C$ into the domain of the original energy function E :

$$\tilde{E} : \mathbb{R}^n \rightarrow \bar{\mathbb{R}} := \mathbb{R} \cup \{\infty\} \quad \tilde{E}(u) = \begin{cases} E(u) & \text{if } u \in C, \\ \infty & \text{else.} \end{cases}$$

We can then re-write (2.2) as

$$\hat{u} \in \arg \min_{u \in \mathbb{R}^n} \tilde{E}(u).$$

2.2 Convex Sets

We have already seen some convex sets as a result of convex functions

Lemma 4. *If E is a convex function, then for any $\beta \in \mathbb{R}$, its level set $\{u : E(u) \leq \beta\}$ is either convex or empty.*

Proof. Let $u, v \in \text{dom}(E)$ with $E(u) \leq \beta$ and $E(v) \leq \beta$, by convexity of E we have $\theta u + (1 - \theta)v \in \text{dom}(E)$ and

$$E(\theta u + (1 - \theta)v) \leq \theta E(u) + (1 - \theta)E(v) \leq \theta\beta + (1 - \theta)\beta = \beta.$$

□

Lemma 5. *Let E be a convex function, then its **epigraph** $\text{epi}(E) = \{(u, \beta) : E(u) \leq \beta\}$ is a convex set.*

Proof. Let $(u, \alpha), (v, \beta) \in \text{epi}(E)$, then $u, v \in \text{dom}(E)$ with $E(u) \leq \alpha$ and $E(v) \leq \beta$, by convexity of E we have $\theta u + (1 - \theta)v \in \text{dom}(E)$ and $\theta(u, \alpha) + (1 - \theta)(v, \beta) \in \text{epi}(E)$ because

$$E(\theta u + (1 - \theta)v) \leq \theta E(u) + (1 - \theta)E(v) \leq \theta\alpha + (1 - \theta)\beta.$$

□

To determine if a set is convex, a few properties are useful.

Lemma 6. *Let $C \subset \mathbb{R}^n, D \subset \mathbb{R}^m$ be convex sets and $\mathcal{A} : \mathbb{R}^n \rightarrow \mathbb{R}^m$ be a linear operator, then the following sets are convex*

- Intersection $C \cap D$.
- Sum $C + D = \{u = x + y : x \in C, y \in D\}$ if $n = m$.
- Affine image $\mathcal{A}(C) = \{u \in \mathbb{R}^m : u = \mathcal{A}(x), x \in C\}$
- Inverse affine image $\mathcal{A}^{-1}(D) = \{v \in \mathbb{R}^n : \mathcal{A}(v) \in D\}$

Proof. Left as exercise □

As a result of the previous lemma, the following sets are convex

- Half-space $\{u \in \mathbb{R}^n : \langle a, u \rangle \leq \beta\}$ is convex since linear functions are convex.
- Polytope $\{u \in \mathbb{R}^n : \langle a_i, u \rangle \leq b_i\}$ is convex as an intersection of convex sets.
- Ellipsoid $\{u \in \mathbb{R}^n : \langle Au, u \rangle \leq 1 \text{ with } A \succeq 0\}$ because the function $\langle Au, u \rangle$ is a convex function.

2.3 Convex Functions

In order to determine if a function is convex, it is useful to know some equivalent definitions of convexity.

Theorem 7. Convexity and Epigraphs. *A proper function $E : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$ is convex if and only if its epigraph is convex.*

Proof. We have already seen one direction, the other is an exercise. □

Lemma 8. Jensen's Inequality. *For any convex function E , $u_1, \dots, u_m \in \text{dom}(E)$ and coefficients $\theta_1, \dots, \theta_m \geq 0$ such that $\sum_{i=1}^m \theta_i u_i = 1$ it holds*

$$E\left(\sum_{i=1}^m \theta_i u_i\right) \leq \sum_{i=1}^m \theta_i E(u_i)$$

Proof. By induction on m . The case $m = 2$ is a result of the definition and the general an exercise. □

Corollary 9. *For any u a convex combination of $u_1, \dots, u_m \in \text{dom}(E)$, $E(u) \leq \max_{1 \leq i \leq m} E(u_i)$.*

Corollary 10. *Let $\Delta = \text{Conv}\{u_1, \dots, u_m\}$ be the convex hull of u_1, \dots, u_m , then*

$$\max_{u \in \Delta} E(u) = \max_{1 \leq i \leq m} E(u_i).$$

Lemma 11. *Function $E : C \rightarrow \mathbb{R}$ is convex if and only if C is convex and for all $u, v \in C$, $\beta \geq 0$ such that $u + \beta(u - v) \in C$ it holds that*

$$E(u + \beta(u - v)) \geq E(u) + \beta(E(u) - E(v)).$$

Proof. Let E be convex, we first prove the alternative definition. Given $\beta > 0$ define $\theta = \frac{\beta}{\beta+1} \in (0, 1]$ and $x = u + \beta(u - v)$ such that

$$u = \frac{1}{1+\beta}(x + \beta v) = (1 - \theta)x + \theta v$$

by convexity of E ,

$$\begin{aligned} E(u) &\leq (1 - \theta)E(x) + \theta E(v) = \frac{1}{1+\beta}E(u + \beta(u - v)) + \frac{\beta}{1+\beta}E(v) \\ (1 + \beta)E(u) - \beta E(v) &\leq E(u + \beta(u - v)) \end{aligned}$$

Let us now prove that this alternative definition implies convexity. Given any $u, v \in \text{dom}(E)$, $\theta \in (0, 1]$, define $\beta = \frac{1-\theta}{\theta}$ and $x = \theta u + (1 - \theta)v$ such that

$$u = \frac{1}{\theta}(x - (1 - \theta)v) = x + \beta(x - v)$$

the inequality reads

$$\begin{aligned} E(u) &= E(x + \beta(x - v)) \geq E(x) + \beta[E(x) - E(v)] \\ E(u) &\geq (1 + \beta)E(x) - \beta E(v) = \frac{1}{\theta}E(x) - \frac{1 - \theta}{\theta}E(v) \\ \theta E(u) + (1 - \theta)E(v) &\geq E(\theta u + (1 - \theta)v) \end{aligned}$$

□

Theorem 12. Monotonicity of the gradient *Let $E : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$ be proper and continuously differentiable, then E is convex if and only if for any $u, v \in \text{dom}(E)$*

$$E(v) \geq E(u) + \langle \nabla E(u), v - u \rangle.$$

Proof. Given $u, v \in \text{dom}(E)$, and $\theta \in [0, 1]$, let $u_\theta = \theta u + (1 - \theta)v$. If E is continuously differentiable and satisfies the theorem's inequality, we have

$$\begin{aligned} E(u_\theta) &\geq E(v) + \langle \nabla E(u_\theta), v - u_\theta \rangle = E(v) + \theta \langle \nabla E(u_\theta), v - u \rangle \\ E(u_\theta) &\geq E(u) + \langle \nabla E(u_\theta), u - u_\theta \rangle = E(u) - (1 - \theta) \langle \nabla E(u_\theta), v - u \rangle. \end{aligned}$$

Multiplying the first inequality by $1 - \theta$, the second by θ , and adding the results, we get the inequality that defines a convex function $E(\theta u + (1 - \theta)v) \leq \theta E(u) + (1 - \theta)E(v)$.

We now prove that a convex and continuously differentiable function satisfies the theorem's inequality. Given $u, v \in \text{dom}(E)$, as E is convex for any $\theta \in [0, 1]$

$$E(v) \geq \frac{1}{1 - \theta}[E(u_\theta) - \theta E(u)] = E(u) + \frac{1}{1 - \theta}[E(u_\theta) - E(u)] = E(u) + \frac{1}{1 - \theta}[E(\theta u + (1 - \theta)v) - \theta E(u)]. \quad (2.3)$$

As E is differentiable, the limit when θ tends to 1 exists and we get $E(v) \geq E(u) + \langle \nabla E(u), v - u \rangle$. □

2.3.1 Necessary and Sufficient Optimality Conditions

Theorem 13. *Let $E : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$ be convex. Any local minimum of E is global.*

Proof. Let u^* be a global minimum of E and \bar{u} a local minimum that is not global, that is, $E(u^*) < E(\bar{u})$. By definition of local minimum, there exists an $\epsilon > 0$ such that $E(v) \geq E(\bar{u})$ for any $v \in \text{dom}(E)$ with $\|\bar{u} - v\| < \epsilon$. As $u^*, \bar{u} \in \text{dom}(E)$ convex, $\theta \bar{u} + (1 - \theta)u^* \in \text{dom}(E)$ and

$$E(\theta \bar{u} + (1 - \theta)u^*) \leq \theta E(\bar{u}) + (1 - \theta)E(u^*) < E(\bar{u})$$

As θ tends to 1, $\|\theta \bar{u} + (1 - \theta)u^* - \bar{u}\| < \epsilon$ and this contradicts the definition of \bar{u} as local minimum. □

When the function is differentiable, we can now prove that first-order optimality conditions are sufficient.

Theorem 14. *If $E : \mathbb{R}^n \rightarrow \mathbb{R}$ is continuously differentiable function with $\nabla E(u^*) = 0$ then u^* is the global minimum of $E(x)$.*

Proof. As $\nabla E(u^*) = 0$, the inequality $E(v) \geq E(u^*) + \langle \nabla E(u^*), v - u^* \rangle \forall v \in \text{dom}(E)$ gives us the condition $E(v) \geq E(u^*)$ that characterizes a global minimum. \square

When the objective function is two-times differentiable, we can also characterize convexity in terms of the Hessian.

Theorem 15. *Two times continuously differentiable function $E: \mathbb{R}^n \rightarrow \mathbb{R}$ is convex if and only if for any $u \in \mathbb{R}^n$ we have $\nabla^2 E(u) \succeq 0$.*

Proof. This is part of an exercise sheet. \square

As a result, for any matrix A symmetric and positive semi-definite, the quadratic function $E(u) = \alpha + \langle a, u \rangle + \langle a, Au \rangle$ is convex because $\nabla^2 E(u) = A \succeq 0$.

2.3.2 Analytic Properties of Convex Functions

The behavior of convex functions at the boundary of their domain can be out of control. To prevent this case, we ask the functions to be closed.

Definition Closed convex function. A convex function is closed if its epigraph is closed.

Lemma 16. *If E is convex and closed, all its level sets are closed.*

Proof. For each β , the level-set $\{u : E(u) = \beta\} = \text{epi}(E) \cap \{(x, t) : t = \beta\}$ can be described as the intersection of the epigraph of E , which is closed and convex, and the closed and convex set $\{(x, t) : t = \beta\}$. \square

If E is convex and continuous and its domain $\text{dom}(E)$ is closed, then E is closed. The converse is not true, a closed convex function is not necessarily continuous. Consider the following examples:

- $E(u) = \frac{1}{u}$ is convex, has an open domain $\text{dom}(E) = \mathbb{R}_{++} = \{u \in \mathbb{R} : u > 0\}$, but is closed because its epigraph $\{(u, t) \in \mathbb{R} \times \mathbb{R}_{++} : \frac{1}{t} \leq u\}$ is closed.
- Function $E(u) = \|u\|$, where $\|\cdot\|$ is any norm, is closed and convex as a result of the triangle inequality and homogeneity properties that define any norm:

$$\|\theta u + (1 - \theta)v\| \leq \|\theta u\| + \|(1 - \theta)v\| = |\theta|\|u\| + |1 - \theta|\|v\| = \theta\|u\| + (1 - \theta)\|v\|$$

The norms more common in computer vision and machine learning are the ℓ_p norms:

$$\|u\| = \left(\sum_{i=1}^n |u_i|^p \right)^{\frac{1}{p}} \quad u \in \mathbb{R}^n.$$

- the Euclidean norm: $|u| = \sqrt{\sum_{i=1}^n u_i^2}$.
- the non-differentiable ℓ_1 norm $\|u\|_1 = \sum_{i=1}^n |u_i|$.
- the ℓ_∞ norm $\|u\|_\infty = \max_{1 \leq i \leq n} |u_i|$.

Any norm defines a system of balls $B_p(u, r) = \{v \in \mathbf{R}^n : \|v - u\|_p \leq r\}$ that are convex.

- the function

$$E(x, y) = \begin{cases} 0 & \text{if } x^2 + y^2 < 1 \\ \phi(x, y) & \text{if } x^2 + y^2 = 1 \end{cases}$$

with domain the unit ball is closed and convex for any $\phi(x, y) > 0$ defined on the unit circle, the boundary of the function domain. Imposing that the function is closed, which implies $\phi(x, y) = 0$, ensures that the function is well-behaved also on the boundary of its domain.

The behavior of convex function at the boundary of their domain can be disappointing, but their behavior in the interior of its domain is very simple.

Theorem 17. *Let function $E: C \subset \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$ be convex, then E is locally bounded at $u \in \text{int dom}(E)$.*

Proof. Let us choose $\epsilon > 0$ such that $u \pm \epsilon e_i \in \text{int dom}(E)$ $i = 1, \dots, n$, where e_i is the i -th coordinate vector of \mathbb{R}^n and define $\hat{\epsilon} = \frac{\epsilon}{\sqrt{n}}$. A simple drawing show us that

$$B(u, \hat{\epsilon}) \subset \Delta = \text{Conv}\{u \pm \epsilon e_i \quad i = 1, \dots, n\}.$$

From the corollary to Jensen's inequality, we find a local bound M to E

$$M = \max_{v \in B(u, \hat{\epsilon})} E(v) \leq \max_{v \in \Delta} E(v) \leq \max_{1 \leq i \leq n} E(u \pm \epsilon e_i)$$

□

Theorem 18. Continuity of Convex Functions *If $E: \mathbb{R}^n \rightarrow \mathbb{R} \cup \{\infty\}$ is convex, then E is locally Lipschitz and hence continuous on $\text{int}(\text{dom}(E))$.*

Let us first define Lipschitz continuity.

Definition A function $E: \mathbb{R}^n \rightarrow \mathbb{R}^m$ is **Lipschitz continuous** with *Lipschitz constant* L if for all $u, v \in \text{dom}(E)$

$$\|E(u) - E(v)\|_2 \leq L\|u - v\|_2$$

A function is **locally Lipschitz continuous** if for every $u \in \text{dom}(E)$ there exists $\epsilon > 0$ such that $f|_{B(\epsilon, u)}$ is Lipschitz continuous

Proof. Let $B(u_0, \epsilon) \subset \text{dom}(E)$ and $M = \sup_{u \in B(u_0, \epsilon)} E(u) < \infty$.

Consider $v \in B(u_0, \epsilon)$, $v \neq u_0$ and define

$$\alpha = \frac{1}{\epsilon}\|v - u_0\| \qquad z = u_0 + \frac{1}{\alpha}(v - u_0)$$

It is clear that $\|z - u_0\| = \epsilon$, $\alpha \leq 1$, and $v = \alpha z + (1 - \alpha)u_0$. By convexity of E then

$$E(v) \leq \alpha E(z) + (1 - \alpha)E(u_0) \leq E(u_0) + \alpha(M - E(u_0)) = E(u_0) + \frac{M - E(u_0)}{\epsilon}\|v - u_0\|$$

Now define $y = u_0 + \frac{1}{\alpha}(u_0 - v)$ with $\|y - u_0\| = \epsilon$ and $v = u_0 + \alpha(u_0 - y)$. We have

$$E(v) \geq E(u_0) + \alpha(E(u_0) - E(y)) \geq E(u_0) - \alpha(M - E(u_0)) = E(u_0) - \frac{M - E(u_0)}{\epsilon}\|v - u_0\|$$

As a result of the 2 inequalities

$$|E(v) - E(u)| \leq \frac{M - E(u_0)}{\epsilon}\|v - u_0\|.$$

□

2.3.3 Examples of Convex Functions

The next statements significantly increases our possibilities of constructing convex functions.

Lemma 19. *Given a closed convex function ϕ and a linear operator $\mathcal{A}: \mathbb{R}^m \rightarrow \mathbb{R}^n$, then $E(u) = \phi(\mathcal{A}(u))$ is closed and convex with*

$$\text{dom}(E) = \{u \in \mathbb{R}^m : \mathcal{A}(u) \in \text{dom}(\phi)\}.$$

Proof. Let $\mathcal{A}(u) = Au + b = x \in \text{dom}(\phi)$ and $\mathcal{A}(v) = Av + b = y \in \text{dom}(\phi)$, then by convexity of ϕ for any $\theta \in [0, 1]$ we have $\theta x + (1 - \theta)y \in \text{dom}(\phi)$ and

$$E[\theta u + (1 - \theta)v] = \phi[\mathcal{A}(\theta u + (1 - \theta)v)] = \phi[\theta(Au + b) + (1 - \theta)(Av + b)] \leq \theta\phi(Au + b) + (1 - \theta)\phi(Av + b) = \theta E(u) + (1 - \theta)E(v).$$

This proves convexity of E . The closedness of its epigraph follows from continuity of the linear operator \mathcal{A} . \square

Lemma 20. *Given two convex function E_1, E_2 and $\alpha_1, \alpha_2 > 0$, then $E = \alpha_1 E_1 + \alpha_2 E_2$ is convex with $\text{dom}(E) = \text{dom}(E_1) \cap \text{dom}E_2$.*

Proof. Let $u, v \in \text{dom}(E_1) \cap \text{dom}E_2$ and $\theta \in [0, 1]$, by convexity of each E_1, E_2 we have

$$\begin{aligned} \alpha_1 E_1(\theta u + (1 - \theta)v) + \alpha_2 E_2(\theta u + (1 - \theta)v) &\leq \alpha_1 \theta E_1(u) + \alpha_1 (1 - \theta)E_1(v) + \alpha_2 \theta E_2(u) + \alpha_2 (1 - \theta)E_2(v) \\ &= \theta[\alpha_1 E_1(u) + \alpha_2 E_2(u)] + (1 - \theta)[\alpha_1 E_1(v) + \alpha_2 E_2(v)]. \end{aligned} \quad (2.4)$$

This proves the convexity of E . \square

Taking into account that the following 1-dimensional functions are convex:

$$\begin{aligned} E(u) &= \exp(u) \\ E(u) &= |u|^p \quad p > 1 \\ E(u) &= |x| - \log(1 + |x|) \end{aligned}$$

the previous lemmas imply that the following multi-dimensional functions are convex:

$$\begin{aligned} E(u) &= \sum_{i=1}^n \exp(\alpha + \langle u, a_i \rangle) \\ E(u) &= |\langle u, a_i \rangle - b_i|^p \quad p > 1 \end{aligned}$$

Lemma 21. *Given two closed and convex function E_1, E_2 , then $E(u) = \max\{E_1(u), E_2(u)\}$ is closed and convex with $\text{dom}(E) = \text{dom}(E_1) \cup \text{dom}(E_2)$.*

Proof. The epigraph is closed and convex because it is the intersection of two closed convex sets

$$\text{epi}(E) = \{(u, t) : u \in \text{dom}(E_1) \cup \text{dom}(E_2), E_1(u) \leq t, E_2(u) \leq t\} = \text{epi}(E_1) \cap \text{epi}(E_2).$$

\square

We have an even more general result.

Theorem 22. *Let D be some set, not necessarily convex or finite dimensional, and*

$$E(u) = \sup_{y \in C} \phi(u, y) \quad \phi \text{ closed and convex in } u \quad \forall y \in D,$$

then E is closed and convex with $\text{dom}(E) = \{u \in \cap_{y \in D} \text{dom}(\phi(\cdot, y)) : \exists \gamma \in \mathbb{R} \text{ s.t. } \phi(u, y) \leq \gamma \forall y \in D\}$.

Proof. We first show the definition of the domain. If u belongs to $\{u \in \cap_{y \in D} \text{dom}(\phi(\cdot, y)) : \exists \gamma \in \mathbb{R} \text{ s.t. } \phi(u, y) \leq \gamma \forall y \in D\}$, then $E(u) < \infty$ and $u \in \text{dom}(E)$. If u does not belong to this set, then there exists a sequence $\{y_k\}$ such that $\phi(u, y_k) \rightarrow \infty$ and u does not belong to $\text{dom}(E)$.

$(u, t) \in \text{epi}(E)$ if and only if for all $y \in D$, we have $u \in \text{dom}(\phi(\cdot, y))$ and $\phi(u, y) \leq t$. As a results $\text{epi}(E) = \cap_{y \in D} \text{epi}(\phi(\cdot, y))$ is closed and convex as the intersection of closed and convex sets. \square

As a result of this lemma, the function $E^*(y) = \sup_{u \in \text{dom}(E)} \langle u, y \rangle - E(u)$ is convex for any E .

2.4 Existence and Uniqueness of Minimizers

It only makes sense to try to solve an optimization problem if it has a solution. Specially, if the solution is the limit of a relaxation sequence that is computed through costly iterative algorithm that might never converge. To show that a convex problem has a minimizer, we will see that it satisfies the necessary conditions to frame the problem in the more general framework of lower semi-continuous functions. This section explains the tools that we will use from this framework.

Definition Lower semi-continuity. A function $E : \mathbb{R}^n \rightarrow \mathbb{R}$ is lower semi-continuous (l.s.c.), if for all u it holds that

$$\liminf_{v \rightarrow u} E(v) \geq E(u).$$

Theorem 23. Let $E : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$ be l.s.c. and let there exist an α such that the sublevelset

$$\{u \in \mathbb{R}^n \mid E(u) \leq \alpha\}$$

is nonempty and bounded, then there exists

$$\hat{u} \in \arg \min_u E(u).$$

Proof. Remember that the infimum is the largest lower bound on all possible values of $E(u)$ and consider a sequence $(u_k)_k$ such that $E(u_k) \rightarrow \inf_u E(u)$.

We distinguish two cases: For $\alpha = \inf_u E(u)$ the non-emptiness of S_α yields the assertion. For $\alpha > \inf_u E(u)$ it holds that from some sufficiently large k_0 on, we will have $u_k \in S_\alpha$. Since S_α is bounded there exists a convergent subsequence $u_{k_l} \rightarrow \bar{u}$. Due to the lower semi-continuity we find

$$\inf_u E(u) = \lim_{k \rightarrow \infty} E(u_k) = \lim_{l \rightarrow \infty} E(u_{k_l}) \geq E(\bar{u}).$$

Since by definition $\inf_u E(u) \leq E(\bar{u})$ we obtain equality and hence there exists $\bar{u} \in \operatorname{argmin}_u E(u)$. \square

Theorem 24. Equivalence of l.s.c. and closedness. For $E : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$ the following two statements are equivalent

- E is lower semi-continuous (l.s.c.).
- E is closed (its epigraph is closed).

Proof. Let E be closed and assume that E is not l.s.c. Then there exists a point u^0 and a sequence $(u_k)_k$ with $\lim_k u_k = u^0$ such that

$$\liminf_k E(u_k) < E(u^0).$$

In particular, there exists $\alpha \in \mathbb{R}$ and a subsequence $(u_{k_l})_{k_l}$ such that

$$E(u_{k_l}) \leq \alpha < E(u^0) \quad \forall k \tag{2.5}$$

Obviously, $(u_{k_l}, \alpha) \in \operatorname{epi}(E)$ for all k_l and $(u_{k_l}, \alpha) \rightarrow (u^0, \alpha)$, but according to (2.5) $(u^0, \alpha) \notin \operatorname{epi}(E)$, which contradicts the closedness of E .

To prove the other direction of the claim, let E be l.s.c. and assume that E is not closed. Then there exists a sequence $(u_k, \alpha_k) \in \operatorname{epi}(E)$ with $(u_k, \alpha_k) \rightarrow (u^0, \alpha^0) \notin \operatorname{epi}(E)$. We find

$$\liminf_k E(u_k) \leq \lim_k \alpha_k = \alpha^0 < E(u^0).$$

On the other hand, due to E being l.s.c. we have $E(u^0) \leq \liminf_k E(u_k)$, which is a contradiction. \square

2.4.1 Existence of Minimizers of Convex Functions

Definition Coercivity. A function $E : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{\infty\}$ is called coercive if $E(v_n) \rightarrow \infty$ for all sequences $(v_n)_n$ with $\|v_n\| \rightarrow \infty$.

It is easy to prove that coercivity implies existence of a bounded sublevelset by contradiction. We have now all the tools to prove existence of minimizers of convex functions.

Theorem 25. Existence of a Minimizer *Let $E : \mathbb{R}^n \rightarrow \mathbb{R}$ be convex and coercive, then an element $\hat{u} \in \arg \min_u E(u)$ exists.*

Proof. As $\text{dom}(E) = \mathbb{R}^n$ and E convex, E is Lipschitz continuous, and thus continuous. At the same time, as E is coercive, there exists a non-empty bounded sublevelset, and we can apply the theorem on the existence of minimizers for lower semi-continuous functions to prove existence of a minimizer. \square

Theorem 26. Uniqueness. *If $E : \mathbb{R}^n \rightarrow \bar{\mathbb{R}}$ is strictly convex, then there exists at most one local minimum which is the unique global minimum.*

Proof. Assume there are 2 global minima u, v with $u \neq v$, $E(u) = E(v)$, then any $\theta \in [0, 1]$ we have

$$E(\theta u + (1 - \theta)v) < \theta E(u) + (1 - \theta)E(v),$$

which contradicts the definition of u, v as global minima. \square

2.5 Subdifferentials

2.5.1 Supporting Hyperplanes

Up to now we were describing properties of convex functions in terms of function values or their gradients. When the function is not differentiable, we need to define a direction that acts as the gradient of differentiable functions and points onto the direction of maximum ascent. In convex analysis such directions are defined by supporting hyperplanes.

Definition Let C be a convex set. We say that hyperplane

$$\mathcal{H}(g, \gamma) = \{u \in \mathbb{R}^n : \langle g, u \rangle = \gamma, \quad g \neq 0\}$$

is supporting to C if any $u \in C$ satisfies $\langle g, u \rangle \leq \gamma$.

We say that the hyperplane $\mathcal{H}(g, \gamma)$ separates a point u_0 from C if

$$\langle g, u \rangle \leq \gamma \leq \langle g, u_0 \rangle \quad \forall u \in C.$$

Now we can enunciate two separation theorems necessary to define gradient-like directions for non-differentiable functions.

Theorem 27. Separating Hyperplane Theorem *Let C be a closed convex set and $u_0 \notin C$. Then there exists a hyperplane $\mathcal{H}(g, \gamma)$ that strictly separates u_0 from C .*

Proof. See Boyd and Vandenberghe, *Convex Optimization Theory*, pp 46–49. \square

The next separation theorem deals with boundary points of convex sets.

Theorem 28. Supporting Hyperplane Theorem *Let C be a closed convex set and u_0 in the boundary of C . Then there exists a hyperplane $\mathcal{H}(g, \gamma)$ supporting to C and passing through u_0 .*

Proof. See Boyd and Vandenberghe, *Convex Optimization Theory*, pp 50–51. \square

2.5.2 The Subdifferential

We now have all the tools to introduce the notion of subdifferential that extends the notion of gradient to non-differentiable functions.

Definition Subdifferential. Let $E : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$ be convex, the subdifferential of E at u is

$$\partial E(u) = \{p \in \mathbb{R}^n \mid E(v) - E(u) - \langle p, v - u \rangle \geq 0, \forall v \in \mathbb{R}^n\}$$

- Elements of $\partial E(u)$ are called subgradients.
- If $\partial E(u) \neq \emptyset$, we call E subdifferentiable at u .
- By convention, $\partial E(u) = \emptyset$ for $u \notin \text{dom}(E)$.

The subdifferential ∂E is necessary because subgradients are not unique. Consider for example a function as friendly-looking as the absolute value at zero:

$$\forall g \in [-1, 1], \quad E(u) = |u| \geq gu = E(0) + \langle g, u - 0 \rangle$$

As a result, the subdifferential at 0 contains the interval $\partial E(0) = [-1, 1]$. In general $\partial E(u)$ is a set. Form its definition as a set of linear constraints, we can easily see that it is closed and convex, in this case the interval $[-1, 1]$.

2.5.3 Subdifferentiability and Convexity

The subdifferentiability of a function is important because it implies its convexity.

Theorem 29. *If for any $u \in \text{dom}(E)$ the subdifferential $\partial E(u)$ is non-empty, then E is a convex function.*

Proof. Given $u, v \in \text{dom}(E)$, and $\theta \in [0, 1]$, let $u_\theta = \theta u + (1 - \theta)v$. As the subdifferential $\partial E(u_\theta)$ is non-empty, we can pick $g \in \partial E(u_\theta)$ satisfying

$$\begin{aligned} E(u_\theta) &\geq E(v) + \langle g, v - u_\theta \rangle = E(v) + \theta \langle g, v - u \rangle \\ E(u_\theta) &\geq E(u) + \langle g, u - u_\theta \rangle = E(u) - (1 - \theta) \langle g, v - u \rangle. \end{aligned}$$

Multiplying the first inequality by $1 - \theta$, the second by θ , and adding the results, we get the inequality that defines a convex function $E(\theta u + (1 - \theta)v) \leq \theta E(u) + (1 - \theta)E(v)$. \square

The converse statement is also true.

Theorem 30. *If E is a closed convex function and $u \in \text{int}(\text{dom}(E))$, then $\partial E(u)$ is a non-empty bounded set.*

Proof. Note that the point $(E(u), u)$ belongs to the boundary of $\text{epi}(E)$, which is convex. As a result, there exists a hyperplane $\mathcal{H} = (g, \gamma)$ supporting to $\text{epi}(E)$ at $(E(u), u)$:

$$\gamma \tau + \langle g, u \rangle \leq \gamma E(u) + \langle g, u \rangle \quad \forall (u, \tau) \in \text{epi}(E)$$

Without loss of generality, we can assume $\|g\|^2 + \gamma^2 = 1$. We can determine the sign of γ by checking the inequality for any point in the epigraph. In particular for any $\tau \geq E(u)$, we have $(u, \tau) \in \text{epi}(E)$ that results in $\gamma > 0$.

To find a subgradient $p \in \partial E(u)$, we will use that a convex function is locally upper bounded in the interior of its domain. That is, there is some $\epsilon > 0, M > 0$ such that $B(u, \epsilon) \subset \text{dom}(E)$ and

$$E(v) - E(u) \leq M\|v - u\| \quad \forall v \in B(u, \epsilon)$$

For any v from this ball, the supporting hyperplane equation reads

$$\langle g, v - u \rangle \leq \gamma(E(v) - E(u)) \leq \gamma M \|v - u\|$$

In particular, if we choose $v = u + \epsilon g$ we get $\|g\|^2 \leq M\gamma\|d\|$. Plugging now the condition $\|g\|^2 + \gamma^2 = 1$ we get

$$\gamma \geq \frac{1}{\sqrt{1 + M^2}}.$$

If we choose $p = \frac{g}{\gamma}$ we obtain

$$E(v) \geq E(u) + \langle p, v - u \rangle \quad \forall v \in \text{dom}(E)$$

and p is a subgradient of E at u . Finally, to show that the subdifferential is bounded we assume that $p \neq 0$ and consider the point $v = u + \epsilon \frac{p}{\|p\|}$ such that

$$\epsilon \|p\| = \langle p, v - u \rangle \leq E(v) - E(u) \leq M \|v - u\| = M\epsilon$$

Thus, $\partial E(u)$ is bounded by M . □

The conditions of this theorem cannot be relaxed. For instance, the function $E(u) = -\sqrt{u}$ is convex and closed in its domain $\{u: u \geq 0\}$, but its subdifferential does not exist at the only point (0) that is not in its interior. This is just another reminder that considering the interior of the domain for convex functions is important.

To conclude this section, let us point out to the property of the subgradients that makes it important for optimization.

Theorem 31. Optimality Condition. $0 \in \partial E(\hat{u})$ if and only if $\hat{u} \in \arg \min_{u \in \mathbb{R}^n} E(u)$.

Proof. If $0 \in \partial E(\hat{u})$, by definition of the subgradient

$$E(u) \geq E(\hat{u}) + \langle 0, u - \hat{u} \rangle = E(\hat{u}) \quad \forall u \in \text{dom}(E)$$

and we conclude that \hat{u} is a minimizer of E . On the other hand, if $E(u) \geq E(\hat{u})$ for all $u \in \text{dom}(E)$, then 0 satisfies the condition of subgradient of E at \hat{u} . □

2.5.4 Alternative Definitions of Subgradients

The supporting hyperplane theorem appears on the proof of the “subdifferentiability” theorem because subgradients can be interpreted in terms of supporting hyperplanes.

Theorem 32. Geometric interpretation of Subgradients. Any subgradient $p \in \partial E(u)$ represents a non-vertical supporting hyperplane to $\text{epi}(E)$ at $(u, E(u))$.

Proof. Let $p \in \partial E(u)$. Then, by definition of subgradient,

$$\begin{aligned} E(v) - E(u) - \langle p, v - u \rangle &\geq 0 && \forall v \in \mathbb{R}^n \\ \alpha - E(u) - \langle p, v - u \rangle &\geq 0 && \forall (v, \alpha) \in \text{epi}(E) \\ \left\langle \begin{bmatrix} -p \\ 1 \end{bmatrix}, \begin{bmatrix} v \\ \alpha \end{bmatrix} - \begin{bmatrix} u \\ E(u) \end{bmatrix} \right\rangle &\geq 0 && \forall (v, \alpha) \in \text{epi}(E). \end{aligned}$$

As a result, the non-vertical hyperplane $\mathcal{H} = (g, \gamma)$ with $g = (-p, 1)$ and $\gamma = \langle p, u \rangle - E(u)$ supports $\text{epi}(E)$ at $(u, E(u))$. □

Apart from this geometric interpretation, it is useful to compute the subdifferential of a differentiable function to understand why it is a generalization of the gradient. The next theorem does that.

Theorem 33. Subdifferential of Differentiable Functions. *Let the convex function $E : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{\infty\}$ be differentiable at $u \in \text{int}(\text{dom}(E))$. Then*

$$\partial E(u) = \{\nabla E(u)\}.$$

Proof. The subdifferential $\partial E(u)$ of some convex E at $u \in \text{dom}(f)$ is given as

$$\{p \in \mathbb{R}^n : E(z) - E(u) - \langle p, z - u \rangle \geq 0, \forall z \in \text{dom}(f)\}.$$

Since $u \in \text{int}(\text{dom}(E))$, we find that for all $v \in \mathbb{R}^n$, $z = u \pm \epsilon v \in \text{dom}(E)$ for ϵ small enough. Therefore, it holds that

$$E(u + \epsilon v) \geq E(u) + \epsilon \langle p, v \rangle, \quad E(u - \epsilon v) \geq E(u) - \epsilon \langle p, v \rangle,$$

for all $v \in \mathbb{R}^n$ and ϵ small enough. This implies that

$$\lim_{\epsilon \rightarrow 0} \frac{E(u + \epsilon v) - E(u)}{\epsilon} \geq \langle p, v \rangle, \quad \lim_{\epsilon \rightarrow 0} \frac{E(u) - E(u - \epsilon v)}{\epsilon} \leq \langle p, v \rangle,$$

which means

$$\langle \nabla E(u), v \rangle \geq \langle p, v \rangle, \quad \langle \nabla E(u), v \rangle \leq \langle p, v \rangle,$$

i.e.

$$\langle \nabla E(u) - p, v \rangle = 0$$

for all $v \in \mathbb{R}^n$. For the particular choice of $v := \nabla E(u) - p$ we find $p = \nabla f(u)$. The above concludes the proof if we can show that $\partial f(u)$ is non-empty, which follows from the Theorem on Subdifferentiability. \square

2.5.5 Subdifferential Rules

In the same way that the gradient of a differentiable function is only defined for points in the interior of the domain, the subdifferential of a proper convex function is always defined for points in the relative interior of its domain.

The relative interior of a set is a refinement of the concept of the interior that is useful when dealing with low-dimensional sets embedded in higher-dimensional spaces. Intuitively, the relative interior of a set contains all points that are not on the “edge” of the set, relative to the smallest subspace in which this set lies. When the set is convex, the definition takes the following simple form:

Definition Relative Interior of Convex Sets The relative interior of a convex set C is defined as

$$\text{ri}(C) := \{x \in C \mid \forall y \in C, \exists \lambda > 1, \text{ s.t. } \lambda x + (1 - \lambda)y \in C\}$$

As mentioned earlier, the subdifferentiability of convex functions can be guaranteed for points that are not necessarily in the interior of the domain, but that are in its relative interior. To better understand this difference, consider the line segment $I = [-1, 1]$ as a convex subset of the Euclidean plane $I \subset \mathbb{R}^2$. The interior of I is empty with the Euclidean topology of \mathbb{R}^2 , but its relative interior is the open line segment $\text{ri}(I) = (0, 1)$.

One key property of the relative interior is that it is not empty for convex sets.

Theorem 34. *Let C be a non-empty convex set, then $\text{ri}(C)$ is not empty.*

Now that we understand where subdifferentials exist, we can learn the rules that guide their computation.

Theorem 35. Sum Rule. *Let E_1, E_2 be convex functions, then $\partial(E_1 + E_2)(u) = \partial E_1(u) + \partial E_2(u)$ for all $u \in \text{ri}(\text{dom}(E_1)) \cap \text{ri}(\text{dom}(E_2))$.*

Proof. See Nesterov, *Introductory Lectures on Convex Optimization*, Lemma. 3.1.9. \square

Theorem 36. Chain Rule *Given the linear operator $A \in \mathbb{R}^{m \times n}$ and the convex function $E : \mathbb{R}^m \rightarrow \mathbb{R} \cup \{\infty\}$, then $\partial(E \circ A)(u) = A^* \partial E(Au)$ for all $u \in \text{ri}(\text{dom}(E)) \cap \text{range}(A)$.*

Proof. See Nesterov, *Introductory Lectures on Convex Optimization*, Nesterov, Lemma. 3.1.8. \square

L -smooth functions

Assuming only differentiability of the objective function we cannot get many reasonable properties of minimization processes. We usually have to impose some additional assumptions on the magnitude of the derivatives. In optimization these kind of assumptions are presented in the form of a Lipschitz condition for a derivative of certain order. Among them, we will make heavy use of L -smoothness.

Definition L -smooth function A differentiable function $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is L -smooth is

$$\|\nabla f(x) - \nabla f(y)\| \leq L\|x - y\| \quad \forall x, y \in \mathbb{R}^n.$$

If a function is twice-continuously differentiable, a sufficient condition for L -smoothness is the following:

Lemma 37. *A twice-continuously differentiable function f is L -smooth if and only if $\|\nabla^2 f(x)\| \leq L \forall x \in \mathbb{R}^n$.*

Proof. Let us first prove that any twice-continuously differentiable function with bounded Hessian is smooth. Given any $x, y \in \mathbb{R}^n$, we have the componentwise inequality

$$\begin{aligned} \nabla f(y) &= \nabla f(x) + \int_0^1 \nabla^2 f(x + \tau(y-x))(y-x) d\tau \\ &= \nabla f(x) + \left(\int_0^1 \nabla^2 f(x + \tau(y-x)) d\tau \right) (y-x) \end{aligned}$$

Re arranging terms and using Cauchy-Schwarz inequality,

$$\begin{aligned} \|\nabla f(y) - \nabla f(x)\| &= \left\| \left(\int_0^1 \nabla^2 f(x + \tau(y-x)) d\tau \right) (y-x) \right\| \\ &\leq \left\| \int_0^1 \nabla^2 f(x + \tau(y-x)) d\tau \right\| \|y-x\| \\ &\leq \int_0^1 \|\nabla^2 f(x + \tau(y-x))\| d\tau \|y-x\| \leq L\|y-x\|. \end{aligned}$$

Let us now prove the other direction, that is, that a twice-continuously differentiable function that is L -smooth has bounded Hessian. As f is twice-continuously differentiable, we have

$$\left\| \left(\int_0^\alpha \nabla^2 f(x + \tau s) d\tau \right) s \right\| = \|\nabla f(x + \alpha s) - \nabla f(x)\| \leq \alpha L \|s\|$$

Dividing this inequality by α and tending $\alpha \rightarrow 0$ we obtain $\|\nabla^2 f(x)\| \leq L$. \square

The next statement is important for the geometric interpretation of L -smooth functions.

Lemma 38. *If f is L -smooth, then for any $x, y \in \mathbb{R}^n$*

$$|f(y) - f(x) - \langle \nabla f(x), y - x \rangle| \leq \frac{L}{2} \|y - x\|^2.$$

Proof. For all $x, y \in \mathbb{R}^n$, we have

$$\begin{aligned} f(y) &= f(x) + \int_0^1 \langle \nabla f(x + \tau(y-x)), y-x \rangle d\tau \\ &= f(x) + \langle \nabla f(x), y-x \rangle + \int_0^1 \langle \nabla f(x + \tau(y-x)) - \nabla f(x), y-x \rangle d\tau \end{aligned}$$

Re-arranging terms we get

$$\begin{aligned} |f(y) - f(x) - \langle \nabla f(x), y-x \rangle| &= \left| \int_0^1 \langle \nabla f(x + \tau(y-x)) - \nabla f(x), y-x \rangle d\tau \right| \\ &\leq \int_0^1 |\langle \nabla f(x + \tau(y-x)) - \nabla f(x), y-x \rangle| d\tau \\ &\leq \int_0^1 \|\nabla f(x + \tau(y-x)) - \nabla f(x)\| \|y-x\| d\tau = \int_0^1 \tau L \|x-y\|^2 d\tau = \frac{1}{2} L \|x-y\|^2 \end{aligned}$$

□

Geometrically, we can draw the following picture. Given a differentiable L -smooth function f and $x_0 \in \mathbb{R}^n$, we can define two quadratic functions

$$\begin{aligned} \phi_1(x) &= f(x_0) + \langle \nabla f(x_0), x-x_0 \rangle - u \frac{L}{2} \|x-x_0\|^2 \\ \phi_2(x) &= f(x_0) + \langle \nabla f(x_0), x-x_0 \rangle + u \frac{L}{2} \|x-x_0\|^2 \end{aligned}$$

that upper and lower bound the function

$$\phi_1(x) \leq f(x) \leq \phi_2(x) \quad \forall x \in \mathbb{R}^n.$$

Definition Strong Convexity A function $E : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$ is called *strongly convex* with constant m or m -strongly convex if $f(x) - \frac{m}{2} \|x\|_2^2$ is still convex.

Theorem 39. For a continuously differentiable function f , the following are equivalent:

1. $f(x) - \frac{m}{2} \|x\|^2$ is convex
2. $f(y) \geq f(x) + \langle \nabla f(x), y-x \rangle + \frac{m}{2} \|y-x\|^2$
3. $\langle \nabla f(x) - \nabla f(y), x-y \rangle \geq m \|y-x\|^2$
4. if f is twice-continuously differentiable, $\nabla^2 f(x) \succeq m \cdot I$

Proof. See Ryu, Boyd, *A Primer on Monotone Operator Methods*, Appendix A

□

Chapter 3

Fixed-Point Iterations

Convex optimization problems come in so many shapes and sizes that the algorithms developed to solve them form a zoo. Each algorithm exploits a particular feature of the convexity of the objective function or the constraint set to find the solution of the problem. As a result, we traditionally also analyze the convergence of each algorithm and its properties in a case by case manner.

It is possible to interpret many of these algorithms as fixed-point iterations in a unified manner and analyze their convergence with the same approach. To do so, we first need to formulate the optimization problem as finding a zero of a monotone operator. This problem is converted into the problem of finding a fixed point of a function and solved by the fixed point iteration algorithm. Different choices of the monotone operator and fixed point function result in different well-known algorithms.

This new view on many classic algorithms provides a convenient strategy to analyze their convergence with a single approach. The price to pay, however, is an additional level of abstraction that might at first seem disconnected from intuitive algorithms like gradient descent. Be patient, and read on.

The material of this chapter is taken mostly from: Ryu and Boyd, *Primer on Monotone Operator Methods*, 2016.

3.1 Nonexpansive mappings and contractions

Definition A function $F : \mathbb{R}^n \rightarrow \mathbb{R}^n$ is a **contraction** if it is Lipschitz continuous with constant $L < 1$, that is,

$$\|F(x) - F(y)\| \leq L\|x - y\| \quad \forall x, y \in \text{dom}(F).$$

When $L = 1$, we say that F is a **nonexpansive** operator.

In other words, mapping a pair of points by a contraction reduces the distance between them; mapping them by a nonexpansive operator does not increase the distance between them. See Figure 3.1.

Intuitively, it is useful to keep in mind an exemplary contraction and an exemplary nonexpansive operator. You can think of a contraction as a “zoom-out” that reduces the distance between two points, and think of a nonexpansive operator as a rotation of the coordinate plane. It is then only natural to see that the combination of two zoom-outs (contractions) is still a zoom-out, while the combination of a zoom-out and a rotation is also a zoom-out. The following lemma describes this.

Lemma 40. *Convex combinations as well as compositions of nonexpansive operators are nonexpansive.*

Proof. If $F_1 : \mathbb{R}^n \rightarrow \mathbb{R}^n$ has Lipschitz constant L_1 and $F_2 : \mathbb{R}^n \rightarrow \mathbb{R}^n$ has Lipschitz constant L_2 , then the composition F_2F_1 has Lipschitz constant L_2L_1 . Indeed, let $x, y \in \text{dom}(F_1)$ such that $F_1x, F_1y \in \text{dom}(F_2)$



(a) Contraction Mapping

(b) Nonexpansive Mapping

Fig. 3.1: Illustration of a contractive and a nonexpansive mapping F on two points. Source: Ryu and Boyd, *Primer on Monotone Operator Methods*, 2016

then

$$\|F_2F_1x - F_2F_1y\| \leq L_2\|F_1x - F_1y\| \leq L_2L_1\|x - y\|$$

As a result, the composition of nonexpansive operators is nonexpansive, and the composition of a contraction and a nonexpansive operator is a contraction.

Similarly, if $\alpha_1, \alpha_2 \in \mathbb{R}$, then $\alpha_1F_1 + \alpha_2F_2$ has Lipschitz constant $|\alpha_1|L_1 + |\alpha_2|L_2$. Indeed, let $x, y \in \text{dom}(F_1) \cap \text{dom}(F_2)$, then

$$\begin{aligned} \|(\alpha_1F_1 + \alpha_2F_2)x - (\alpha_1F_1 + \alpha_2F_2)y\| &\leq \|\alpha_1F_1x - \alpha_1F_1y\| + \|\alpha_2F_2x - \alpha_2F_2y\| \\ &\leq |\alpha_1|L_1\|x - y\| + |\alpha_2|L_2\|x - y\| \\ &\leq (|\alpha_1|L_1 + |\alpha_2|L_2)\|x - y\|. \end{aligned} \quad (3.1)$$

As a result, a weighted average of nonexpansive operators $\theta F_1 + (1 - \theta)F_2$ with $\theta \in [0, 1]$ is also nonexpansive. If one of them is a contraction and $\theta \in (0, 1)$, then the weighted average is a contraction. \square

Contractions are important for us because they have a single fixed point and we can use this property to design iterative algorithms that converge to it.

Theorem 41. *If F is nonexpansive and $\text{dom}(F) = \mathbb{R}^n$, then its set of fixed points*

$$\{x \in \text{dom}(F) : x = F(x)\}$$

is closed and convex. If F is a contraction and $\text{dom}(F) = \mathbb{R}^n$, it has exactly one fixed point.

Proof. Let $F: \mathbb{R}^n \rightarrow \mathbb{R}^n$ be nonexpansive and denote by X the set of its fixed points. Note that we can also define $X = (I - F)^{-1}(\{0\})$, where I is the identity function. From this definition, X is closed because it is the preimage of a continuous function $(F - I)$ on a closed set $(\{0\})$. To show that it is convex, let $x, y \in \text{dom}(F)$ and $\theta \in [0, 1]$ and define $z = \theta x + (1 - \theta)y$. We will show that $z \in X$. As F is nonexpansive

$$\begin{aligned} \|Fz - x\| &= \|Fz - Fx\| \leq \|z - x\| = (1 - \theta)\|x - y\| \\ \|Fz - y\| &= \|Fz - Fy\| \leq \|z - y\| = \theta\|x - y\| \\ \|x - y\| &\leq \|Fz - x\| + \|Fz - y\| \leq \|x - y\| \end{aligned}$$

The last triangle inequality tells us that Fz is on the line segment between x and y . In particular, as $\|Fz - y\| = \theta\|x - y\|$, we have $Fz = \theta x + (1 - \theta)y = z$ and z is a fixed point of F .

Let $x, y \in X$ be again fixed points of F and F be now a contraction with Lipschitz constant $L < 1$, then

$$\|x - y\| = \|Fx - Fy\| \leq L\|x - y\|.$$

This is a contradiction unless $x = y$, which implies that there is a single fixed-point. \square

There are few examples of contractions that are common in convex optimization. Most of the time we work with nonexpansive operators. Among them, a particular type called averaged operator, is specially useful and common.

Definition An operator G is **averaged** if $G = (1 - \alpha)I + \alpha R$ for some $\alpha \in (0, 1)$ and nonexpansive R .

G is nonexpansive because it is a convex combination of nonexpansive operators (the identity I is nonexpansive). Moreover, it is easy to see that G has the same fixed points as R .

$$u^* = Ru^* \Leftrightarrow (1 - \alpha)u^* + \alpha u^* = (1 - \alpha)u^* + \alpha Ru^* \Leftrightarrow u^* = [(1 - \alpha)I + \alpha R]u^* = Gu^* \quad (3.2)$$

We will use this property to design algorithms that find fixed points of nonexpansive operators R and are parametrized by $\alpha \in (0, 1)$.

Properties of Averaged Operators

Lemma 42. *If a function $G : \mathbb{R}^n \rightarrow \mathbb{R}^n$ is averaged with respect to $\alpha \in (0, 1)$, then it is also averaged with respect to any other parameter $\tilde{\alpha} \in (0, \alpha)$.*

Proof. Since G is averaged with respect to α there exists a nonexpansive operator R such that $G = \alpha I + (1 - \alpha)R$. We find

$$\begin{aligned} G &= \alpha I + (1 - \alpha)R \\ &= \tilde{\alpha}I + (\alpha - \tilde{\alpha})I + (1 - \alpha)R \\ &= \tilde{\alpha}I + (1 - \tilde{\alpha}) \underbrace{\left(\frac{\alpha - \tilde{\alpha}}{1 - \tilde{\alpha}}I + \frac{1 - \alpha}{1 - \tilde{\alpha}}R \right)}_{=: \tilde{R}}. \end{aligned}$$

And \tilde{R} is still nonexpansive because

$$\begin{aligned} \|\tilde{R}(u) - \tilde{R}(v)\| &\leq \frac{\alpha - \tilde{\alpha}}{1 - \tilde{\alpha}}\|u - v\| + \frac{1 - \alpha}{1 - \tilde{\alpha}}\|R(u) - R(v)\| \\ &\leq \frac{\alpha - \tilde{\alpha}}{1 - \tilde{\alpha}}\|u - v\| + \frac{1 - \alpha}{1 - \tilde{\alpha}}\|u - v\| \\ &= \|u - v\|. \end{aligned}$$

□

Lemma 43. *If $G_1 : \mathbb{R}^n \rightarrow \mathbb{R}^n$ and $G_2 : \mathbb{R}^n \rightarrow \mathbb{R}^n$ are averaged, then $G_2 \circ G_1$ is also averaged.*

Proof. Let $G_1 = \alpha_1 I + (1 - \alpha_1)R_1$ and $G_2 = \alpha_2 I + (1 - \alpha_2)R_2$ for nonexpansive operators R_1 and R_2 . Then

$$\begin{aligned} G_2(G_1)(u) &= \alpha_2 G_1(u) + (1 - \alpha_2)R_2(G_1(u)) \\ &= \alpha_1 \alpha_2 u + \alpha_2(1 - \alpha_1)R_1(u) + (1 - \alpha_2)R_2(G_1(u)) \\ &= \alpha_1 \alpha_2 u + (1 - \alpha_1 \alpha_2) \left(\frac{\alpha_2(1 - \alpha_1)}{1 - \alpha_1 \alpha_2} R_1(u) + \frac{(1 - \alpha_2)}{1 - \alpha_1 \alpha_2} R_2(G_1(u)) \right). \end{aligned}$$

Since the concatenation of nonexpansive operators is nonexpansive, and convex combinations of nonexpansive operators are nonexpansive, we conclude that $G_2 \circ G_1$ is averaged. □

It is possible to determine if an operator is averaged without explicitly finding its decomposition into a convex combination of the identity and a nonexpansive operator. We do so through the notion of firmly nonexpansive operators.

Definition A function $G : \mathbb{R}^n \rightarrow \mathbb{R}^n$ is called **firmly nonexpansive**, if for all $u, v \in \mathbb{R}^n$ it holds that

$$\|G(u) - G(v)\|_2^2 \leq \langle G(u) - G(v), u - v \rangle.$$

Lemma 44. *A function $G : \mathbb{R}^n \rightarrow \mathbb{R}^n$ is firmly nonexpansive if and only if G is averaged with $\alpha = \frac{1}{2}$.*

Proof. First, let G be averaged with $\alpha = 1/2$, i.e., $G = \frac{1}{2}I + \frac{1}{2}R$ for some nonexpansive operator $R = 2G - I$. As R is nonexpansive, we have

$$\begin{aligned} \|u - v\|_2^2 &\geq \|R(u) - R(v)\|^2 = \|2G(u) - 2G(v) - (u - v)\|^2 \\ &= 4\|G(u) - G(v)\|^2 - 4\langle G(u) - G(v), u - v \rangle + \|u - v\|^2, \end{aligned}$$

which implies $\langle G(u) - G(v), u - v \rangle \geq \|G(u) - G(v)\|_2^2$ and shows that G is firmly nonexpansive.

Second, let G be firmly nonexpansive and define $R = 2G - I$, then

$$\begin{aligned} \|R(u) - R(v)\|^2 &= \|2G(u) - 2G(v) - (u - v)\|^2 \\ &= 4\|G(u) - G(v)\|^2 - 4\langle G(u) - G(v), u - v \rangle + \|u - v\|^2 \\ &\leq \|u - v\|^2, \end{aligned}$$

which shows that R is nonexpansive, i.e., $G = \frac{1}{2}I + \frac{1}{2}R$ is averaged with $\alpha = 1/2$. \square

3.2 Fixed-point Iterations

We are now ready to discuss the main algorithm of this chapter.

Definition Let $G : \mathbb{R}^n \rightarrow \mathbb{R}^n$, and $u^0 \in \mathbb{R}^n$ be a starting point, the **fixed-point or Picard iteration** is

$$u^{k+1} = G(u^k).$$

As the name suggests, the fixed-point iteration is used to find a fixed point u of G . Using this iteration to solve an optimization problem involves two steps: 1) find a suitable G whose fixed points are solutions to the problem at hand, 2) show that the iteration converges to a fixed point. For this second step, we show two simple conditions that guarantee convergence.

Theorem 45. Banach fixed-point theorem. *If the update rule $G : \mathbb{R}^n \rightarrow \mathbb{R}^n$ is a contraction with Lipschitz constant $L < 1$, then the fixed-point iteration converges to the unique fixed-point \hat{u} of G with*

$$\|u^k - \hat{u}\| \leq L^k \|u^0 - \hat{u}\|.$$

Theorem 46. Krasnosel'skii-Mann Theorem. *If the operator $G : \mathbb{R}^n \rightarrow \mathbb{R}^n$ is averaged and has a fixed-point, then the iteration*

$$u^{k+1} = G(u^k)$$

converges to a fixed point of G for any starting point $u^0 \in \mathbb{R}^n$.

Proof. We'll make use of the identity

$$\|(1 - \theta)a + \theta b\|^2 = (1 - \theta)\|a\|^2 + \theta\|b\|^2 - \theta(1 - \theta)\|a - b\|^2,$$

which holds for any $\theta \in \mathbb{R}$, $a, b \in \mathbb{R}^n$. It can be verified by expanding both sides as a quadratic function of θ . The first two terms correspond to the definition of convexity for function $\|\cdot\|^2$, the third one improves this bound.

Because G is averaged, there exists a non-expansive mapping $T : \mathbb{R}^n \rightarrow \mathbb{R}^n$ such that $G = (1-\theta)I + \theta T$. Recall that T has the same fixed points as F . We consider the fixed point iteration

$$u^{k+1} = G(u^k) = (1-\theta)u^k + \theta T u^k.$$

Let U be the (nonempty) set of fixed-points of G and $u^* \in U$, we have then $G(u^*) = u^*$ and

$$\begin{aligned} \|u^{k+1} - u^*\|^2 &= \|(1-\theta)(u^k - u^*) + \theta(Tu^k - u^*)\|^2 \\ &= (1-\theta)\|u^k - u^*\|^2 + \theta\|Tu^k - u^*\|^2 - \theta(1-\theta)\|Tu^k - u^k\|^2 \\ &= (1-\theta)\|u^k - u^*\|^2 + \theta\|Tu^k - Tu^*\|^2 - \theta(1-\theta)\|Tu^k - u^k\|^2 \quad (*) \\ &\leq (1-\theta)\|u^k - u^*\|^2 + \theta\|u^k - u^*\|^2 - \theta(1-\theta)\|Tu^k - u^k\|^2 \\ &= \|u^k - u^*\|^2 - \theta(1-\theta)\|Tu^k - u^k\|^2 \end{aligned}$$

This shows that the distance to the solution set decreases at each step. We call this property Fejèr monotonicity.

Applying the inequality k times yields

$$\|u^{k+1} - u^*\|^2 \leq \|u^0 - u^*\|^2 - \theta(1-\theta) \sum_{j=0}^k \|Tu^j - u^j\|^2$$

and hence

$$\sum_{j=0}^k \|Tu^j - u^j\|^2 \leq \frac{\|u^0 - u^*\|^2 - \|u^{k+1} - u^*\|^2}{\theta(1-\theta)} \leq \frac{\|u^0 - u^*\|^2}{\theta(1-\theta)}.$$

As the upper bound does not depend on k , the series of non-negative terms remains bounded as $k \rightarrow \infty$ and we conclude that $\|Tu^k - u^k\| \rightarrow 0$ as $k \rightarrow \infty$.

From that we can also estimate a convergence rate of the fixed-point residual:

$$\min_{j=0 \dots k} \|Tu^j - u^j\|^2 \leq \frac{\|u^0 - u^*\|^2}{(k+1)\theta(1-\theta)},$$

Since the iterates $\{u^k\}_{k=1}^\infty$ lie in the compact set

$$C = \{v \mid \|v - u^*\| \leq \|u^0 - u^*\|\},$$

there exists at least one subsequence $\{u^{k_l}\}_{l=1}^\infty$ which converges to some point \hat{u} .

Since $Tu^{k_l} - u^{k_l} \rightarrow 0$, we have $Gu^{k_l} - u^{k_l} = (G - I)u^{k_l} \rightarrow 0$ and, as $G - I$ is Lipschitz continuous because T is nonexpansive, we have that $G\hat{u} = \hat{u}$ and the subsequence converges to a point in $\hat{u} \in U$.

As (*) holds for any point from $u^* \in U$, we can apply it the point \hat{u} our subsequence converges to. We know that for the iterates of the original sequence the distance to this point is monotonically decreasing,

$$\|u^{k+1} - \hat{u}\| \leq \|u^k - \hat{u}\|.$$

Since a subsequence $\{u^{k_l}\}_{l=1}^\infty$ of $\{u^k\}_{k=1}^\infty$ is converging to \hat{u} , and $\|u^k - \hat{u}\|$ is monotonically decreasing, we have convergence of the entire sequence to \hat{u} . \square

3.3 Gradient Descent as an Averaged Operator

Given a differentiable convex function $E : \mathbb{R}^n \rightarrow \mathbb{R}$, consider the problem

$$u \in \arg \min_{u \in \mathbb{R}^n} E(u).$$

The first-order optimality conditions of the problem characterize the solution u^* by

$$\nabla E(u^*) = 0 \iff u^* = (I - \tau \nabla E)u^*$$

for any $\tau \neq 0$. The fixed-point iteration for this setup is

$$u^{k+1} = u^k - \tau \nabla E(u^k).$$

This algorithm is called **gradient descent** with a constant step size $\tau > 0$. To guarantee convergence of this fixed-point iteration, we need to determine under which conditions $(I - \tau \nabla E)$ is a contraction or an averaged operator. To this purpose, we will use the following result.

Theorem 47. Baillon-Haddad theorem. *A continuously differentiable convex function $E : \mathbb{R}^n \rightarrow \mathbb{R}$ is L -smooth if and only if $\frac{1}{L}\nabla E$ is firmly nonexpansive, i.e.*

$$\langle \nabla E(u) - \nabla E(v), u - v \rangle \geq \frac{1}{L} \|\nabla E(u) - \nabla E(v)\|_2^2$$

for all $u, v \in \mathbb{R}^n$.

Proof. See Nesterov, *Introductory Lectures on Convex Optimization*, Theorem 2.1.5. \square

We can now determine the conditions under which gradient descent with a constant step size converges.

Theorem 48. *If $E : \mathbb{R}^n \rightarrow \mathbb{R}$ has a minimizer, is convex, and L -smooth, then the gradient descent iteration with constant step size $\tau \in (0, \frac{2}{L})$ converges to a minimizer.*

Proof. We will show that the fixed-point operator of gradient descent $G(u) = u - \tau \nabla E(u)$ is averaged.

By Baillon-Haddad theorem, we know that $\frac{1}{L}\nabla E$ is firmly non-expansive, or equivalently, averaged with $\alpha = 1/2$. Let $\frac{1}{L}\nabla E = \frac{1}{2}(I + T)$ for a non-expansive T , it holds

$$G(u) = u - \tau L \frac{1}{L} \nabla E(u) = \left(1 - \frac{L\tau}{2}\right) I + \frac{L\tau}{2}(-T)$$

It is clear that if T is non-expansive, $(-T)$ is also non-expansive, and consequently G is averaged for $\tau \in (0, \frac{2}{L})$. \square

Theorem 49. *If $E : \mathbb{R}^n \rightarrow \mathbb{R}$ is strongly convex with parameter m and strongly smooth with parameter L , then the gradient descent iteration with constant step size $\tau \in (0, \frac{2}{L})$ converges to the unique minimizer u^* with geometric convergence rate*

$$\|u^k - u^*\| \leq c^k \|u^0 - u^*\|.$$

Proof. We will show that $(I - \tau \nabla E)$ is Lipschitz with parameter $c = \max\{|1 - \tau m|, |1 - \tau L|\}$. To simplify the proof, we will assume that E is twice continuously differentiable although the result is still true without this assumption. If E is twice continuously differentiable, we have

- $D(I - \tau \nabla E) = I_n - \tau \nabla^2 E$, where I_n is the identity matrix
- m strong convexity is equivalent to $\nabla^2 E \succeq mI_n$
- L -smoothness corresponds to $\nabla^2 E \preceq LI_n$

Putting these together, we have

$$\begin{aligned} (1 - \tau L)I_n &\preceq D(I - \tau \nabla E) \preceq (1 - \tau m)I_n \\ \|D(I - \tau \nabla E)\| &\leq \max\{|1 - \tau m|, |1 - \tau L|\} \\ (I - \tau \nabla E) &\text{ has Lipschitz constant } c = \max\{|1 - \tau m|, |1 - \tau L|\}. \end{aligned} \tag{3.3}$$

As a result, $(I - \tau \nabla E)$ is a contraction for $\tau \in (0, \frac{2}{L})$ and the fixed-point iteration converges to the unique fixed point of the contraction with the geometric rate c^k by Banach fixed-point theorem. \square

3.4 Projected Gradient Descent

Definition Projection For a (nonempty) closed convex set $C \subset \mathbb{R}^n$,

$$\pi_C(v) = \operatorname{argmin}_{u \in C} \|u - v\|_2^2$$

is called the projection of v onto the set C .

In plain English, the projection of a point v onto C is the point in C that is closest to v . As a result, if $v \in C$ then $\pi_C(v) = v$ and the reverse is also true: $\pi_C(v) = v$ if and only if $v \in C$.

As the projection is defined in terms of a minimization problem, it is natural to wonder if the optimization problem has a solution and whether this solution is unique. The convexity of the set C , gives us a positive answer.

Theorem 50. Existence and Uniqueness of the Projection *For any (nonempty) closed convex set $C \subset \mathbb{R}^n$ and any v the projection $\pi_C(v)$ exists and is single valued.*

Proof. To show that $\pi_C(v)$ is not empty, we define

$$E(u) = \begin{cases} \|u - v\|_2^2 & \text{if } u \in C \\ \infty & \text{otherwise} \end{cases}.$$

As C is not empty, we can pick $v_0 \in C$ and define the sublevel set $S_{E(v_0)} = \{u \in \mathbb{R}^n : E(u) \leq E(v_0)\}$. As $v_0 \in S_{E(v_0)}$, the sublevel set is not empty. It is also bounded because any $u \in S_{E(v_0)}$, satisfies

$$\|u\| \leq \|u - v\| + \|v\| = \sqrt{E(u)} + \|v\| \leq \sqrt{E(v_0)} + \|v\|.$$

The closedness of C implies the closedness of $\operatorname{epi}(E)$, and we have already seen that closed functions are l.s.c. As a result, we can use Theorem 23 to prove existence of a minimizer of $E(u)$. The minimizer is unique because $E(u)$ is strictly convex as a result of the convexity of C and the strict convexity of $\|u - v\|_2^2$. \square

As a result of this theorem, although $\pi_C(v)$ is by definition a set, we usually identify $\pi_C(v)$ with the single element in the set.

It is useful to know the form of the projection operator for some common sets. For instance:

- $C = \{u \in \mathbb{R}^n \mid \|u\|_2 \leq 1\}$
- $C = \{u \in \mathbb{R}^n \mid \|u\|_\infty := \max_i |u_i| \leq 1\}$
- $C = \{u \in \mathbb{R}^n \mid u_i \in [a, b]\}$
- $C = \{u \in \mathbb{R}^n \mid u_i \geq a\}$
- $C = \{u \in \mathbb{R}^n \mid \|u\|_1 = \sum_i |u_i|\}$

Projection operators are necessary in optimization to solve problems subject to a closed convex constraint set C

$$u^* \in \operatorname{argmin}_{u \in C} E(u), \tag{3.4}$$

When the objective function E is also convex, the optimization problem is convex, and we know under which condition it has a solution but we do not know how to solve it yet. The projected gradient algorithm is the first step towards this goal.

The projected gradient descent algorithm builds on gradient descent to find a solution of (3.4) when E is convex and L -smooth. To this goal, let us look at the gradient descent update rule

$$u^{k+1} = u^k - \tau \nabla E(u^k).$$

The problem with this update for solving problem (3.4) is that, even if $u^k \in C$, the update u^{k+1} might lie outside the feasible set C . Gradient projection solves this by simply projecting every iteration back to the feasible set with $u^{k+1} = \pi_C(u^k - \tau \nabla E(u^k))$.

Definition Gradient Projection Algorithm Let $C \subset \mathbb{R}^n$ be a nonempty closed convex set and let $E : \mathbb{R}^n \rightarrow \mathbb{R} \in C^1(\mathbb{R}^n)$. Then, for $u^0 \in C$

$$u^{k+1} = \pi_C(u^k - \tau \nabla E(u^k))$$

is called the *gradient projection* algorithm.

Similar to gradient descent, we can write the gradient projection algorithm as a fixed point iteration of an operator

$$G(u) = \pi_C(u - \tau \nabla E(u))$$

to analyze its convergence. In particular, we need to determine under which conditions G is an averaged operator or a contraction to use Banach or Krasnosel'skii-Mann theorems to prove its convergence.

From the analysis of gradient descent, we know that if E is L -smooth and $\tau \in (0, \frac{2}{L})$ the operator

$$G_1(u) = u - \tau \nabla E(u)$$

is averaged. If we know recollect that the composition of averaged operators is also averaged, we only need to show that the projection π_C is averaged. In fact, we will see that it is firmly nonexpansive, and therefore, it is averaged with $\alpha = \frac{1}{2}$.

Theorem 51. *The projection π_C onto a nonempty closed convex set $C \subset \mathbb{R}^n$ is firmly nonexpansive, i.e. it meets*

$$\langle u - v, \pi_C(u) - \pi_C(v) \rangle \geq \|\pi_C(u) - \pi_C(v)\|^2 \quad \forall u, v \in \mathbb{R}^n.$$

Proof. Let δ_C be the indicator function of the convex set C , which is defined as

$$\delta_C(v) = \begin{cases} 0 & \text{if } v \in C \\ \infty & \text{otherwise} \end{cases}.$$

As C is not empty, closed, and convex, δ_C is proper, closed, and convex. We can now write

$$\pi_C(u) = \operatorname{argmin}_{z \in \mathbb{R}^n} \delta_C(z) + \|z - u\|_2^2$$

From the optimality conditions of the optimization problem we have

$$\begin{aligned} u - \pi_C(u) &\in \partial \delta_C(\pi_C(u)) \\ v - \pi_C(v) &\in \partial \delta_C(\pi_C(v)). \end{aligned}$$

At the same time, recall the definition of the subgradient of function E

$$E(z) - E(x) \geq \langle p, z - x \rangle \quad \forall z \quad p \in \partial E(x).$$

If we apply this inequality with $E = \delta$ at the points $x = \pi_C(u)$ and $x = \pi_C(v)$, we have

$$\begin{aligned}\delta_C(z) &= \delta_C(z) - \delta_C(\pi_C(u)) \geq \langle u - \pi_C(u), z - \pi_C(u) \rangle \quad \forall z \\ \delta_C(z) &= \delta_C(z) - \delta_C(\pi_C(v)) \geq \langle v - \pi_C(v), z - \pi_C(v) \rangle \quad \forall z.\end{aligned}\tag{3.5}$$

If we choose $z = \pi_C(v)$ for the first inequality and $z = \pi_C(u)$ for the second and add both inequalities, we have $\delta_C(z) = 0$ and

$$\begin{aligned}0 &\geq \langle u - \pi_C(u) + \pi_C(v) - v, \pi_C(v) - \pi_C(u) \rangle \\ 0 &\geq \langle u - v, \pi_C(v) - \pi_C(u) \rangle + \|\pi_C(v) - \pi_C(u)\|^2 \\ \langle u - v, \pi_C(u) - \pi_C(v) \rangle &\geq \|\pi_C(u) - \pi_C(v)\|^2.\end{aligned}$$

□

We can now state the main convergence result of projected gradient algorithm.

Theorem 52. *For an L -smooth energy E that has a minimizer and a choice $\tau \in (0, \frac{2}{L})$ the gradient projection converges to a solution of*

$$u^* \in \arg \min_{u \in C} E(u)\tag{3.6}$$

with convergence rate $\mathcal{O}(1/k)$.

The convergence rate $\mathcal{O}(1/k)$ of the vanilla projected gradient algorithm is suboptimal, but it can be improved to $\mathcal{O}(1/k^2)$ with acceleration techniques that introduce an extrapolation step exploiting the L -smoothness of E . We can also improve this rate if our objective function is m -strongly convex.

Theorem 53. *For E being L -smooth and m -strongly convex and $\tau \in (0, \frac{2}{L})$ the gradient projection algorithm converges to the (unique) global minimizer u^* with $E(u^k) - E(u^*) \in \mathcal{O}(c^k)$ with $c < 1$.*

Proof. Recall that the composition of a non-expansive operator with a contraction is a contraction. As a result, whenever $G_1(u) = u - \tau \nabla E(u)$ is a contraction, the gradient projection operator $\pi_C(G_1(u))$ is a contraction and we can use Banach fixed-point theorem to prove convergence with a linear rate. □

3.5 Proximal Gradient Descent

Definition Proximal Operator Given a closed, proper, convex function $E : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{\infty\}$, the mapping $\text{prox}_E : \mathbb{R}^n \rightarrow \mathbb{R}^n$ defined as

$$\text{prox}_E(v) := \underset{u \in \mathbb{R}^n}{\operatorname{argmin}} E(u) + \frac{1}{2} \|u - v\|^2$$

is called the proximal operator or proximal mapping of E .

The proximal operator is a generalization of the projection. Indeed, given a nonempty, closed, convex set C , the projection π_C is the proximal operator of the indicator function

$$\delta_C(v) = \begin{cases} 0 & \text{if } v \in C \\ \infty & \text{otherwise} \end{cases}.$$

Many of the properties of the projection are inherited by the proximal operator.

Lemma 54. *Given a closed, proper, and convex function $E: \mathbb{R}^n \rightarrow \bar{\mathbb{R}}$ and $u \in \mathbb{R}^n$ there exists a unique proximal $\text{prox}_E(u)$.*

Proof. We have already proved the existence of minimizers of convex functions that are coercive, that is, functions that have bounded sublevel sets. In the case of the proximal operator, it is easy to see that $E(u) + \frac{1}{2} \|u - v\|^2$ is convex if E is convex and that it has bounded sublevel sets because

$$E(u) + \frac{1}{2} \|u - v\|^2 < \gamma \Rightarrow \|u - v\|^2 < 2\gamma.$$

Uniqueness is a consequence of the strong convexity of $(1/2) \|u - v\|^2$, which implies strong convexity of $E(u) + (1/2) \|u - v\|^2$ \square

Similarly, the proximal operator is also firmly nonexpansive, as the next theorem shows.

Theorem 55. *The proximal operator prox_E for a closed, proper, convex function E is firmly nonexpansive, that is, it satisfies*

$$\langle u - v, \text{prox}_E(u) - \text{prox}_E(v) \rangle \geq \|\text{prox}_E(u) - \text{prox}_E(v)\|^2 \quad \forall u, v \in \mathbb{R}^n.$$

Proof. Let $x = \text{prox}_E(u)$ and $y = \text{prox}_E(v)$, the optimality conditions then

$$\begin{aligned} x &= \min_z E(z) + \|z - u\|^2 \Rightarrow u - x \in \partial E(x) \\ x &= \min_z E(z) + \|z - v\|^2 \Rightarrow v - y \in \partial E(y). \end{aligned} \tag{3.7}$$

At the same time, by definition of the subgradient we have

$$\begin{aligned} E(z) - E(x) &\geq \langle \partial E(x), z - x \rangle = \langle u - x, z - x \rangle \quad \forall z \\ E(z) - E(y) &\geq \langle \partial E(y), z - y \rangle = \langle v - y, z - y \rangle \quad \forall z, \end{aligned} \tag{3.8}$$

where in the second step of the inequalities we have used that $u - x \in \partial E(x)$ and $v - y \in \partial E(y)$. If we now choose $z = y$ for the first inequality and $z = x$ for the second, we have

$$\begin{aligned} E(y) - E(x) &\geq \langle u - x, y - x \rangle \\ E(x) - E(y) &\geq \langle v - y, x - y \rangle \end{aligned}$$

adding both inequalities we have

$$\begin{aligned} 0 &\geq \langle u - x + y - v, y - x \rangle \\ 0 &\geq \langle u - v, y - x \rangle + \|y - x\|^2 \\ \langle u - v, x - y \rangle &\geq \|x - y\|^2. \\ \langle u - v, \text{prox}_E(u) - \text{prox}_E(v) \rangle &\geq \|\text{prox}_E(u) - \text{prox}_E(v)\|^2. \end{aligned}$$

\square

For many common convex objective functions, the proximal operators is simple to compute and has a closed-form expression.

- Quadratic functions

$$f(x) = \frac{1}{2} \|Ax - b\|^2, \quad \text{prox}_{\tau f}(v) = (I + \tau A^T A)^{-1}(v - \tau b)$$

- Euclidean norm

$$f(x) = \|x\|, \quad \text{prox}_{\tau f}(v) = \begin{cases} (1 - \tau/\|v\|)v & \text{if } \|v\| \geq \tau \\ 0 & \text{otherwise.} \end{cases}$$

- ℓ_1 -norm (cf. exercise sheet 3)

$$f(x) = \|x\|_1, \quad (\text{prox}_{\tau f}(v))_i = \begin{cases} v_i + \tau & \text{if } v_i < -\tau \\ 0 & \text{if } |v_i| \leq \tau \\ v_i - \tau & \text{if } v_i > \tau. \end{cases}$$

In the same way as we used the projection to generalize gradient descent to constrained minimization problems, we use the proximal operator to generalize gradient descent to optimization problems of the form

$$E(u) = E_1(u) + E_2(u),$$

where both E_1 and E_2 are proper, closed, and convex and satisfy

- $E_1 : \mathbb{R}^n \rightarrow \mathbb{R}$ is L -smooth.
- $E_2 : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{\infty\}$ has an easy-to-evaluate proximal operator.

In this case, can generalize the projected gradient algorithm by taking gradient descent steps on E_1 and proximal steps on E_2 . This strategy is known as proximal gradient method.

Definition Proximal Gradient Method For a closed, proper, convex function $E_1, E_2 : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{\infty\}$, where E_1 is differentiable, and given an initial point $u^0 \in \mathbb{R}^n$ and a step size τ , the algorithm

$$u^{k+1} = \text{prox}_{\tau E_2}(u^k - \tau \nabla E_1(u^k)), \quad k = 0, 1, 2, \dots,$$

is called the *proximal gradient method* or *forward-backward splitting*.

For a constant E_2 , the proximal gradient method reduces to gradient descent, while for $E_2 = \delta_C$ it reduces to projected gradient descent. The case we have not seen, constant E_1 , results in the *proximal point algorithm*.

We can again prove analyze the convergence of the proximal gradient method as the fixed-point iteration of the operator

$$G(u) = \text{prox}_{\tau E_2}(u - \tau \nabla E_1(u)).$$

We do so by determining the conditions under which G is averaged or a contraction.

Theorem 56. For closed, proper, convex functions E_1 and E_2 , with E_1 L -smooth and having a minimizer u^* of $E(u) = E_1(u) + E_2(u)$, the proximal gradient method with constant step size $\tau \in (0, \frac{2}{L})$ converges to u^* with rate $E(u^k) - E(u^*) \in \mathcal{O}(1/k)$.

Proof. The operator G is averaged for $\tau \in (0, \frac{2}{L})$ because it is the composition of two averaged operators, the gradient-descent operator $G_1(u) = u - \tau \nabla E_1(u)$ and the proximal operator $\text{prox}_{\tau E_2}$. Indeed, we have already seen that if E_1 is L -smooth closed, proper, convex, then G_1 is averaged for $\tau \in (0, \frac{2}{L})$. At the same time, $\text{prox}_{\tau E_2}$ is averaged with $\alpha = \frac{1}{2}$ because it is firmly nonexpansive. The convergence rate results from particularizing the analysis of the convergence rate of fixed-point iterations of an averaged operator to the proximal gradient. \square

As we have done with gradient descent and the projected gradient algorithms, we can obtain linear convergence with some additional assumptions on the objective function. To this goal, we need to determine under which conditions the proximal operator is not only nonexpansive but a contraction.

Theorem 57. *If the proper, closed function E is m -strongly convex, then $\text{prox}_{\tau E} : \mathbb{R}^n \rightarrow \mathbb{R}^n$ is a contraction.*

As the composition of a contraction with a nonexpansive mapping is a contraction, $\text{prox}_{\tau E_2}(u - \tau \nabla E_1)$ is a contraction whenever $(u - \tau \nabla E_1)$ or $\text{prox}_{\tau E}$ are a contraction. This condition translates into strong convexity of E_1 and E_2 .

Corollary 58. *If E_1 is L -smooth, $\tau \in (0, \frac{2}{L})$, and either E_1 or E_2 is strongly convex, then the proximal gradient method converges linearly, i.e., $\|u^k - u^*\|_2^2 \in \mathcal{O}(c^k)$ for some $c < 1$.*

Proof. This is an immediate result of Banach fixed-point iteration theorem. □

Chapter 4

Duality

Functional transforms can shed light into an optimization problem by presenting it from a different perspective. In convex analysis, this new perspective is the result of duality. In a nutshell, every convex problem has a dual version, and their solutions are related by a linear transform.

We will start by introducing the convex conjugate of a function $E(u)$, which is convex, and then show how conjugating twice a convex functions give us the original function reformulated as a maximization problem over a dual variable p . As a result, the minimization problem over u can be reformulated into a saddle-point problem

$$\min_u E(u) \iff \min_u \max_p \mathcal{L}(u, p).$$

Under certain conditions, we will see that we can additionally swap min and max

$$\min_u E(u) \iff \min_u \max_p \mathcal{L}(u, p) \iff \max_p \min_u \mathcal{L}(u, p).$$

If we can now analytically solve the interior minimization problem

$$u^* = \operatorname{argmin}_u \mathcal{L}(u, p)$$

then we have found a closed-form formulation of the dual problem

$$\min_u E(u) \iff \max_p \mathcal{L}(u^*, p).$$

The dual problem usually has complementary properties to the primal (the original) one. For instance, if the primal problem is the composition of affine transforms with simple convex functions, the dual tends to have a larger number of optimization variables but has simple proximal operators and is well suited to a proximal gradient method. Duality is also useful when the primal problem is not differentiable but is m -strongly convex, because its dual is $\frac{1}{m}$ -smooth and can be solved with simple gradient descent. Duality increases the set of problems that we can solve with the algorithms we have studied so far, or simply let us do it more efficiently, with lower memory requirements, or operation counts. This forms the first part of this chapter.

When looking at dual problem does not help us solve the primal one better, we look at the saddle point problem to try to combine the best of the primal and the dual worlds. This results in the primal-dual algorithms that we will cover in the second part of this chapter.

4.1 Convex Conjugate

Definition Let $E : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{\infty\}$ be any function, not necessarily convex, we define its *convex conjugate* to be

$$E^*(p) = \sup_{u \in \mathbb{R}^n} [\langle u, p \rangle - E(u)].$$

As the name suggests, the convex conjugate of a function is convex, even when the function is not. In fact, it is also closed.

Lemma 59. Convexity of the Convex Conjugate *The convex conjugate*

$$E^*(p) = \sup_{u \in \mathbb{R}^n} (\langle u, p \rangle - E(u)).$$

of any proper function $E : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{\infty\}$ is convex. If the function is closed, its conjugate is also closed.

Proof. The convexity of E^* results from its definition as a supremum. We can prove it, by proving the convexity of its epigraph. Let $(p, \alpha) \in \text{epi}(E^*)$, $(q, \beta) \in \text{epi}(E^*)$, by definition of epigraph we have

$$\alpha \geq \sup_{u \in \mathbb{R}^n} [\langle u, p \rangle - E(u)] \quad (4.1)$$

$$\beta \geq \sup_{u \in \mathbb{R}^n} [\langle u, q \rangle - E(u)]. \quad (4.2)$$

Multiplying the first inequality by $\theta \in [0, 1]$ and the second by $(1 - \theta)$ and adding them up, we obtain

$$\begin{aligned} \theta\alpha + (1 - \theta)\beta &\geq \sup_{u \in \mathbb{R}^n} [\langle u, \theta p \rangle - \theta E(u)] + \sup_{u \in \mathbb{R}^n} [\langle u, (1 - \theta)q \rangle - (1 - \theta)E(u)] \\ &\geq \sup_{u \in \mathbb{R}^n} [\langle u, \theta p \rangle - \theta E(u) + \langle u, (1 - \theta)q \rangle - (1 - \theta)E(u)] \\ &\geq \sup_{u \in \mathbb{R}^n} [\langle u, \theta p + (1 - \theta)q \rangle - E(u)] \end{aligned}$$

which shows that $\theta(p, \alpha) + (1 - \theta)(q, \beta) \in \text{epi}(E^*)$. To show that it is closed, we will show that its epigraph is the intersection of an arbitrary number of closed convex sets $\text{epi}(E^*) = \{(p, \alpha) \in \mathbb{R}^n \times \mathbb{R} : \langle u, p \rangle - E(u) \leq \alpha \forall u\} = \bigcap_u \text{epi}E_u$ where $E_u(\cdot) = \langle u, \cdot \rangle - E(u)$. Since each $\text{epi}E_u$ is closed, and any arbitrary intersection of closed sets is closed, $\text{epi}(E^*)$ is closed. \square

We will mostly study convex conjugates of convex functions. Some of them are classic examples that is useful to know instead of derive every single time.

- If $E(u) = \frac{1}{2}\|u\|^2$, the optimality conditions of

$$\sup_u \langle p, u \rangle - \frac{1}{2}\|u\|^2$$

show that the supremum is attained at $\hat{u} = p$, where $\langle p, \hat{u} \rangle - \frac{1}{2}\|\hat{u}\|^2 = \frac{1}{2}\|p\|^2$. This yields $E^*(p) = \frac{1}{2}\|p\|^2$.

- If $\|\cdot\|_*$ is the dual norm of $\|\cdot\|$, the convex conjugate of $E(u) = \|u\|$ is $E^*(p) = \begin{cases} 0 & \text{if } \|p\|_* \leq 1, \\ \infty & \text{else.} \end{cases}$

Recall the definition of the dual norm: given $p \in \mathbb{R}^n$

$$\|p\|_* = \sup\{\langle p, u \rangle : u \in \mathbb{R}^n, \|u\| \leq 1\}.$$

As a result, if $\|p\|_* > 1$, there exists $x \in \mathbb{R}^n$ with $\|x\| < 1$ such that $\langle x, p \rangle > 1$ and $\langle p, x \rangle - \|x\| > 0$. Now, define $z = tx$, we have

$$\sup_{u \in \mathbb{R}^n} \langle p, u \rangle - \|u\| \geq \sup_t \langle p, tx \rangle - \|tx\| = \sup_{t>0} t[\langle p, x \rangle - \|x\|] = \infty.$$

Conversely, if $\|p\|_* \leq 1$, we have $\langle p, \frac{u}{\|u\|} \rangle \leq 1$ for all u , which implies $\langle p, u \rangle \leq \|u\|$ for all u . Therefore $u = 0$ is the value that maximizes $\langle p, u \rangle - \|u\|$ with maximum value 0.

In particular, we have

- The conjugate of $E(u) = \|u\|_2$ is $E^*(p) = \begin{cases} 0 & \text{if } \|p\|_2 \leq 1, \\ \infty & \text{else.} \end{cases}$
- The conjugate of $E(u) = \|u\|_1$ is $E^*(p) = \begin{cases} 0 & \text{if } \|p\|_\infty \leq 1, \\ \infty & \text{else.} \end{cases}$
- The conjugate of $E(u) = \|u\|_\infty$ is $E^*(p) = \begin{cases} 0 & \text{if } \|p\|_1 \leq 1, \\ \infty & \text{else.} \end{cases}$

- The convex conjugate of the indicator function of the unit ball $E(u) = \begin{cases} 0 & \text{if } \|u\| \leq 1, \\ \infty & \text{else.} \end{cases}$ is the dual norm $E^*(p) = \|p\|_*$. Indeed

$$\sup_{u \in \mathbb{R}^n} \langle p, u \rangle - E(u) = \sup_{\|u\| \leq 1} \langle p, u \rangle = \|p\|_*.$$

In particular, we have

- $E(u) = \begin{cases} 0 & \text{if } \|u\|_2 \leq 1, \\ \infty & \text{else.} \end{cases}$ leads to $E^*(p) = \|p\|_2$.
- $E(u) = \begin{cases} 0 & \text{if } \|u\|_\infty \leq 1, \\ \infty & \text{else.} \end{cases}$ leads to $E^*(p) = \|p\|_1$.
- $E(u) = \begin{cases} 0 & \text{if } \|u\|_1 \leq 1, \\ \infty & \text{else.} \end{cases}$ leads to $E^*(p) = \|p\|_\infty$.

Now that we know some basic conjugates, it is useful to investigate how the conjugation affects some basic operations like linear composition, scaling or affine transforms to increase our repertoire of conjugates.

- **Scalar multiplication** :If $E(u) = \alpha \tilde{E}(u)$

$$E^*(p) = \sup_u \langle p, u \rangle - \alpha \tilde{E}(u) = \alpha \sup_u \langle \frac{p}{\alpha}, u \rangle - \tilde{E}(u) = \alpha \tilde{E}^*(p/\alpha).$$

- **Separable sum**: If $E(u_1, u_2) = E_1(u_1) + E_2(u_2)$

$$\begin{aligned} E^*(p) &= \sup_{p=(p_1, p_2)} \langle p_1, u_1 \rangle + \langle p_2, u_2 \rangle - E_1(u_1) - E_2(u_2) \\ &= \sup_{p_1} \langle p_1, u_1 \rangle - E_1(u_1) + \sup_{p_2} \langle p_2, u_2 \rangle - E_2(u_2) \\ &= E_1^*(p_1) + E_2^*(p_2). \end{aligned}$$

- **Sum rule**: If E_1, E_2 are closed, convex, proper and $E(u) = E_1(u) + E_2(u)$

$$\begin{aligned} E^*(p) &= \sup_p \langle p, u \rangle - E_1(u) - E_2(u) \\ &= \sup_{p=p_1+p_2} \langle p_1, u \rangle - E_1(u) + \langle p_2, u \rangle - E_2(u) \\ &= \inf_{p=p_1+p_2} \sup_{p_1} \langle p_1, u \rangle - E_1(u) + \sup_{p_2} \langle p_2, u \rangle - E_2(u) \\ &= \inf_{p=p_1+p_2} E_1^*(p_1) + E_2^*(p_2). \end{aligned}$$

Where we have used that

$$\sup_{p=p_1+p_2} F(p_1, p_2) = \inf_{p=p_1+p_2} \sup_{p_1} \sup_{p_2} F(p_1, p_2)$$

- **Translation:** If $E(u) = \tilde{E}(u - b)$

$$\begin{aligned} E^*(p) &= \sup_u \langle p, u \rangle - \tilde{E}(u - b) = \sup_u \langle p, b \rangle + \langle p, u - b \rangle - \tilde{E}(u - b) \\ &= \sup_{u-b} \langle p, b \rangle + \langle p, u - b \rangle - \tilde{E}(u - b) \\ &= \langle p, b \rangle + \tilde{E}^*(p). \end{aligned}$$

- **Additional affine functions:** If $E(u) = \tilde{E}(u) + \langle b, u \rangle + a$

$$\begin{aligned} E^*(p) &= \sup_u \langle p, u \rangle - \tilde{E}(u) - \langle b, u \rangle - a = \sup_u -a + \langle p - b, u \rangle - \tilde{E}(u) \\ &= -a + \tilde{E}^*(p - b). \end{aligned}$$

4.2 Duality Theorems

Theorem 60 (Fenchel-Young Inequality). *Let E be proper, convex and closed, $u \in \text{dom}(E) \subset \mathbb{R}^n$, and $p \in \mathbb{R}^n$, then*

$$E(u) + E^*(p) \geq \langle u, p \rangle.$$

Equality holds if and only if $p \in \partial E(u)$.

Proof. The inequality follows immediately from the definition of the conjugate

$$E(u) + E^*(p) = E(u) + \sup_v \langle v, p \rangle - E(v) \geq E(u) + \langle u, p \rangle - E(u) = \langle u, p \rangle.$$

To show the equality statement, we will show the remaining inequality

$$E(u) + E^*(p) \leq \langle u, p \rangle,$$

or, in other words,

$$E(u) + \langle p, z \rangle - E(z) \leq \langle u, p \rangle, \quad \forall z.$$

Rewritten, the above is nothing but

$$E(z) - E(u) - \langle p, z - u \rangle \geq 0, \quad \forall z,$$

which is simply the definition of the subgradient $p \in \partial E(u)$. □

Theorem 61 (Biconjugate). *Let $E : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{\infty\}$ be proper, convex and closed, then $E^{**} = E$.*

Proof. We'll show an incomplete proof that only considers the relative interior but gives already a quick intuition of why the statement makes sense. For the full proof, please check Rockafellar's book *Convex Analysis*, Theorem 12.2.

First of all, note that $E^{**}(u) \leq E(u)$ because

$$E^{**}(u) = \sup_p \langle p, u \rangle - E^*(p) = \sup_p \langle p, u \rangle - \sup_v [\langle p, v \rangle - E(v)] \leq \sup_p \langle p, u \rangle - [\langle p, u \rangle - E(u)] = E(u).$$

If E is subdifferentiable at u , let $q \in \partial E(u)$. Fenchel-Young inequality tells us that $E(u) + E^*(q) = \langle u, q \rangle$ and, by definition of the supremum, we have

$$E^{**}(u) = \sup_p \langle p, u \rangle - E^*(p) \geq \langle q, u \rangle - E^*(q) = E(u).$$

Combining both inequalities, we have $E^{**}(u) = E(u)$. □

This give us the first hint on how to reformulate a convex problem into a friendlier version: simply changing E by its biconjugate E^{**} and checking if it is *in some way simpler* to solve. There are many ways to measure simplicity, a useful one is the subdifferential that appears in the optimality conditions of the problem.

Lemma 62. Subgradient of convex conjugate *Let E be proper, convex and closed, then the following two conditions are equivalent:*

- $p \in \partial E(u)$
- $u \in \partial E^*(p)$

Proof. Let $p \in \partial E(u)$, by the Fenchel-Young Inequality we know that

$$E(u) + E^*(p) = \langle u, p \rangle.$$

On the other hand, $E = E^{**}$ such that

$$E^{**}(u) + E^*(p) = \langle u, p \rangle,$$

and the Fenchel-Young Inequality tells us that $u \in \partial E^*(p)$. Similarly, $u \in \partial E^*(p)$ implies $p \in \partial E(u)$. \square

This let us prove an important property, that the convex conjugate of a strongly convex function is smooth. This is important because it means that we can non-differentiable strongly convex problems with a descent technique through its conjugate.

Theorem 63. Conjugation of strongly convex functions *If $E : \mathbb{R}^n \rightarrow \bar{\mathbb{R}}$ is proper, closed and m -strongly convex, then E^* is proper, closed, convex and $1/m$ -smooth.*

Proof. For a proper, closed, strongly convex function, the supremum in the definition of E^* is attained. To see this, let us re-write

$$\max_u \langle u, p \rangle - E(u) = - \min_u E(u) - \langle u, p \rangle$$

Now the convexity of E let us substitute the minimum by an infimum

$$\max_u \langle u, p \rangle - E(u) = - \inf_u E(u) - \langle u, p \rangle = \sup_u \langle u, p \rangle - E(u) = E^*(p).$$

The strong convexity of E also means that the minimum is unique. The optimality condition immediately yields that this minimum is attained for $p \in \partial E(u)$, i.e. for $u \in \partial E^*(p)$. Since the optimal u is unique, the subdifferential $\partial E^*(p)$ is single valued for all p , which yields the differentiability of E^* .

As E is m -strongly convex, $E - \frac{m}{2} \|\cdot\|^2$ is convex. As a result

$$\langle u - v, p - q \rangle \geq m \|u - v\|^2 \quad \forall p \in \partial E(u), q \in \partial E(v),$$

or in other words

$$\langle \nabla E^*(p) - \nabla E^*(q), p - q \rangle \geq m \|\nabla E^*(p) - \nabla E^*(q)\|^2 \quad \forall p, q.$$

Now, by Cauchy-Schwarz inequality we have

$$\|\nabla E^*(p) - \nabla E^*(q)\| \|p - q\| \geq \langle \nabla E^*(p) - \nabla E^*(q), p - q \rangle \geq m \|\nabla E^*(p) - \nabla E^*(q)\|^2 \quad \forall p, q$$

which implies $\frac{1}{m}$ -smoothness of E^* as

$$\|\nabla E^*(p) - \nabla E^*(q)\| \leq \frac{1}{m} \|p - q\| \quad \forall p, q.$$

\square

Theorem 64 (Fenchel's Duality¹). *Let $G : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{\infty\}$ and $F : \mathbb{R}^m \rightarrow \mathbb{R} \cup \{\infty\}$ be proper, closed, convex functions and let there exist a $u \in \text{ri}(\text{dom}(G))$ such that $Ku \in \text{ri}(\text{dom}(F))$. Then*

\inf_u	$G(u) + F(Ku)$	"Primal"
= $\inf_u \sup_q$	$G(u) + \langle q, Ku \rangle - F^*(q)$	"Saddle point"
= $\sup_q \inf_u$	$G(u) + \langle q, Ku \rangle - F^*(q)$	"Saddle point"
= \sup_q	$-G^*(-K^*q) - F^*(q)$	"Dual"

Proof. Partial proof: Let us assume a minimum is attained at some \hat{u} . Our assumptions let us apply the sum rule of the subgradient to compute the optimality conditions of the primal problem

$$\hat{u} \in \underset{u}{\text{argmin}} \ G(u) + F(Ku) \tag{4.3}$$

$$0 \in q + K^*\hat{p} \qquad q \in \partial G(\hat{u}) \qquad \hat{p} \in \partial F(K\hat{u}) \tag{4.4}$$

$$q \in -K^*\hat{p} \qquad \hat{u} \in \partial G^*(q) \qquad K\hat{u} \in \partial F^*(\hat{p}) \tag{4.5}$$

$$\hat{u} \in \partial G^*(-K^*\hat{p}) \qquad K\hat{u} \in \partial F^*(\hat{p}) \tag{4.6}$$

If we know take $\hat{u} \in \partial G^*(-K^*\hat{p})$ and $K\hat{u} \in \partial F^*(\hat{p})$ and use it to write

$$0 = K\hat{u} - K\hat{u} \in -K\partial G^*(-K^*\hat{p}) - \partial F^*(\hat{p})$$

we see that \hat{p} satisfies the optimality conditions of the dual problem $\hat{p} \in \arg \max_p \ -G^*(-K^*p) - F^*(p)$. Moreover, the optimal solution pair (\hat{u}, \hat{p}) satisfies

$$K\hat{u} \in \partial F^*(\hat{p}) \qquad -K^*\hat{p} \in \partial G(\hat{u}) \tag{4.7}$$

$$K\hat{u} \in \partial F^*(\hat{p}) \qquad \hat{u} \in \partial G^*(-K^*\hat{p}) \tag{4.8}$$

$$\hat{p} \in \partial F(K\hat{u}) \qquad \hat{u} \in \partial G^*(-K^*\hat{p}) \tag{4.9}$$

$$\hat{p} \in \partial F(K\hat{u}) \qquad -K^*\hat{p} \in \partial G(\hat{u}). \tag{4.10}$$

□

This immediately give us the following alternative characterizations of the solution of our problem.

Corollary 65. *Let the assumptions from Fenchel's Duality Theorem hold. If there exists a pair $(u, q) \in \mathbb{R}^n \times \mathbb{R}^n$ such that one of the following four equivalent conditions are met*

1. $-K^*q \in \partial G(u), \quad q \in \partial F(Ku),$
2. $-K^*q \in \partial G(u), \quad Ku \in \partial F^*(q),$
3. $u \in \partial G^*(-K^*q), \quad q \in \partial F(Ku),$
4. $u \in \partial G^*(-K^*q), \quad Ku \in \partial F^*(q),$

Then u solves the primal and q solves the dual optimization problem.

¹C.f. Rockafellar, *Convex Analysis*, Section 31

4.3 Applications of Duality in Vision

We covered in class how to solve problem with the descent algorithms that we know by looking at their dual.

The classic TV denoising model minimizes a combination of least-squares data term $G(u) = \frac{1}{2}\|u - f\|^2$ that penalizes the deviations of u to the input image f and the TV regularizer $F(Ku) = \alpha\|Ku\|_{2,1}$ that penalizes non-smooth images in terms of the discrete gradient operator K . To find the dual formulation of the problem, we substitute $F(Ku)$ by its biconjugate

$$F^{**}(Ku) = \max_q \langle Ku, q \rangle - \iota_{\|\cdot\|_{2,\infty} \leq 1}(q)$$

to obtain

$$\begin{aligned} \min_u \frac{1}{2}\|u - f\|_2^2 + \alpha\|Ku\|_{2,1} &= \min_u \max_{\|q\|_{2,\infty} \leq 1} \frac{1}{2}\|u - f\|_2^2 + \alpha\langle Ku, q \rangle \\ &= \max_{\|q\|_{2,\infty} \leq 1} \min_u \frac{1}{2}\|u - f\|_2^2 + \alpha\langle Ku, q \rangle \end{aligned}$$

Now the inner minimization problem obtains its optimum at $u = f - \alpha K^*q$ and the remaining problem in q becomes

$$\begin{aligned} \max_{\|q\|_{2,\infty} \leq 1} \frac{1}{2}\|f - \alpha K^*q - f\|_2^2 + \alpha\langle K(f - \alpha K^*q), q \rangle &= \max_{\|q\|_{2,\infty} \leq 1} \frac{1}{2}\|\alpha K^*q\|_2^2 + \alpha\langle Kf, q \rangle - \|\alpha D^*q\|_2^2 \\ &= \max_{\|q\|_{2,\infty} \leq 1} -\frac{1}{2}\|\alpha K^*q\|_2^2 + \alpha\langle Kf, q \rangle \\ &= \max_{\|q\|_{2,\infty} \leq 1} -\frac{1}{2}\|\alpha K^*q - f\|_2^2. \end{aligned}$$

Since we prefer minimizations over maximizations, we write

$$\begin{aligned} \hat{q} &= \operatorname{argmax}_{\|q\|_{2,\infty} \leq 1} -\frac{1}{2}\|\alpha K^*q - f\|_2^2 \\ \hat{q} &= \operatorname{argmin}_{\|q\|_{2,\infty} \leq 1} \frac{1}{2}\left\|K^*q - \frac{f}{\alpha}\right\|_2^2 \\ \hat{q} &= \operatorname{argmin}_{q \in C} \frac{1}{2}\left\|K^*q - \frac{f}{\alpha}\right\|_2^2 \quad \text{where } C = \{q \in \mathbb{R}^{nm \times 2c} \mid \|q\|_{2,\infty} \leq 1\} \end{aligned}$$

This is a minimization of a convex, proper, closed, L -smooth function over a convex set C that we can solve with gradient projection as follows:

$$q^{k+1} = \pi_C \left(q^k - \tau D \left(K^* q^k - \frac{f}{\alpha} \right) \right).$$

Working with the dual of a convex problem does not solve all of our problems and a simple modification to the TV denoising problem is enough to render our trick useless. Consider TV- ℓ^1 denoising model

$$\begin{aligned} &\inf_u \|u - f\|_1 + \alpha\|Du\|_{2,1} \\ &= \inf_u \sup_q \|u - f\|_1 + \alpha\langle q, Du \rangle - \delta_{\|\cdot\|_{2,\infty} \leq 1}(q) \\ &= \sup_q \inf_u \|u - f\|_1 + \alpha\langle q, Du \rangle - \delta_{\|\cdot\|_{2,\infty} \leq 1}(q) \\ &= \sup_q \left(-\sup_u \langle -\alpha D^*q, u \rangle - \|u - f\|_1 \right) - \delta_{\|\cdot\|_{2,\infty} \leq 1}(q) \\ &= \sup_q \langle \alpha D^*q, f \rangle - \delta_{\|\cdot\|_{\infty} \leq 1}(-\alpha D^*q) - \delta_{\|\cdot\|_{2,\infty} \leq 1}(q) \end{aligned}$$

The problem did not become easier because neither the primal nor the dual have simple proximal operators or satisfy all the smoothness assumptions that descent algorithms require. To overcome this situations, we will work directly with the saddle-point formulation of the problem.

4.4 Primal-Dual Algorithms

Primal-dual techniques solve minimization problems of the form

$$\min_u G(u) + F(Ku) \quad (4.11)$$

by solving the equivalent saddle-point problem

$$\min_u \max_p G(u) + \langle p, Ku \rangle - F^*(p). \quad (4.12)$$

The algorithm that we will study uses an interpretation of the proximal mapping as implicit gradient descent. This generalization results from the optimality conditions of the proximal update

$$\begin{aligned} u^{k+1} &= \text{prox}_{\tau E}(u^k) \\ u^{k+1} &= \underset{u}{\operatorname{argmin}} \tau E(u) + \frac{1}{2} \|u - u^k\|^2 \\ 0 &\in \tau \partial E(u^{k+1}) + u^{k+1} - u^k \\ u^{k+1} &\in u^k - \tau \partial E(u^{k+1}). \end{aligned} \quad (4.13)$$

If E is differentiable, we obtain an update rule

$$u^{k+1} = u^k - \tau \nabla E(u^{k+1})$$

that coincides to an implicit discretization of gradient descent, where the gradient is computed at the next iterate. The **proximal-point algorithm** uses this interpretation to generalize gradient descent to non-differentiable functions with a simple proximal operator.

We will use this algorithm to solve the saddle-point problem (4.12). The energy that we want to minimize over u and maximize over p is called the Lagrangian

$$\mathcal{L}(u, p) := G(u) + \langle p, Ku \rangle - F^*(p). \quad (4.14)$$

If we alternate implicit ascent steps in p with implicit descent steps in u , we obtain the following update rules:

$$\begin{aligned} p^{k+1} &= \text{prox}_{-\sigma \mathcal{L}(u^k, \cdot)}(p^k) \\ u^{k+1} &= \text{prox}_{\tau \mathcal{L}(\cdot, p^{k+1})}(u^k) \end{aligned}$$

We can exploit the structure of the Lagrangian, inherited from the decomposition of the energy as the sum of convex functions $G(u) + F(Ku)$, to formulate the dual update in terms of the conjugate of F as follows:

$$\begin{aligned} p^{k+1} &= \text{prox}_{-\sigma \mathcal{L}(u^k, \cdot)}(p^k), \\ &= \underset{p}{\operatorname{argmin}} \frac{1}{2} \|p - p^k\|^2 + \sigma F^*(p) - \sigma \langle Ku^k, p \rangle \\ &= \underset{p}{\operatorname{argmin}} \frac{1}{2} \|p - p^k - \sigma Ku^k\|^2 + \sigma F^*(p) \\ &= \text{prox}_{\sigma F^*}(p^k + \sigma Ku^k). \end{aligned}$$

Similarly, we can formulate the primal update in terms of the function G as follows:

$$\begin{aligned}
u^{k+1} &= \text{prox}_{\tau \mathcal{L}(\cdot, p^{k+1})}(u^k), \\
&= \underset{u}{\text{argmin}} \frac{1}{2} \|u - u^k\|^2 + \tau G(u) + \tau \langle Ku, p^{k+1} \rangle \\
&= \underset{u}{\text{argmin}} \frac{1}{2} \|u - u^k + \tau K^* p^{k+1}\|^2 + \tau G(u) \\
&= \text{prox}_{\tau G}(u^k - \tau K^* p^{k+1}).
\end{aligned}$$

This results in a very temping update rule

$$\begin{aligned}
p^{k+1} &= \text{prox}_{\sigma F^*}(p^k + \sigma K u^k), \\
u^{k+1} &= \text{prox}_{\tau G}(u^k - \tau K^* p^{k+1})
\end{aligned}$$

for a primal-dual algorithm. However, decoupling the updates of the primal and dual variables violates the optimality conditions of the saddle-point problem. To overcome this limitation, we need to substitute u^k in the dual update by an extrapolated variable \bar{u} . The result is the primal-dual hybrid gradient method.

Definition We will call the iteration

$$\begin{aligned}
p^{k+1} &= \text{prox}_{\sigma F^*}(p^k + \sigma K \bar{u}^k), \\
u^{k+1} &= \text{prox}_{\tau G}(u^k - \tau K^* p^{k+1}), \\
\bar{u}^{k+1} &= 2u^{k+1} - u^k.
\end{aligned} \tag{PDHG}$$

the **Primal-Dual Hybrid Gradient (PDHG) Method**.

As we will see, the Primal-Dual Hybrid Gradient method converges if $\tau\sigma < \frac{1}{\|K\|^2}$. It is actually easy to show that it converges to a minimizer of the original energy because a fixed-point (u^*, p^*, \bar{u}^*) of PDHG satisfies:

$$\begin{aligned}
\bar{u}^* &= u^*, \\
p^* &= \text{prox}_{\sigma F^*}(p^* + \sigma K u^*) \\
&= \underset{p}{\text{argmin}} \sigma F^*(p) + \frac{1}{2} \|p - p^* - \sigma K u^*\|^2 \iff 0 \in \sigma \partial F^*(p^*) - \sigma K u^*, \\
u^* &= \text{prox}_{\tau G}(u^* - \tau K^* p^*) \\
&= \underset{u}{\text{argmin}} \tau G(u) + \frac{1}{2} \|u - u^* + \tau K^* p^*\|^2 \iff 0 \in \tau \partial G(u^*) + \tau K^* p^*.
\end{aligned} \tag{4.15}$$

Combining these expressions we obtain the optimality conditions of Corollary 65

$$-K^* p^* \in \partial G(u^*) \qquad K \bar{u}^* = \partial F^*(p^*). \tag{4.16}$$

To analyze the convergence of the algorithm, it is useful to modify the order of the primal and dual updates to

$$\begin{aligned}
u^{k+1} &= \text{prox}_{\tau G}(u^k - \tau K^* p^k) \\
\bar{u}^{k+1} &= 2u^{k+1} - u^k \\
p^{k+1} &= \text{prox}_{\sigma F^*}(p^k + \sigma K \bar{u}^{k+1}).
\end{aligned} \tag{4.17}$$

and expand the proximal operators in terms of subgradients. For the primal variable, we have

$$\begin{aligned} u^{k+1} &= \text{prox}_{\tau G}(u^k - \tau K^* p^k) = \text{argmin } \tau G(u) + \frac{1}{2} \|u - u^k + \tau K^* p^k\|^2 \\ 0 &\in \partial G(u^{k+1}) + \frac{1}{\tau}(u^{k+1} - u^k) + K^* p^k \\ 0 &\in \partial G(u^{k+1}) - K^* p^{k+1} + \frac{1}{\tau}(u^{k+1} - u^k) + K^*(p^{k+1} - p^k). \end{aligned} \quad (4.18)$$

Similarly, for the dual variable

$$\begin{aligned} p^{k+1} &= \text{prox}_{\sigma F^*}(p^k + \sigma K(2u^{k+1} - u^k)) = \text{argmin } \sigma F^*(p) + \frac{1}{2} \|p - p^k - \sigma K(2u^{k+1} - u^k)\|^2 \\ 0 &\in \partial F^*(p^{k+1}) + \frac{1}{\sigma}(p^{k+1} - p^k) - K(2u^{k+1} - u^k) \\ 0 &\in \partial F^*(p^{k+1}) - K u^{k+1} + \frac{1}{\sigma}(p^{k+1} - p^k) - K(u^{k+1} - u^k). \end{aligned} \quad (4.19)$$

If we combine both expressions into a single update rule and define a set-valued operator T and matrix M as follows

$$\begin{pmatrix} 0 \\ 0 \end{pmatrix} \in \underbrace{\begin{pmatrix} \partial G & K^T \\ -K & \partial F^* \end{pmatrix}}_T \begin{pmatrix} u^{k+1} \\ p^{k+1} \end{pmatrix} + \underbrace{\begin{pmatrix} \frac{1}{\tau} I & -K^T \\ -K & \frac{1}{\sigma} I \end{pmatrix}}_M \begin{pmatrix} u^{k+1} - u^k \\ p^{k+1} - p^k \end{pmatrix}, \quad (4.20)$$

we can write the update rule of PDHG as the customized proximal-point algorithm

$$\begin{aligned} z^{k+1} &= (M + T)^{-1}(M z^k) \\ 0 &\in T(z^{k+1}) + M(z^{k+1} - z^k). \end{aligned} \quad (4.21)$$

The name originates from its similarity with the proximal point algorithm

$$\begin{aligned} u^{k+1} &= \text{prox}_E(u^k) = (I + \tau \partial E)^{-1}(u^k) \\ 0 &\in \partial E(u^{k+1}) + \frac{1}{\tau}(u^{k+1} - u^k). \end{aligned} \quad (4.22)$$

We will build on this algorithm to analyze the convergence of PDHG. In particular, as we showed convergence of the proximal point algorithm by showing that $\text{prox}_E = (I + \tau \partial E)^{-1}$ is firmly nonexpansive, thus averaged, and the Picard iteration converges; we will show convergence of PDHG by showing that T is maximal monotone, thus $(I + T)^{-1}$ is averaged, and the Picard iteration converges. Before we can do this, we need to define set-valued and monotone operators.

4.5 Monotone Operators

The mapping of a set-valued operator defines a relation. In monotone operator theory, a relation, a point-to-set mapping, a multi-valued function, and a correspondence all refer to the same concept.

Definition A relation or set-valued operator T on \mathbb{R}^n is a subset of $\mathbb{R}^n \times \mathbb{R}^n$.

For simplicity we write $T(x)$ and Tx to mean the set $\{y : (x, y) \in T\}$. This is an abuse of notation that is useful when T is a function and $T(x) = y$ is the singleton set $T(x) = \{y\}$.

For instance, the subdifferential of a function $E: \mathbb{R}^n \rightarrow \bar{\mathbb{R}}$ defines the subdifferential relation

$$\partial E = \{(u, g) \in \text{dom}(E) \times \mathbb{R}^n : E(v) \geq E(u) + \langle g, v - u \rangle \quad \forall v \in \mathbb{R}^n\}.$$

The set $\partial E(u)$ is a well-defined closed convex set at any point $u \in \text{dom}(E)$, but it can be empty. We have already seen that when E is convex, $\partial E(u) \neq \emptyset$ for any $u \in \text{ri}(\text{dom}(E))$.

We can extend many notions for functions to relations. For example, if T and S are relations, we define the domain $\text{dom}(T)$, composition TS and the sum $T + S$ as

$$\begin{aligned}\text{dom}(T) &= \{x \in \mathbb{R}^n : T(x) \neq \emptyset\} \\ TS &= \{(x, z) : \exists y \text{ s.t. } (x, y) \in S, (y, z) \in T\} \\ T + S &= \{(x, y + z) : (x, y) \in T, (x, z) \in S\}.\end{aligned}$$

More interesting for us are the inverse relation and the zeros of a relation.

Definition Inverse Relation. The inverse relation of T is defined as

$$T^{-1} = \{(x, y) : (y, x) \in T\}.$$

This always exists, even when T is a function that is not one-to-one. The inverse relation is not quite an inverse in the usual sense, as we can have $T^{-1}T \neq I$ like in the case of the zero function. In fact, we only have $T^{-1}Tx = x$ when T^{-1} is a function and $x \in \text{dom}(T)$.

Definition Zeros of a Relation. When $0 \in T(x)$, we say that x is a zero of T . The zero set of a relation T is

$$T^{-1}(\{0\}) = T^{-1}(0) = \{x : (x, 0) \in T\}.$$

Zeros and inverses are important for us when applied to the subdifferential. The **inverse of subdifferential** $(\partial E)^{-1}$ is defined as

$$\begin{aligned}(u, \hat{u}) \in (\partial E)^{-1} &\iff (\hat{u}, u) \in \partial E \\ &\iff u \in \partial E(\hat{u}) \\ &\iff 0 \in \partial E(\hat{u}) - u \\ &\iff \hat{u} \in \arg \min_v E(v) + \langle u, v \rangle\end{aligned}$$

Definition A relation T on \mathbb{R}^n is **monotone** if it satisfies

$$\langle u - v, x - y \rangle \geq 0 \quad \forall (x, u), (y, v) \in T$$

In multi-valued function notation, monotonicity reads

$$\langle Tx - Ty, x - y \rangle \geq 0 \quad \forall x, y \in \text{dom}(T),$$

where the left-hand side is a subset of \mathbb{R} and the inequality means that it lies in \mathbb{R}_+ .

Definition The relation T is **maximal monotone** if there is no monotone operator that properly contains it as a subset of $\mathbb{R}^n \times \mathbb{R}^n$.

In other words, if the monotone operator T is not maximal, then there is $(x, u) \notin T$ such that $T \cup \{(x, u)\}$ is monotone. We already know some examples of (maximally) monotone operators.

Lemma 66. $E: \mathbb{R}^n \rightarrow \bar{\mathbb{R}}$, then ∂E is a monotone operator. If E is closed convex and proper then ∂E is maximal monotone.

Proof. To prove monotonicity, we add the inequalities that define the subdifferential

$$\begin{aligned} E(y) &\geq E(x) + \langle \partial E(x), y - x \rangle \\ E(x) &\geq E(y) + \langle \partial E(y), x - y \rangle \end{aligned}$$

to obtain $\langle \partial E(x) - \partial E(y), x - y \rangle \geq 0$. This holds even when E is not convex. When E is convex, we need to show that ∂E is maximal, i.e., for any $(x, p) \notin \partial E$ there exists $(y, q) \in \partial E$ such that

$$\langle x - y, p - q \rangle < 0.$$

because $\partial E \cup (x, p)$ is not monotone. Let $y = \arg \min_z E(z) + \frac{1}{2}\|z - (x + p)\|^2$, then

$$\begin{aligned} 0 &\in \partial E(y) + y - x - p \\ x - y &\in q - p \quad \text{for } q \in \partial E(y). \end{aligned}$$

As $(x, p) \notin \partial E$, either $x \neq y$ or $p \neq q$ and we have

$$\langle x - y, p - q \rangle = -\|x - y\|^2 = -\|p - q\|^2 < 0.$$

□

A relation on \mathbb{R} is monotone if it is a curve in \mathbb{R}^2 that is always nondecreasing (it can have flat and vertical portions). If it is a continuous curve with no end points, then it is maximal monotone.

Lemma 67. *A continuous monotone function $F: \mathbb{R}^n \rightarrow \mathbb{R}$ with $\text{dom}(F) = \mathbb{R}^n$ is maximal.*

Proof. Assume by contradiction that there is a pair $(y, v) \notin F$ such that

$$\langle v - F(x), y - x \rangle \geq 0 \quad \forall x \in \text{dom}(F) = \mathbb{R}^n.$$

Given any $z \in \mathbb{R}^n$, the previous inequality should hold true for $x = y - t(z - y)$ for all $t > 0$, that is

$$\begin{aligned} \langle v - F(y - t(z - y)), t(z - y) \rangle &\geq 0 \\ \langle v - F(y - t(z - y)), z - y \rangle &\geq 0. \end{aligned} \tag{4.23}$$

As F is continuous, when $t \rightarrow 0$ we have

$$\langle v - F(y), z - y \rangle \geq 0.$$

As this should hold for an arbitrary z , we have $v = F(y)$ and we obtain the contradiction $(y, v) \in F$. □

Definition The **resolvent** of a relation T on \mathbb{R}^n is

$$R_T = (I + \alpha T)^{-1},$$

where $\alpha \in \mathbb{R}$. The **Cayley operator**, reflection operator, or reflected resolvent of T is defined as

$$C_T = 2R_T - I,$$

where I is the identity function.

Lemma 68. *If T is maximal monotone, then the resolvent $R_T = (I + \alpha T)^{-1}$ with $\alpha > 0$ and the Cayley operator $C_T = 2R_T - I$ are nonexpansive functions.*

Proof. We show first that R_T is nonexpansive. Suppose $(x, u) \in R_T$ and $(y, v) \in R_T$. By definition of R_T , we have

$$(x, u) \in R_T \tag{4.24}$$

$$(u, x) \in (I + \alpha T) \tag{4.25}$$

$$x \in u + \alpha T(u). \tag{4.26}$$

Similarly $y \in v + \alpha T(v)$ and we can combine both into

$$x - y \in u - v + \alpha(T(u) - T(v)). \tag{4.27}$$

Multiplying both sides by $u - v$ and using the monotonicity of T , we get

$$\langle u - v, x - y \rangle = \|u - v\|^2 + \alpha \langle u - v, T(u) - T(v) \rangle \tag{4.28}$$

$$\langle u - v, x - y \rangle \geq \|u - v\|^2. \tag{4.29}$$

Now we apply Cauchy-Schwarz inequality and divide by $\|u - v\|$ to get

$$\|x - y\| \geq \|u - v\| = \|R_T(x) - R_T(y)\| \tag{4.30}$$

which shows that R_T is nonexpansive.

Next, we show that $C_T = 2R_T - I$ is also nonexpansive. Using the inequality (4.29), we get

$$\begin{aligned} \|C_T(x) - C_T(y)\|^2 &= \|2(u - v) - (x - y)\|^2 \\ &= 4\|u - v\|^2 - 4\langle x - y, u - v \rangle + \|x - y\|^2 \\ &\leq \|x - y\|^2, \end{aligned} \tag{4.31}$$

which shows that C_T is nonexpansive. \square

The following properties describe how we can combine monotone operators without losing the monotonicity.

Lemma 69. *If F and G are (maximal) monotone and $\text{dom}(F) \cap \text{dom}(G) \neq \emptyset$, then $F + G$ is (maximal) monotone.*

Theorem 70. Convergence proximal point algorithm *Let T be a maximal monotone operator, and let there exist a z such that $0 \in T(z)$. Then the (generalized) proximal point algorithm*

$$\begin{aligned} z^{k+1} &= (T + I)^{-1}(z^k) \\ 0 &\in T(z^{k+1}) + z^{k+1} - z^k \end{aligned} \tag{4.32}$$

converges to a point \tilde{z} with $0 \in T(\tilde{z})$.

Proof. From Lemma 68 we know that if T is maximal monotone, then the resolvent $R_T = (T + I)^{-1}$ and the Caley operator $C_T = 2R_T - I$ are nonexpansive. Since $R_T = \frac{1}{2}I + \frac{1}{2}C_T$, the resolvent R_T is an averaged operator and the generalized proximal point algorithm

$$z^{k+1} = R_T(z^k) \tag{4.33}$$

is a fixed-point iteration of an averaged operator that converges by Krasnoselskii-Mann Theorem. \square

To apply this theorem to the PDHG algorithm

$$0 \in T(z^{k+1}) + Mz^{k+1} - Mz^k, \quad (4.34)$$

we need to factor the matrix M . If M is symmetric positive definite, we can decompose it as $M = L^T L$ and define the operator $\tilde{T} = L^{-T} T L^{-1}$ whose fixed-point iteration

$$\begin{aligned} z^{k+1} &= (\tilde{T} + I)^{-1}(z^k) \\ 0 &\in \tilde{T}(z^{k+1}) + z^{k+1} - z^k \\ 0 &\in L^{-T} T L^{-1}(z^{k+1}) + z^{k+1} - z^k \\ 0 &\in T L^{-1}(z^{k+1}) + L^T(z^{k+1} - z^k) \\ 0 &\in T(L^{-1}z^{k+1}) + L^T L(L^{-1}z^{k+1} - L^{-1}z^k) \\ 0 &\in T(x^{k+1}) + L^T L(x^{k+1} - x^k) \\ 0 &\in T(x^{k+1}) + M(x^{k+1} - x^k), \end{aligned} \quad (4.35)$$

coincides with the PDHG algorithm update. As a result, we can show convergence of the PDHG algorithm by determining the conditions under which the set-valued operator \tilde{T} is maximal monotone, and the matrix M is symmetric positive definite.

The next lemma shows that if T is maximal monotone, then \tilde{T} is also maximal monotone.

Lemma 71. *If T is (maximal) monotone, then $L^{-T} T L^{-1}$ is (maximal) monotone, too.*

Proof. To show that $L^{-T} T L^{-1}$ is monotone, we need to show the inequality

$$\langle q_u - q_v, u - v \rangle \geq 0$$

holds for all $u, v, q_u \in L^{-T} T L^{-1}(u)$ and $q_v \in L^{-T} T L^{-1}(v)$. That is, for all $u, v, p_u \in T L^{-1}(u)$ and $p_v \in T L^{-1}(v)$

$$\begin{aligned} \langle L^{-T} p_u - L^{-T} p_v, u - v \rangle &\geq 0 \\ \langle p_u - p_v, L^{-1} u - L^{-1} v \rangle &\geq 0. \end{aligned}$$

If we now define $\bar{u} = L^{-1} u, \bar{v} = L^{-1} v$ we have the following inequality

$$\langle p_u - p_v, \bar{u} - \bar{v} \rangle \geq 0$$

for all $\bar{u}, \bar{v}, p_u \in T(\bar{u})$ and $p_v \in T(\bar{v})$ which is a result of the monotonicity of T . The maximal monotonicity can be proven easily by contradiction. \square

We now need to show that T is maximal monotone.

Lemma 72. *Let $G : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{\infty\}$ and $F : \mathbb{R}^m \rightarrow \mathbb{R} \cup \{\infty\}$ be proper, closed, convex functions and matrix $K \in \mathbb{R}^{m \times n}$, then the set valued operator*

$$T = \begin{pmatrix} \partial G & K^T \\ -K & \partial F^* \end{pmatrix} = \begin{pmatrix} \partial G & 0 \\ 0 & \partial F^* \end{pmatrix} + \begin{pmatrix} 0 & K^T \\ -K & 0 \end{pmatrix}.$$

is maximal monotone

Proof. Let $x = (u, p) \in T \in \mathbb{R}^n \times \mathbb{R}^m$ and $y = (v, q) \in T \in \mathbb{R}^n \times \mathbb{R}^m$, by bilinearity of the scalar product we have

$$\begin{aligned}
\langle x - y, Tx - Ty \rangle &= \left\langle \begin{pmatrix} u - v \\ p - q \end{pmatrix}, \begin{pmatrix} \partial G(u) + K^\top p \\ -Ku + \partial F^*(p) \end{pmatrix} - \begin{pmatrix} \partial G(v) + K^\top q \\ -Kv + \partial F^*(q) \end{pmatrix} \right\rangle \\
&= \left\langle \begin{pmatrix} u - v \\ p - q \end{pmatrix}, \begin{pmatrix} \partial G(u) - \partial G(v) + K^\top(p - q) \\ -K(u - v) + \partial F^*(p) - \partial F^*(q) \end{pmatrix} \right\rangle \\
&= \langle u - v, \partial G(u) - \partial G(v) + K^\top(p - q) \rangle + \langle p - q, -K(u - v) + \partial F^*(p) - \partial F^*(q) \rangle \\
&= \langle u - v, \partial G(u) - \partial G(v) \rangle + \langle u - v, K^\top(p - q) \rangle - \langle p - q, K(u - v) \rangle + \langle p - q, \partial F^*(p) - \partial F^*(q) \rangle \\
&= \langle u - v, \partial G(u) - \partial G(v) \rangle + \langle p - q, \partial F^*(p) - \partial F^*(q) \rangle. \tag{4.36}
\end{aligned}$$

As F and G are proper, closed, and convex, their subdifferentials are maximal monotone operators (Lemma 68), that is,

$$\begin{aligned}
\langle u - v, \partial G(u) - \partial G(v) \rangle &\geq 0 \\
\langle p - q, \partial F^*(p) - \partial F^*(q) \rangle &\geq 0. \tag{4.37}
\end{aligned}$$

Adding both inequalities we prove that T is monotone because

$$\langle x - y, Tx - Ty \rangle = \langle u - v, \partial G(u) - \partial G(v) \rangle + \langle p - q, \partial F^*(p) - \partial F^*(q) \rangle \geq 0. \tag{4.38}$$

To show that T is maximal, let us write it as the sum of maximal monotone operators: the subdifferential of convex, proper, closed functions G and F^* , that are maximal monotone by Lemma 68, and an affine a function that is monotone by Lemma 71.

$$T = \begin{pmatrix} \partial G & K^\top \\ -K & \partial F^* \end{pmatrix} = \begin{pmatrix} \partial G & 0 \\ 0 & \partial F^* \end{pmatrix} + \begin{pmatrix} 0 & K^\top \\ -K & 0 \end{pmatrix}. \tag{4.39}$$

□

By construction M is already symmetric, if $\tau\sigma < \frac{1}{\|K\|^2}$ it is also positive definite. We can thus summarize the previous result in the following theorem.

Theorem 73. Convergence PDHG *Let F and G be proper, closed, and convex, and assume that the problem*

$$\min_{u \in \mathbb{R}^n} G(u) + F(Ku) \tag{4.40}$$

has a solution and there exists $u \in \text{ri}(G)$ such that $Ku \in \text{ri}(F)$. Then, the operator

$$T = \begin{pmatrix} \partial G & K^\top \\ -K & \partial F^* \end{pmatrix}$$

is maximally monotone. For $\tau\sigma < \frac{1}{\|K\|^2}$ the matrix

$$M = \begin{pmatrix} \frac{1}{\tau}I & -K^\top \\ -K & \frac{1}{\sigma}I \end{pmatrix}$$

in the PDHG algorithm is positive definite and PDHG converges.

4.6 Applications of Primal-Dual Algorithms in Vision

Let K being a discretization of the multichannel gradient operator, we will use the saddle-point formulation to solve the following problems:

The classic ROF denoising model minimizes a least-squares data term $G(u) = \frac{1}{2}\|u - f\|^2$ and the TV regularizer $F(Ku) = \alpha\|Ku\|_{2,1}$ as follows:

$$\min_u P(u) = \min_u \frac{1}{2}\|u - f\|^2 + \alpha\|Ku\|_{2,1} = \min_u \max_p \frac{1}{2}\|u - f\|^2 + \langle Ku, p \rangle - \iota_{\|\cdot\|_{2,\infty} \leq \alpha}(p).$$

As the proximal steps of both G and F^* are simple, the PDHG updates

$$\begin{aligned} p^{k+1} &= \operatorname{argmin}_p \frac{1}{2}\|p - (p^k + \sigma K \bar{u}^k)\|^2 + \sigma \iota_{\|\cdot\|_{2,\infty} \leq \alpha}(p) \\ u^{k+1} &= \operatorname{argmin}_u \frac{1}{2}\|u - (u^k - \tau K^* p^{k+1})\|^2 + \frac{\tau}{2}\|u - f\|^2 = \frac{u^k - \tau K^* p^{k+1} + \tau f}{1 + \tau} \\ \bar{u}^{k+1} &= 2u^{k+1} - u^k. \end{aligned}$$

are extremely simple to code. They are also efficient and result in an algorithm that parallelizes well.

The TV- L^1 denoising problem keeps the TV regularizer $F(Ku) = \alpha\|Ku\|_{2,1}$ but changes the norm used for the data term G to

$$\min_u P(u) = \min_u \|u - f\|_1 + \alpha\|Ku\|_{2,1} = \min_u \max_p \frac{1}{2}\|u - f\|_1 + \langle Ku, p \rangle - \iota_{\|\cdot\|_{2,\infty} \leq \alpha}(p).$$

as a result, only the proximal step of the primal variable appearing in G changes. In particular, we have

$$\begin{aligned} p^{k+1} &= \operatorname{argmin}_p \frac{1}{2}\|p - (p^k + \sigma K \bar{u}^k)\|^2 + \sigma \iota_{\|\cdot\|_{2,\infty} \leq \alpha}(p) \\ u^{k+1} &= \operatorname{argmin}_u \frac{1}{2}\|u - (u^k - \tau K^* p^{k+1})\|^2 + \tau \|u - f\|_1 \\ \Rightarrow u_i^{k+1} &= f_i + \operatorname{sign}(v_i^k) \max(|v_i^k| - \tau, 0) \\ \bar{u}^{k+1} &= 2u^{k+1} - u^k. \end{aligned}$$

The TV deblurring problem is more interesting because it admits multiple formulations as a saddle-point problem. Given A the convolution operator with a blur kernel, the TV deblurring problem reads

$$\min_u P(u) = \min_u \frac{1}{2}\|Au - f\|^2 + \alpha\|Ku\|_{2,1}.$$

To derive the first PDHG formulation of the problem, we substitute the TV regularizer by its biconjugate as we have done before, that is

$$\min_u P(u) = \min_u \max_p \frac{1}{2}\|Au - f\|^2 + \langle Ku, p \rangle - \iota_{\|\cdot\|_{2,\infty} \leq \alpha}(p).$$

and obtain the following PDHG updates

$$\begin{aligned} p^{k+1} &= \operatorname{argmin}_p \frac{1}{2}\|p - (p^k + \sigma K \bar{u}^k)\|^2 + \sigma \iota_{\|\cdot\|_{2,\infty} \leq \alpha}(p) \\ u^{k+1} &= \operatorname{argmin}_u \frac{1}{2}\|u - (u^k - \tau K^* p^{k+1})\|^2 + \frac{\tau}{2}\|Au - f\|^2 \\ &= (I + \tau A^* A)^{-1}(u^k - \tau K^* p^{k+1} + \tau f) \\ \bar{u}^{k+1} &= 2u^{k+1} - u^k. \end{aligned}$$

In the second case, we introduce an additional dual variable q for the blurred image Au and substitute the data term $\frac{1}{2}\|Au - f\|^2$ by its biconjugate:

$$\begin{aligned} \min_u P(u) &= \min_u \max_{p,q} \langle Au - f, q \rangle - \frac{1}{2}\|q\|^2 + \langle Ku, p \rangle - \iota_{\|\cdot\|_2, \infty \leq \alpha}(p) \\ &= \min_u \max_{p,q} \left\langle \begin{pmatrix} A \\ K \end{pmatrix} u, \begin{pmatrix} q \\ p \end{pmatrix} \right\rangle - \langle f, q \rangle - \frac{1}{2}\|q\|^2 - \iota_{\|\cdot\|_2, \infty \leq \alpha}(p). \end{aligned}$$

Now we have

$$F^*(p, q) = \langle f, q \rangle + \frac{1}{2}\|q\|^2 + \iota_{\|\cdot\|_2, \infty \leq \alpha}(p) \quad G(u) = 0 \quad \tilde{K} = \begin{pmatrix} A \\ K \end{pmatrix}$$

and the PDHG updates are simpler because $F^*(p, q)$ decomposes into independent minimizations in p and q but involve an additional variable

$$\begin{aligned} q^{k+1} &= \operatorname{argmin}_q \frac{1}{2}\|q - (q^k + \sigma A \bar{u}^k)\|^2 + \sigma \langle f, q \rangle + \frac{\sigma}{2}\|q\|^2 \\ p^{k+1} &= \operatorname{argmin}_p \frac{1}{2}\|p - (p^k + \sigma K \bar{u}^k)\|^2 + \sigma \iota_{\|\cdot\|_2, \infty \leq \alpha}(p) \\ u^{k+1} &= u^k - \tau K^* p^{k+1} - \tau A^* q^{k+1} \\ \bar{u}^{k+1} &= 2u^{k+1} - u^k. \end{aligned}$$

This is a common situation: different formulations of the dual or saddle-point problem exist for the same primal problem depending on how the terms in the objective functional are split, and the primal and dual variables are defined. Usually this different formulations result in primal-dual updates with different computational or memory cost. A common case finds two different formulations that trade off memory for computations. The right choice is a matter of the resources and constraints allocated to the problem.